#### Distributed Cloud Market: Who Benefits from Specification Flexibilities?

Arne Ludwig, Stefan Schmid



- 30% CAGR cloud 2013-2018
  - $\rightarrow$  Plethora of cloud providers



Google Cloud Platform

Microsoft

Azure

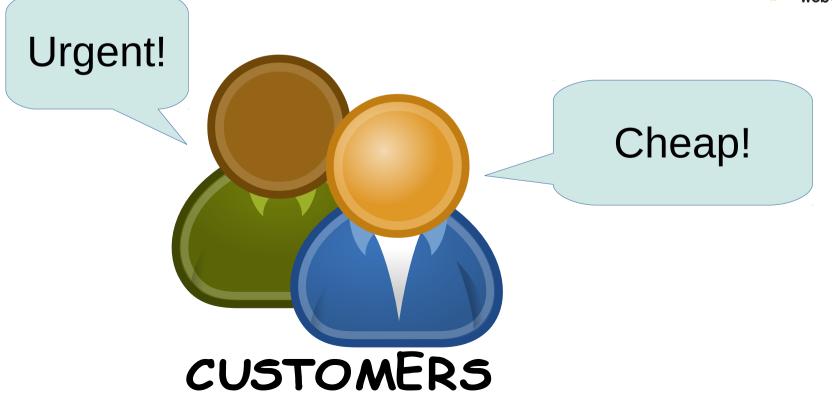


• Customer specifications?

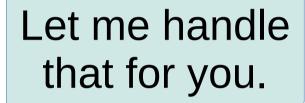


• Customer specifications?

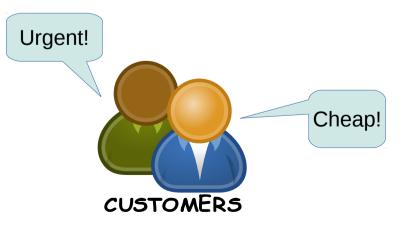


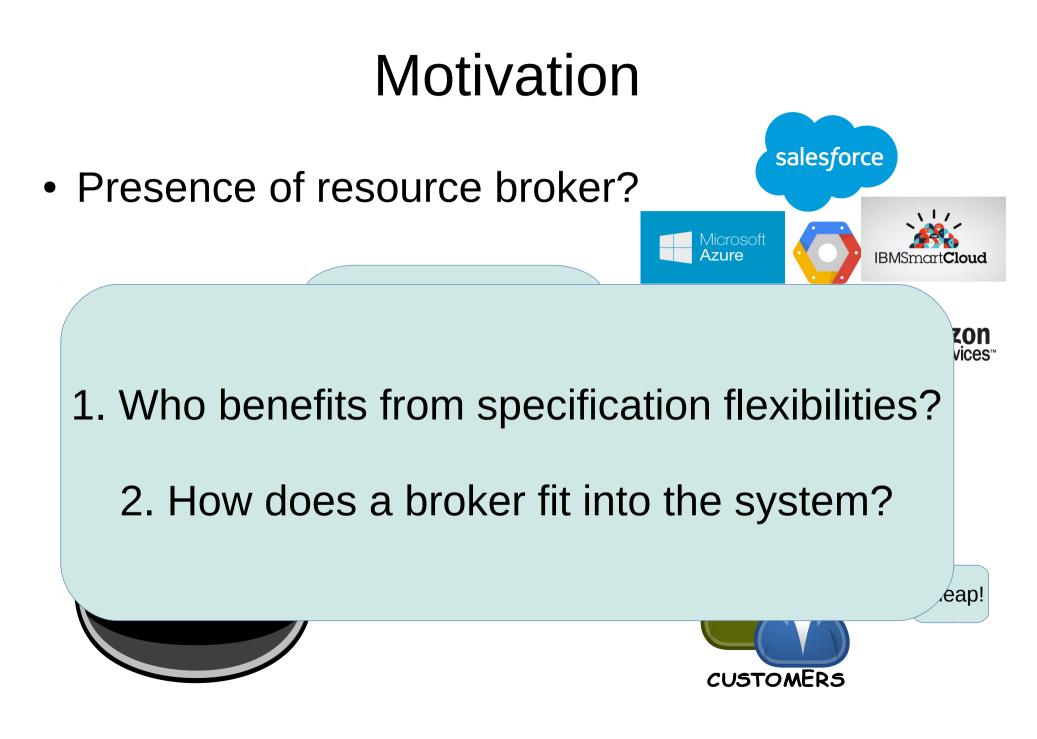


• Presence of resource broker?



