ReNets: Statically-Optimal Demand-Aware Networks

Chen Avin and Stefan Schmid

"We cannot direct the wind, but we can adjust the sails."

Acknowledgements:

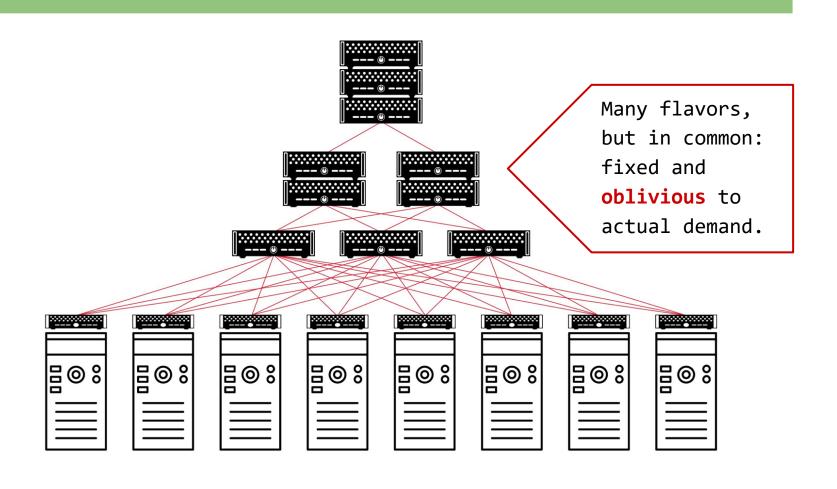






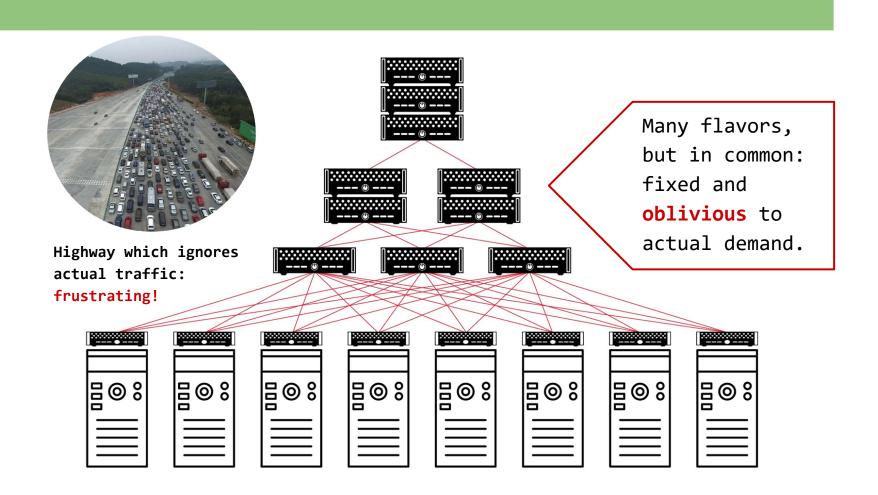
Today's Datacenters

