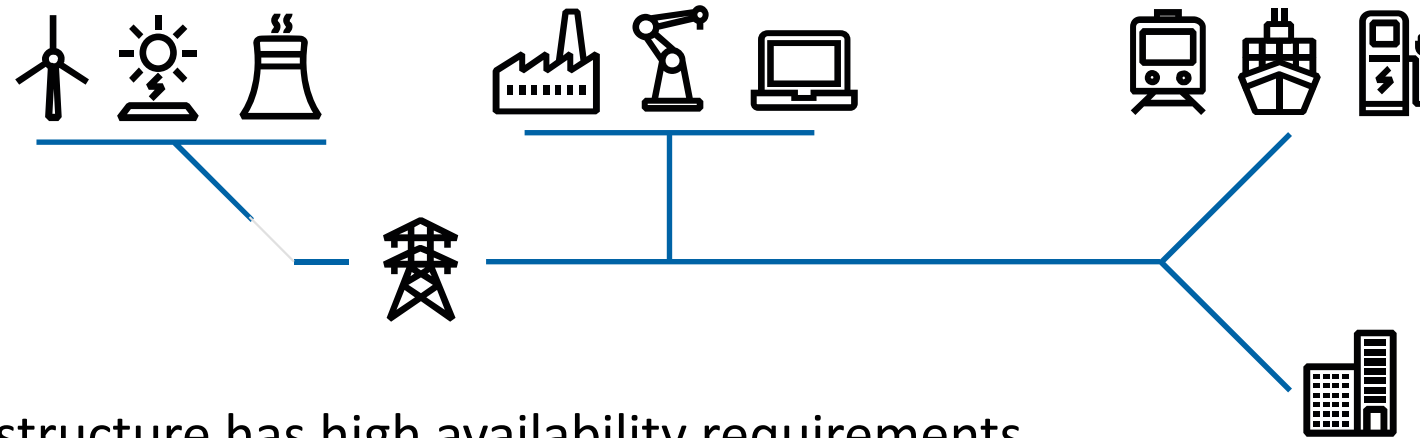


BA: What Can(not) Be Perfectly Rerouted Locally



- Critical network infrastructure has high availability requirements
- Hard real-time requirements on packet routing

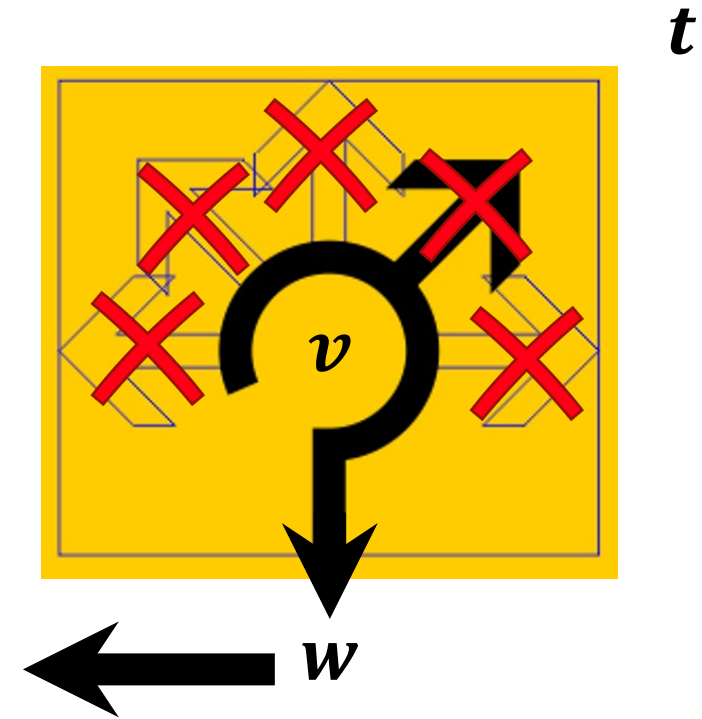
⇒ How to provide dependability guarantee despite edge failures in networks?

⇒ Possible without communication between nodes?

Prior work: Impossible in general [PODC 2012, Feigenbaum et al.]

Model

- Network is a connected undirected graph $G = (V, E)$
- Forwarding at node v may only match on:
 1. Destination t
 2. Incident edge failures $F \cap E(v)$
 3. Incoming port from $E(v) \cup \{\perp\}$
- No packet (header) changes allowed, no communication
- Static routing tables, deterministic behavior
- Goal: Install rules *ahead of time* to reach destination t under any edge failures F (if connected)



Perfect Resiliency on Non-Planar Graphs? Impossible!

- Perfect resiliency on Graph $G \rightarrow$ **Any subgraph** G' of G also allows for perfect resiliency
 - Idea: Take routing on G , fail edges to create G' , routing must still work 😊

- **Contracting** works as well, by a simulation argument



- Essentially: Take perfectly resilient routing function, show that we can make it work after contraction
- Combined: Perfect resiliency on Graph $G \rightarrow$ **Any minor** G' of G as well
 - Definition: G' is minor of G if can be obtained by contracting/subgraphing
- We show $K_5, K_{3,3}$ no perfect resilience \rightarrow **non-planar graphs do not allow for perfect resilience**

BA: What Can(not) Be Perfectly Rerouted Locally

- Perfect resiliency **impossible**:
 - On **some planar** graphs [already with just 7 nodes] and on **all non-planar** graphs
- Perfect resiliency **possible**:
 - On **some planar** graphs and on **all outerplanar** graphs
- In the extended version (link in paper):
 - Results on **more powerful routing** models, on **rule space size**, further open questions

*Full version also appears at SIAM Symposium on
Algorithmic Principles of Computer Systems
[APOCS'21]*