Salesforce Microsoft Azure Microsoft DMSmcrtCloud BMSmcrtCloud BMSmcrtCloud





### Overview

- 1) Horizontal market
- 2) Vertical market (broker)
- 3) Resource buying strategies

## Model

- 3 types of customer
  - Flexible in time
  - Flexible in price (or inflexible in time)
  - Standard
- Provider fixed prices P1 < P2 < P3
- Virtual network requests

 $\rightarrow$  SLAs on server and network resources

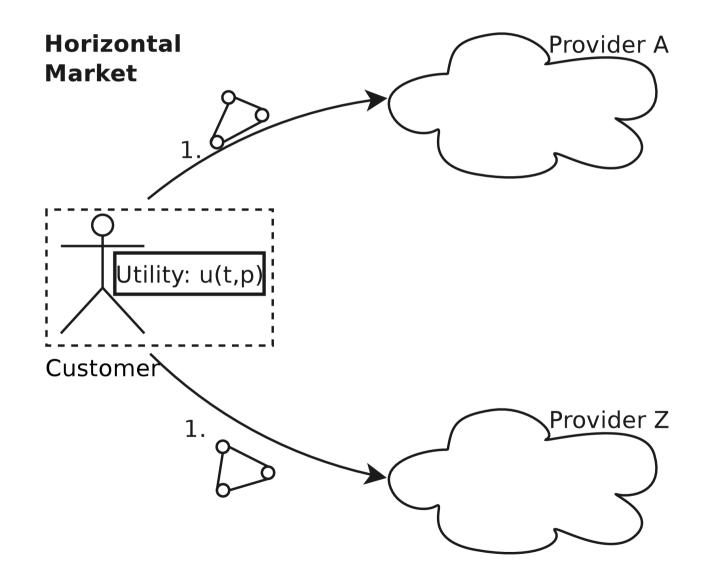
## Model

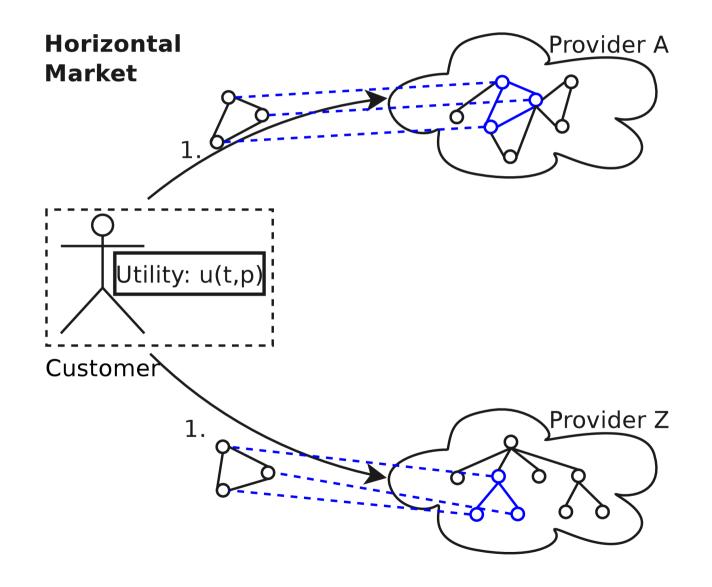
- 3 types of customer
  - Flexible in time  $u(t,p) = -t x \cdot p$
  - Flexible in price (inflexible in time)  $u(t,p) = -x \cdot t p$
  - Standard