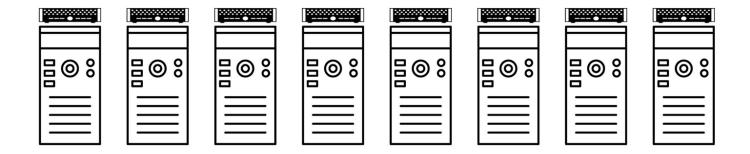
Fixed and Demand-Oblivious Topology

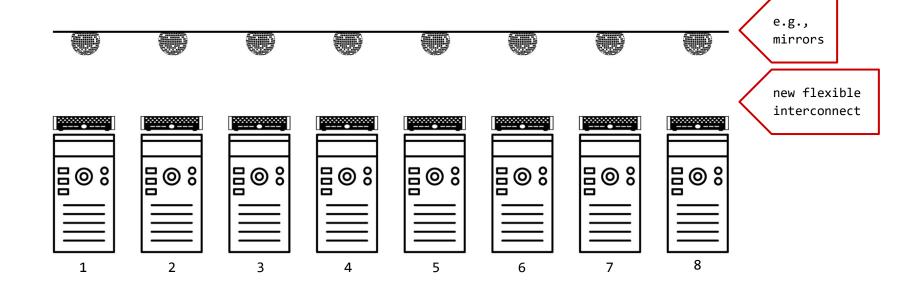


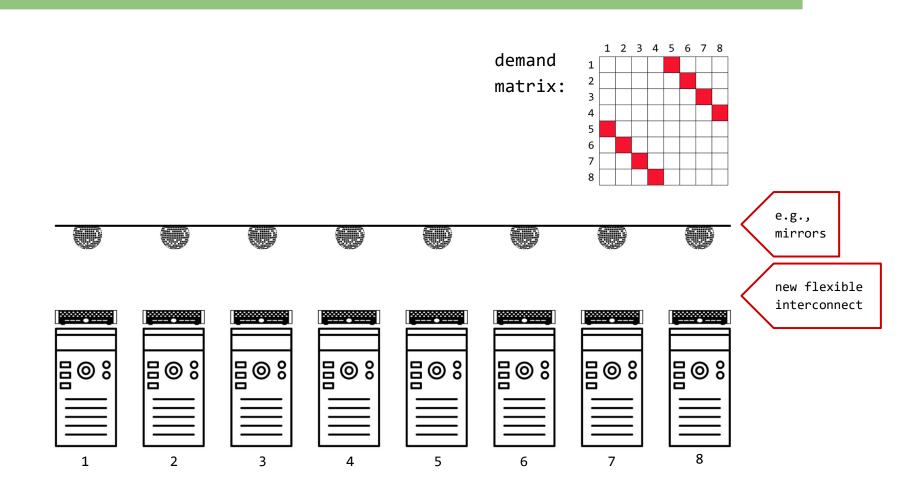
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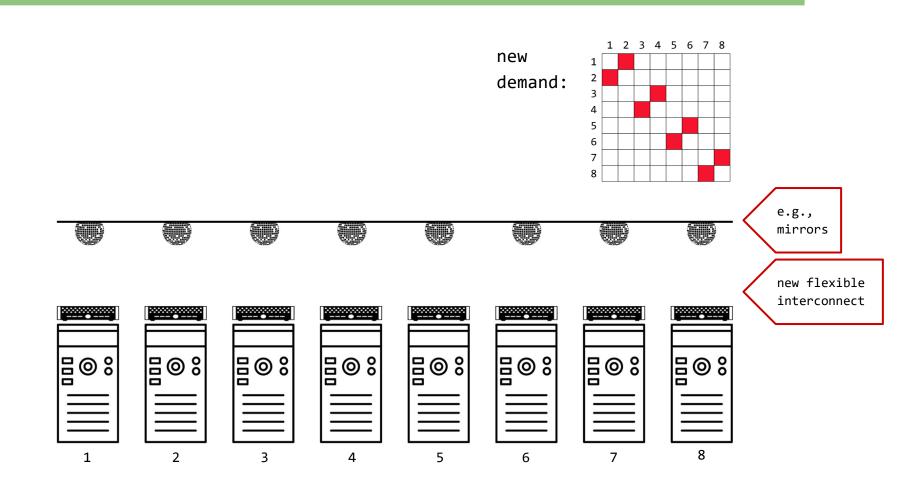