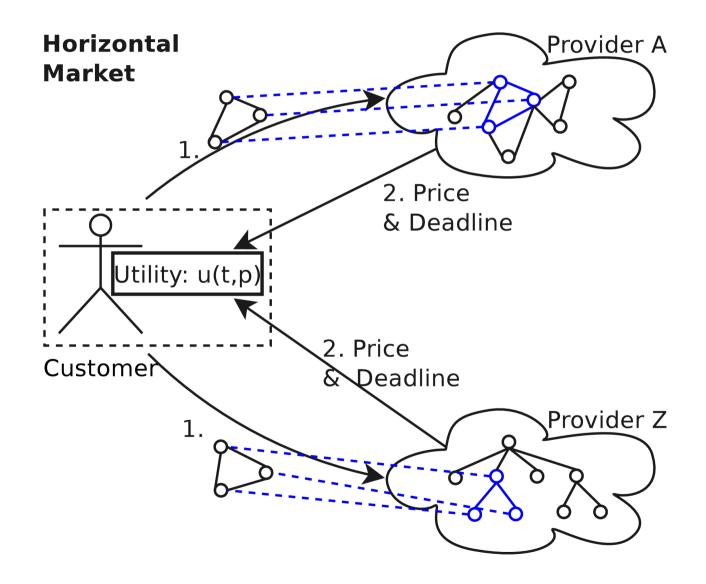
$$u(t,p) = -t - p$$

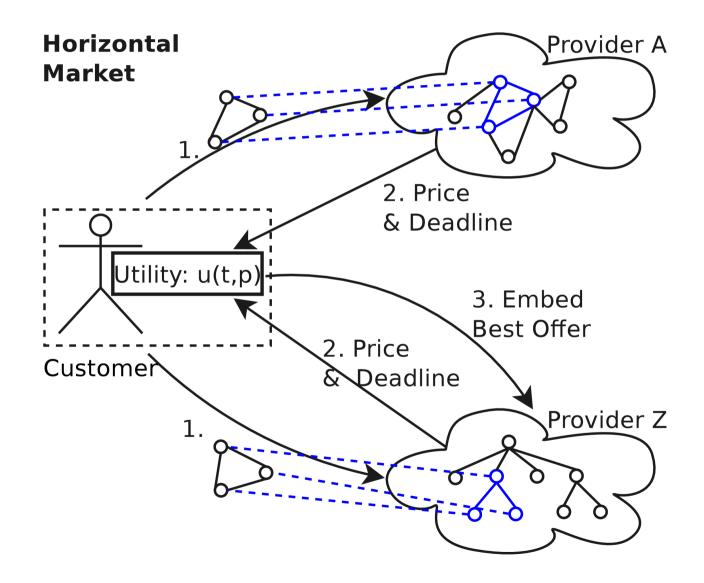
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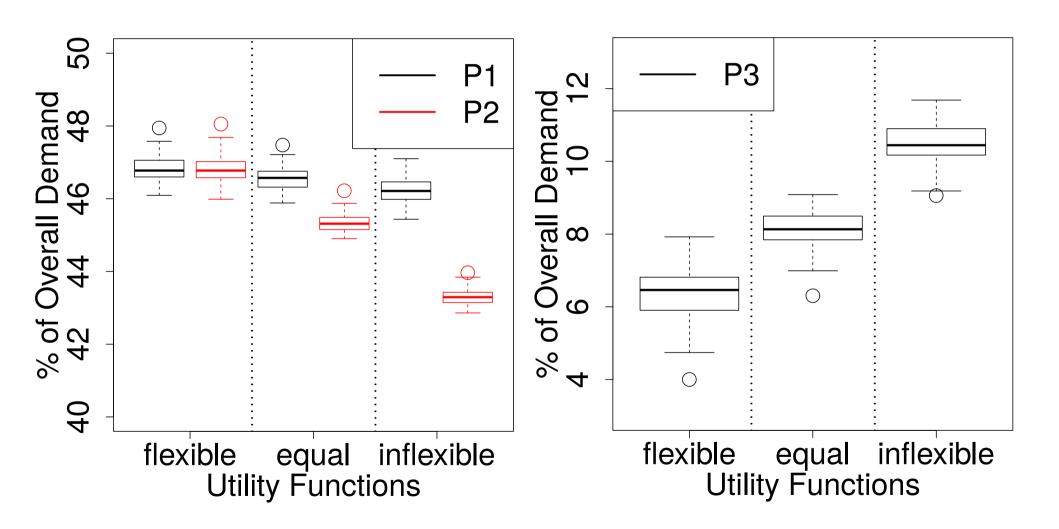




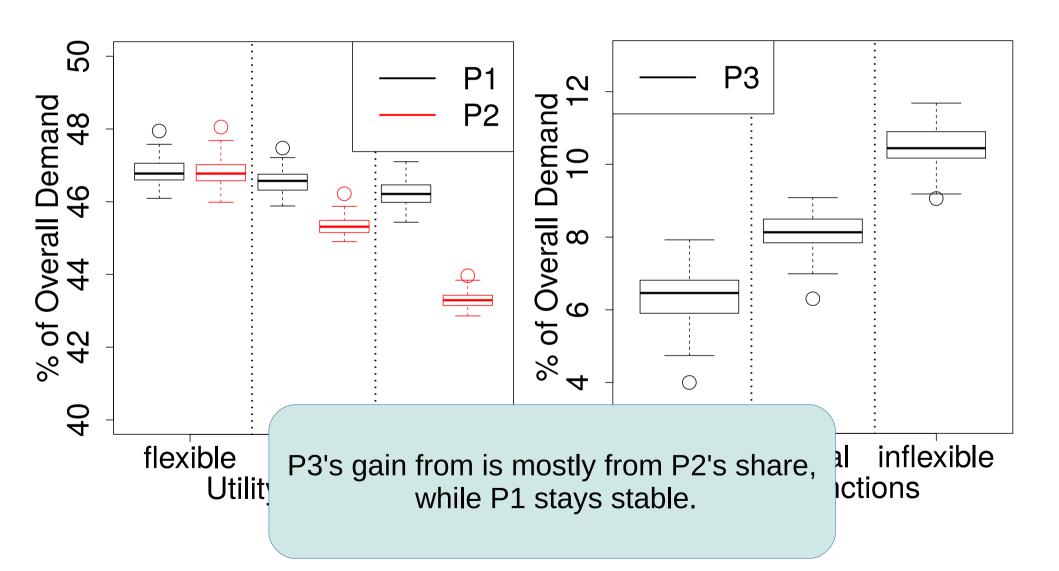


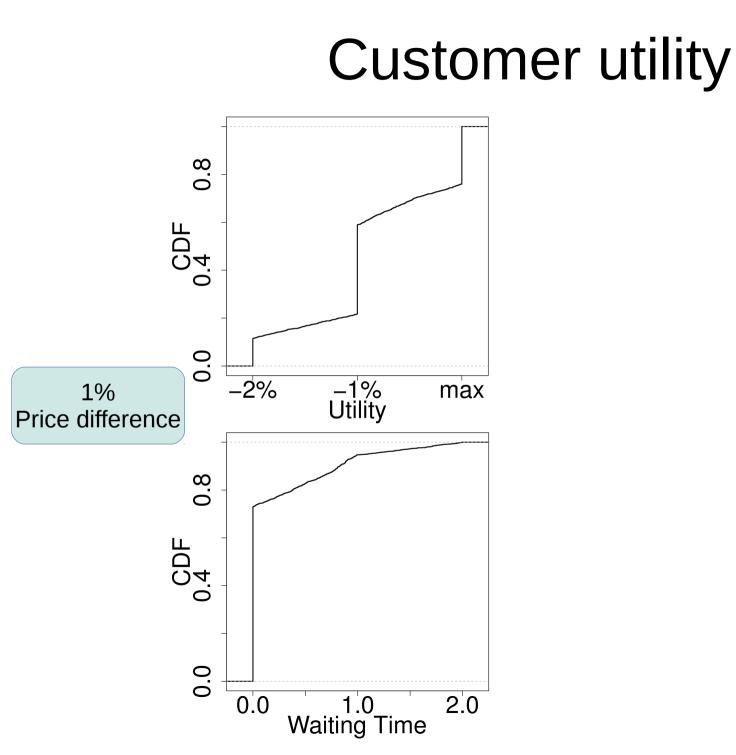


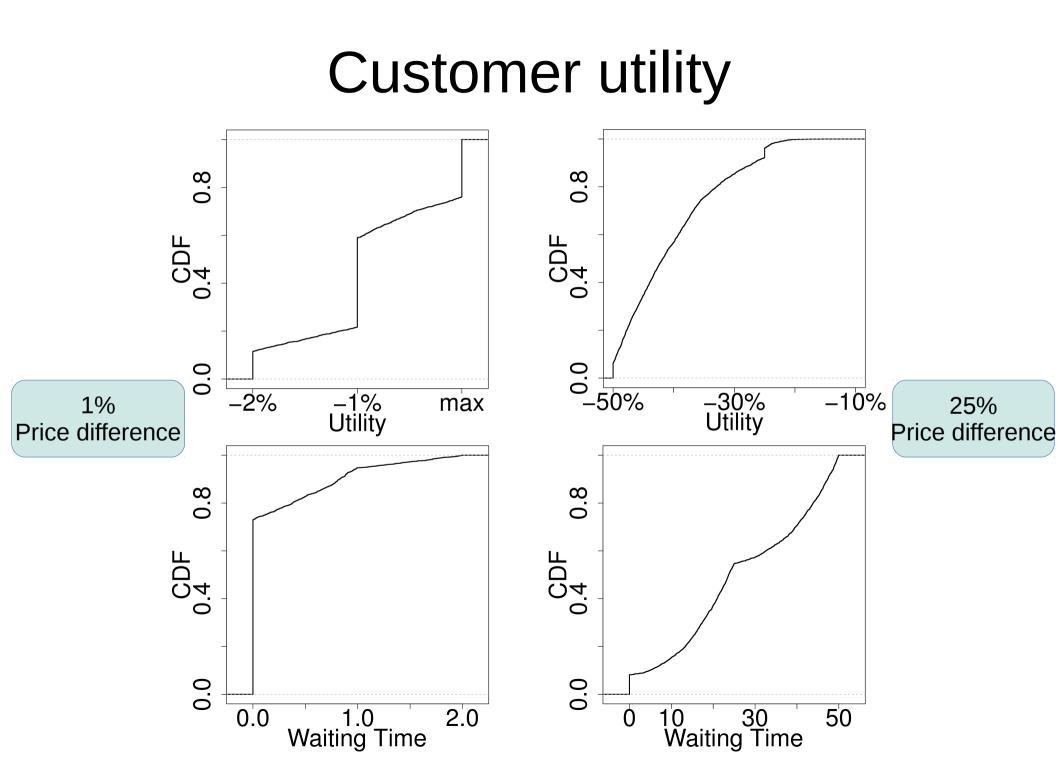
#### **Provider benefits**

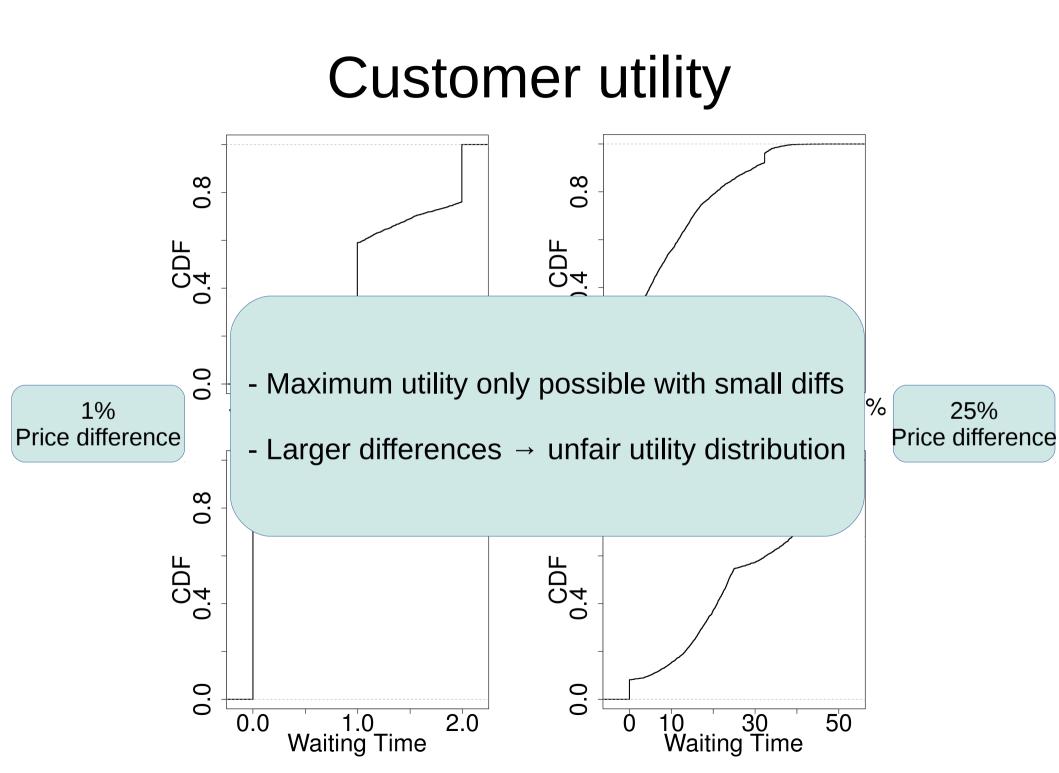


#### **Provider benefits**

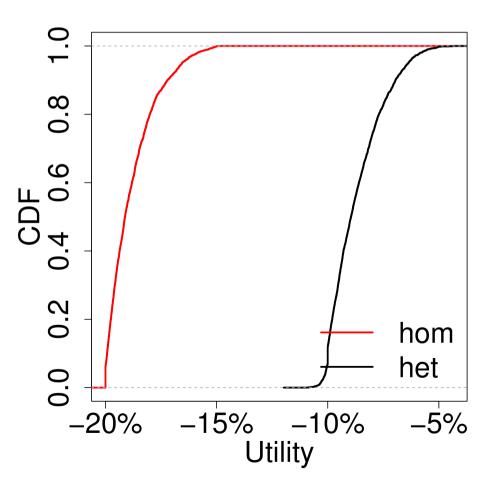