Flexible and Demand-Aware Topologies

2 matrix: 4 Matches demand 5 e.g., mirrors new flexible interconnect **□**◎°

demand

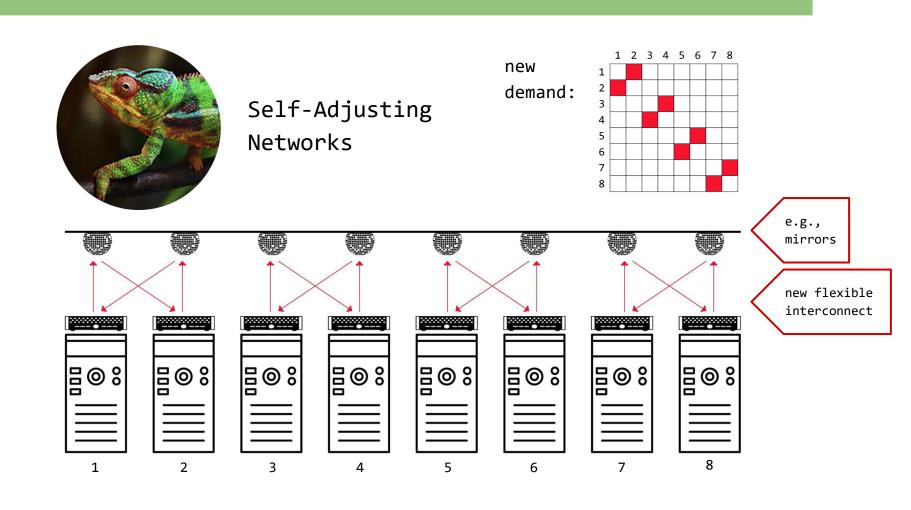
1 2 3 4 5 6 7 8



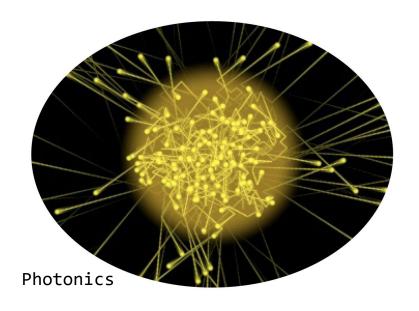
Flexible and Demand-Aware Topologies

new demand: Matches demand 5 e.g., mirrors new flexible interconnect **□**◎°

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Sounds Crazy? Emerging Enabling Technology.



H2020:

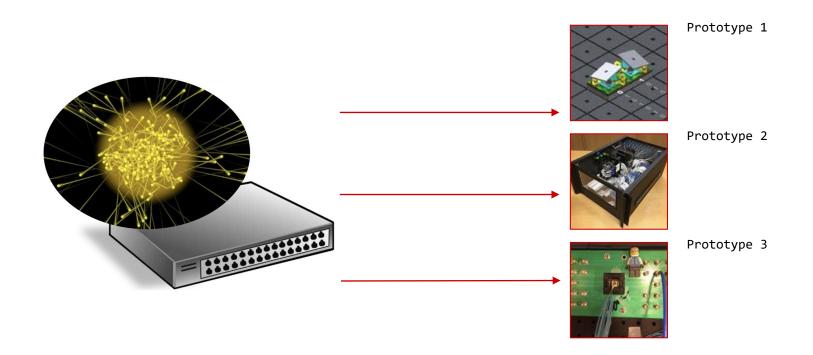
"Photonics one of only five key enabling technologies for future prosperity."

US National Research Council: "Photons are the new Electrons."

Enabler

Novel Reconfigurable Optical Switches

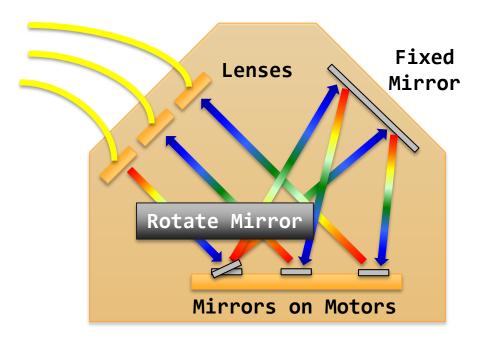
- → Spectrum of prototypes
 - → Different sizes, different reconfiguration times
 - → From our recent ACM SIGCOMM OptSys'19 workshop



Example

Optical Circuit Switch

Optical Circuit Switch rapid adaption of physical layer
→ based on rotating mirrors