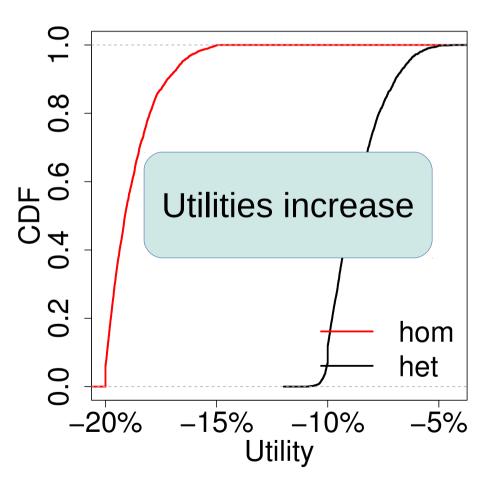




Utilities of (time) flexible user

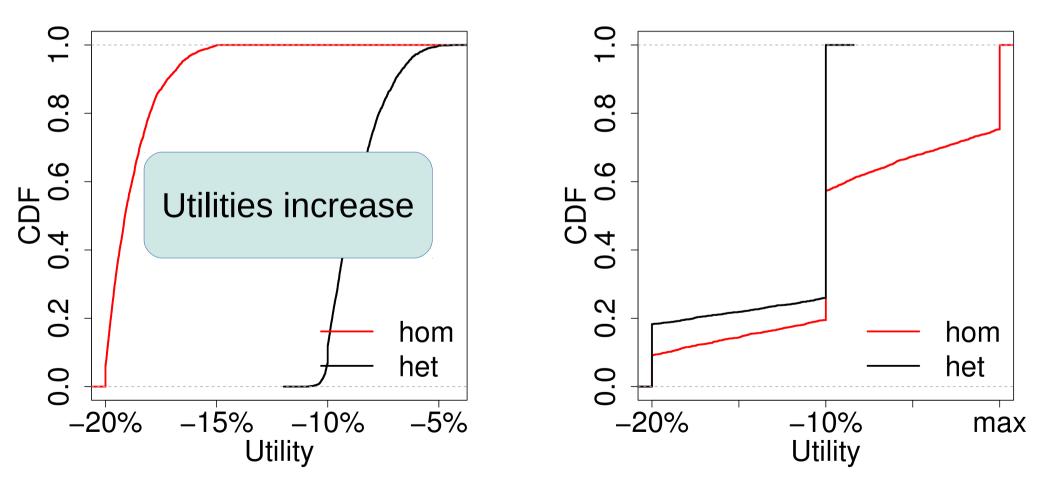


Utilities of (time) flexible user



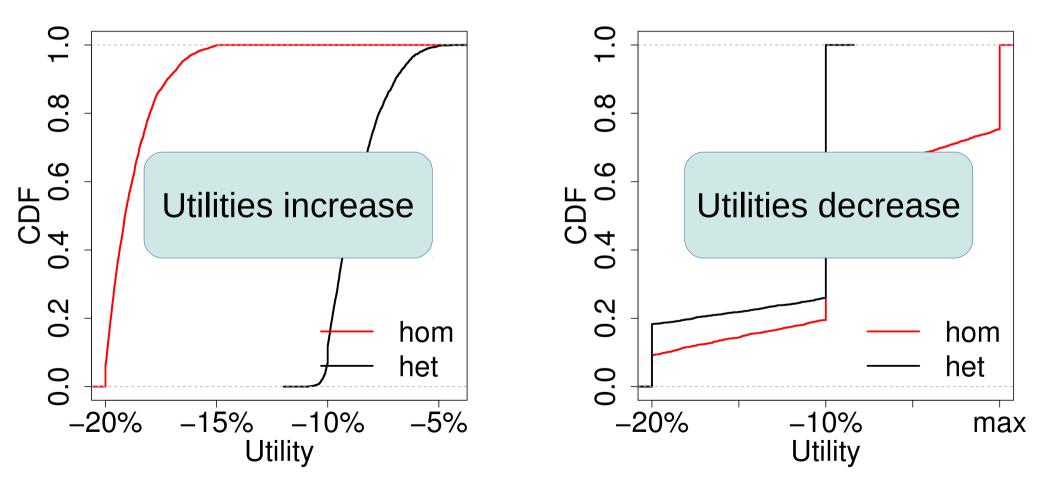
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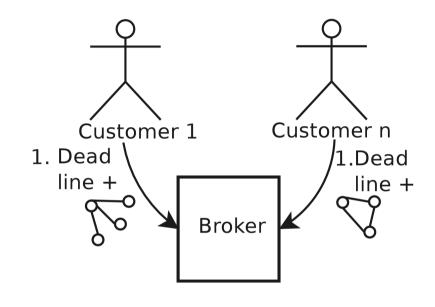
Utilities of (time) inflexible user

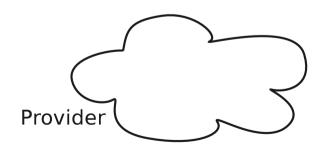


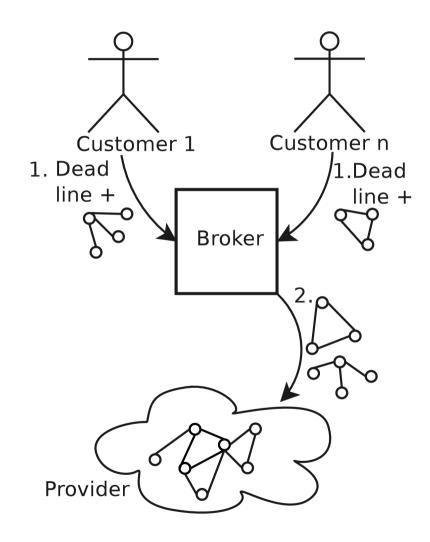
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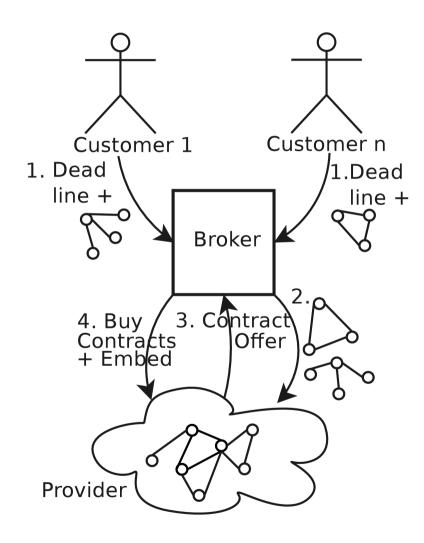
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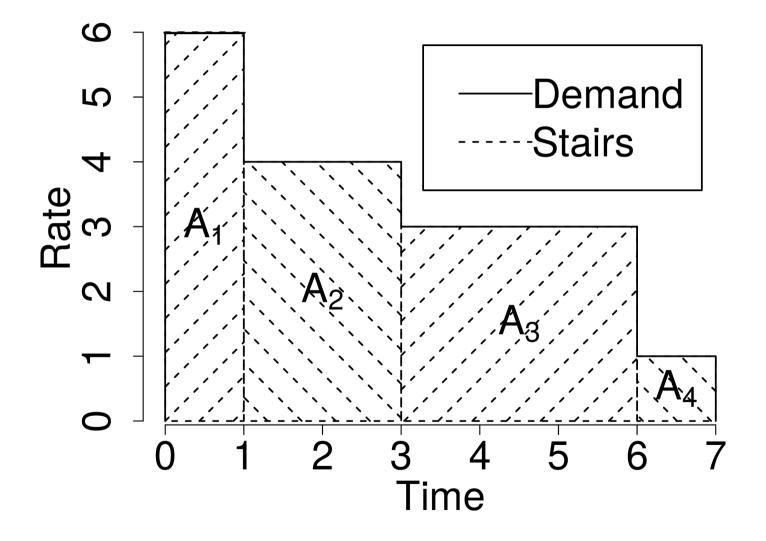