Optical Circuit Switch

By Nathan Farrington, SIGCOMM 2010

Empirical Motivation

Temporal and Spatial Structure



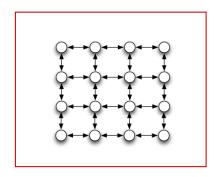
Source: Avin et al., On the Complexity of Traffic Traces and Implications, SIGMETRICS 2020

The Potential

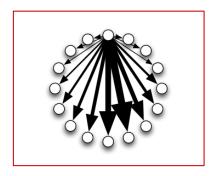
Example: Expected Route Length

- ---> Expected path length: number of hops times demand
- ---> Consider design of constant degree topologies (e.g., 4)
- Note: diameter is at least logarithmic

Demand 1: Low Degree



Demand 2: Skewed

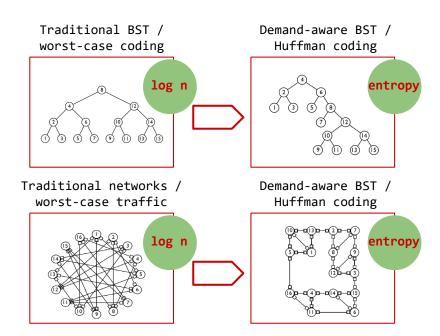


Expected route length in demand-aware network is constant in these cases (while diameter is $\Omega(\log n)$).

Connection to Entropy

→ Achievable expected route length is proportional to conditional entropy of the demand matrix

---> Similar to coding and data structures:



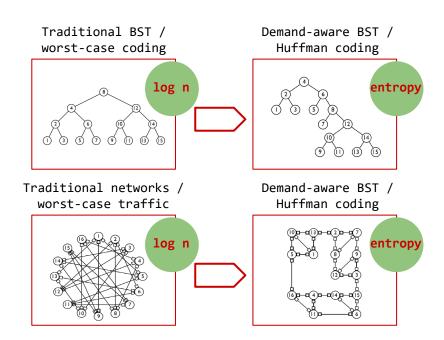


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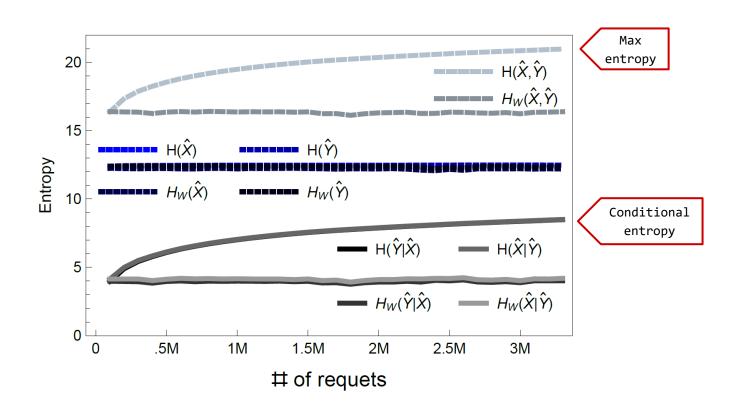




But how to achieve short routes if the demand is not known ahead of time and we have to account for reconfiguration costs?

Example

Facebook's Datacenter Traces

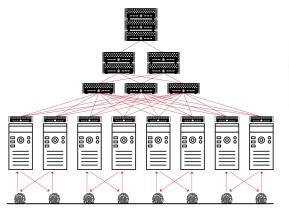


Our Contribution:

ReNet, A Statically Optimal Demand-Aware Network

→ Model: hybrid architecture

- → Fixed network of diameter log n
 plus reconfigurable network
 (constant number of direct links)
- → Segregated routing
- → **Online** sequence of requests: σ = (σ1, σ2, σ3, ...)
- → Global controller





reconfigurable



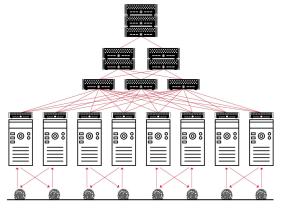
- Objective: Minimize route length
 plus reconfigurations
 - → More specifically: be statically optimal
 - \rightarrow Compared to a fixed algorithm which knows σ ahead of time