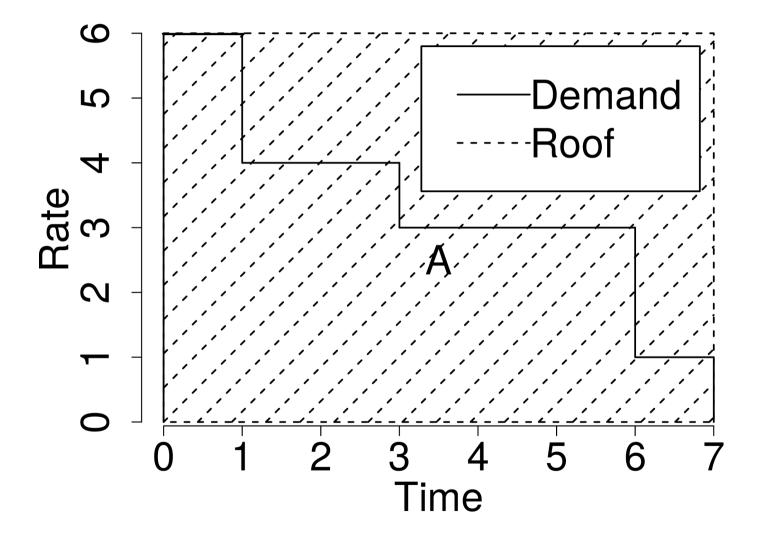


- More resources → smaller price per unit
- Broker can buy (rate, time)"rectangle" contracts
- VNet request have fixed deadline
- Customers are flexible/inflexible regarding the deadline

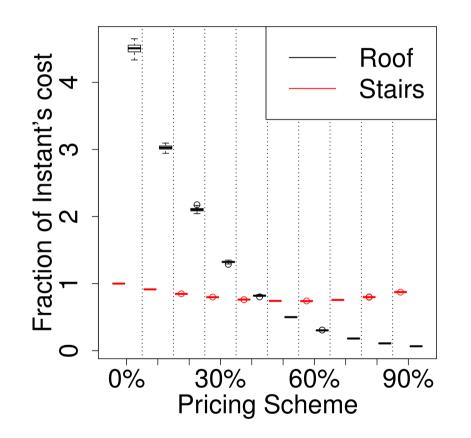
#### Broker strategy: stairs



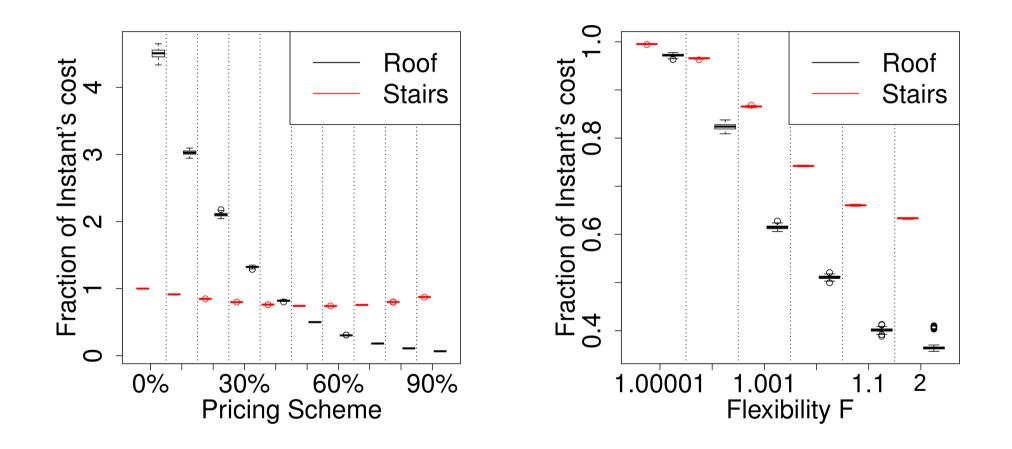
#### Broker strategy: roof



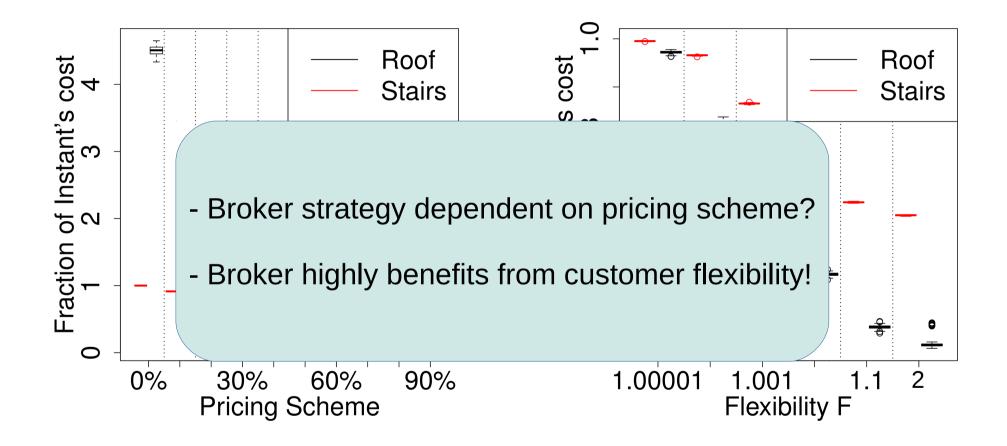
### Results



#### Results



### Results



#### **Outlook: resource allocation**

Is there a better way to buy discounted resources?

## **Related work**

- Ski rental (classical problem):
  - Buy or rent skis?

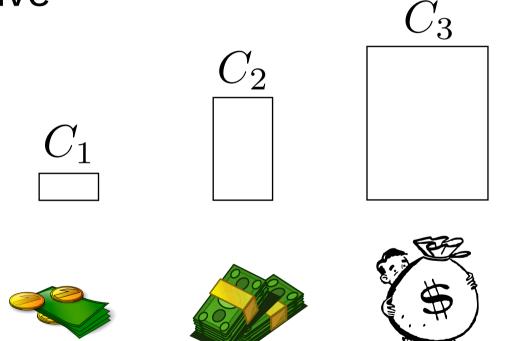


- Unpredictable car usage.
- Discount on longer permits.
- Which duration to buy?



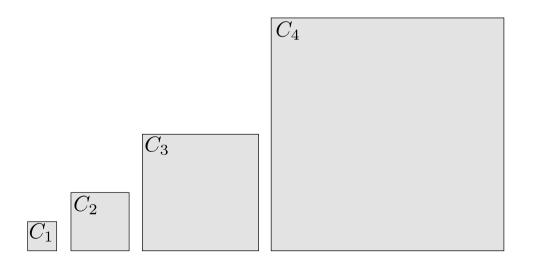
### Model - assumptions

1. *Monotonic Prices*, i.e. larger contracts – more expensive



### Model - assumptions

- 1. Monotonic Prices
- 2. *Multiplicity*, i.e., the next larger contract has a multiple of duration and rate



## Model

- Single resource (generalizable to multiple)
- Complete demand needs to be covered
- One lookahead model

- Objective: minimize overall price
  - Minimize competitive ratio

#### Competitive Strategy for Resource Allocation with Discounts?

## Online algorithm - ON2D

- ON2D mimics the offline algorithm OFF2D:
  - Compare previously bought contracts with the contracts OFF2D buys
  - If current demand is not covered, buy the new contracts

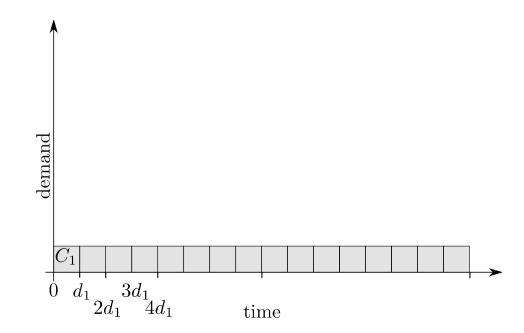
## Upper bound

- 1. Contract independence
- 2. Worst case classification
- 3. Maximum cumulative price

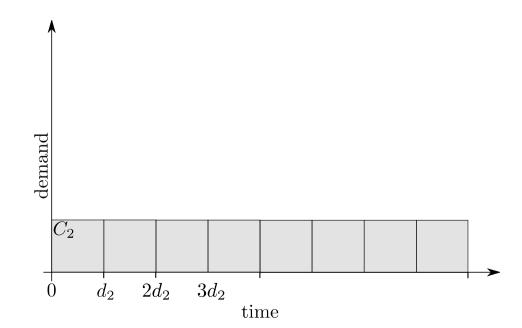
 $\rightarrow$  *k*-competitive (*k* = number of contracts)

- $p_1$  covered by  $\mathcal{C}_i$  , at time t
- $p_2$  covered by  $\mathcal{C}_j$  , at time t
  - → No contract C at time t' < t such that, C covers both  $p_1$  and  $p_2$

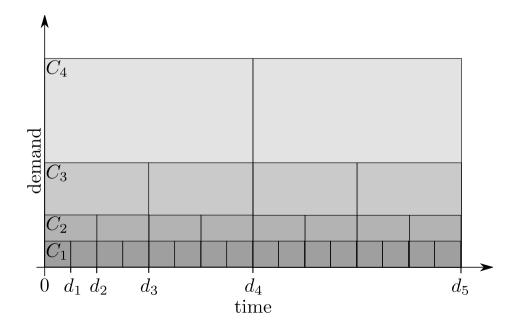
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# Upper bound worst case

• Worst case: OFF2D buys only one contract

- Proof by contradiction:
  - OFF2D buys  $\mathcal{C}_i$  and  $\mathcal{C}_j$ 
    - $\rightarrow$  ON2D buys a couple of contracts to cover both
  - Due to contract independence, the competitive ratio is the max of the competive ratios for both single scenarios

# Upper bound maximum cumulative price

- $\sum price_{ON}(\mathcal{C}_i) \leq i \cdot price(\mathcal{C}_i)$
- Inductive proof in the paper

 $\rightarrow$  ON2D is *k*-competitive, where *k* is the total number of contracts

## Online algorithm - ON2D

- ON2D mimics the offline algorithm OFF2D:
  - Compare previously bought contracts with the contracts OFF2D buys
  - If current demand is not covered, buy the new contracts
- Deterministic online algorithm O(k)
- Almost optimal (lower bound k/3)
- Polynomial offline algorithm

## Conclusion

- Shed light on specification benefits in
  - Horizontal markets
    - Lower priced providers profit the most
    - Flexible users benefit from inflexibles (not vice versa)
  - Vertical markets
    - Broker benefits from user flexibility
- Outlook: How to handle resource allocation (to appear @ICDCS)