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reconfigurable



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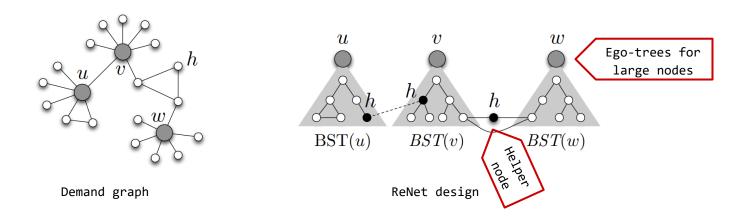


- → Compact routing (constant tables)
- → Local routing (greedy)
- → Arbitrary addressing

The ReNet Algorithm (1)

Algorithmic building blocks:

- 1. Working Set (WS)
 - \rightarrow Nodes keep track of recent communication partners in σ .
- 2. Small/large nodes and Ego-Tree
 - → Nodes with small WS connect to WS directly, nodes with large WS via a self-adjusting binary search tree (e.g., a splay tree)
- 3. Helper nodes to reduce the degree
 - → Large nodes may appear in many ego-trees, so get help of small nodes

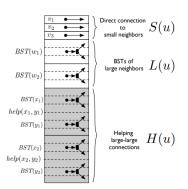


The ReNet Algorithm (2)

Continued:

4. Self adjustments

- → Keep track of WS; when too large: flush-when-full
- 5. Centralized coordination
 - → Fairly **decentralized**: coordinator only needs to keep track of which nodes are large and which small
 - → Nodes inform coordinator when adding node to working set
 - → Coordinator then assigns helper node on demand



Analytical Results (1)

Theorem 1:

For any sparse communication sequence of a certain length, ReNets are statically optimal while ensuring a bounded degree.

- Sparse: subsequences of only involve a linear number of nodes
- → Required to ensure availability of helper nodes (DISC 2017)

Analytical Results (2)

Theorem 2:

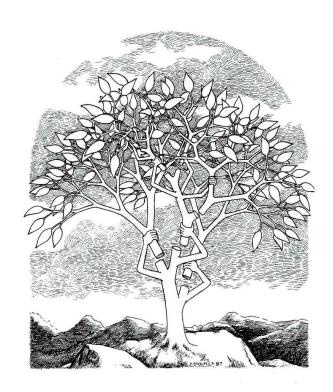
Under certain communication patterns, the amortized cost of ReNet can be significantly lower than the static optimum, i.e., $\Omega(\log n)$.

- Example: consider sequence of $\sigma = (\sigma^{(1)}, \sigma^{(2)}, \sigma^{(3)}, \ldots)$ where each $\sigma^{(i)}$ is of length n log n, sparse and corresponds to different 2-dimensional grid.
- In this example, the cost of ReNet is **constant** for each $\sigma^{(i)}$.
- \longrightarrow Overall, the union of the grids form a uniform pattern, so the cost of the static algorithm is $\log n$ (for constant degree).

Conclusion

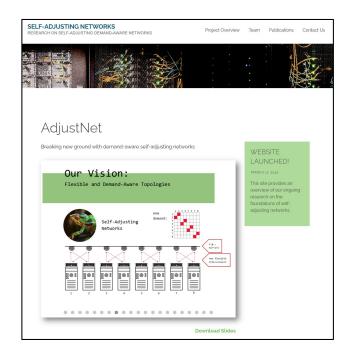
- → ReNet: statically optimal and
 - → compact routing
 - → local routing
 - → arbitrary addressing
- ---> Avenues for future work
 - → dense communication
 - → dynamic optimality

Thank you!



A Self-Adjusting Search Tree by Jorge Stolfi (1987)

Websites



http://self-adjusting.net/ Project website



https://trace-collection.net/ Trace collection website

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