Software Transactional Networking: Concurrent and Consistent Policy Composition

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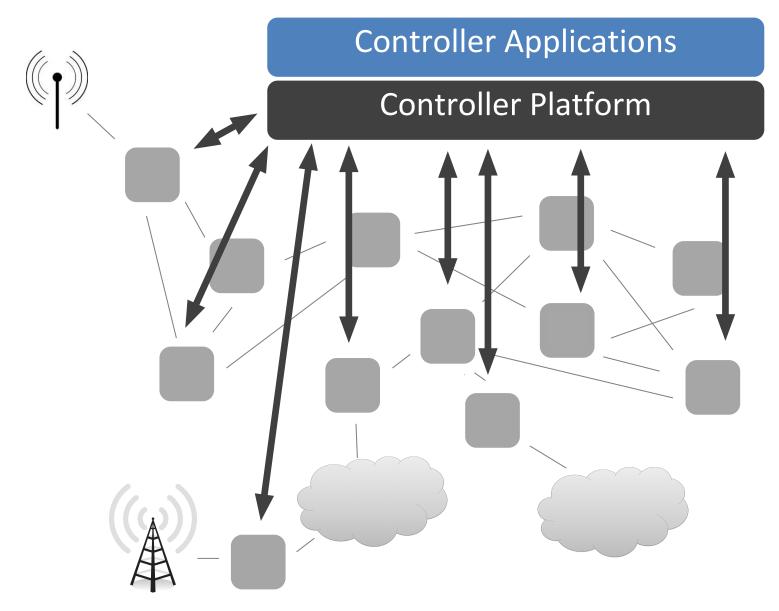
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Software Transactional Networking:

Concurrent and Consistent Policy Composition

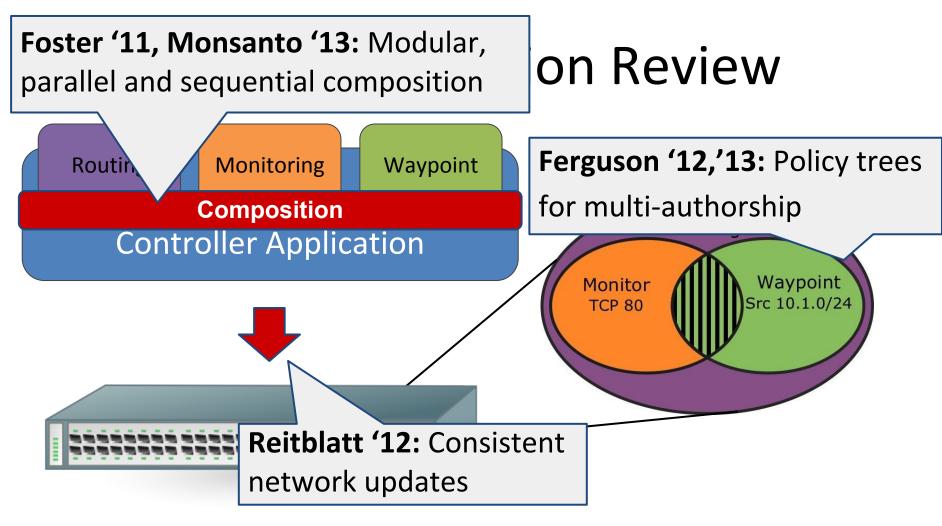
Network Policy Specification



Network-wide Policy: Not Monolithic Policy may originate from **multiple authors**, defined across **multiple functional modules**.

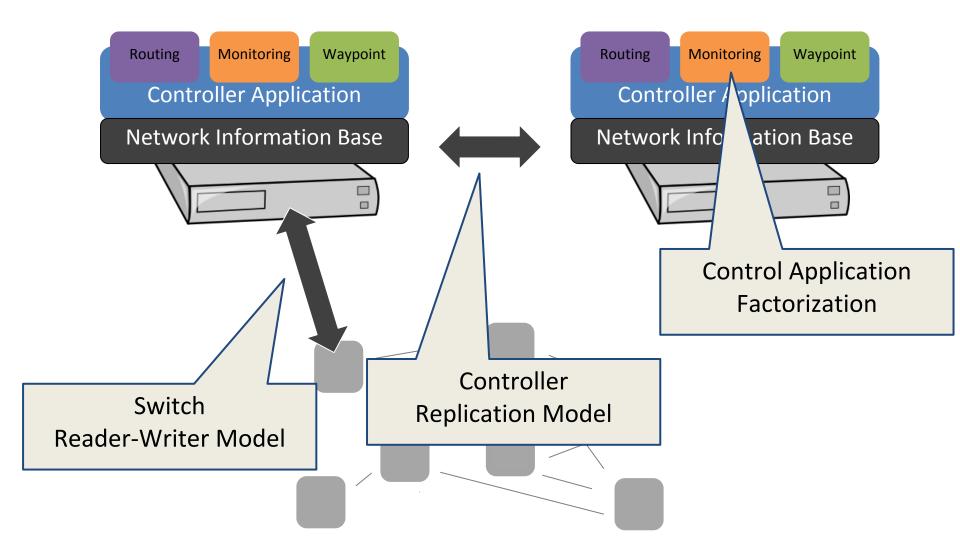


...necessitates policy composition prior to network update.



- 1. Precedence must be defined across policy sources
- 2. Packet forwarding rule priorities must be defined, and respect policy source precedence

Now, consider policy composition in the distributed control plane...



Research Question #1

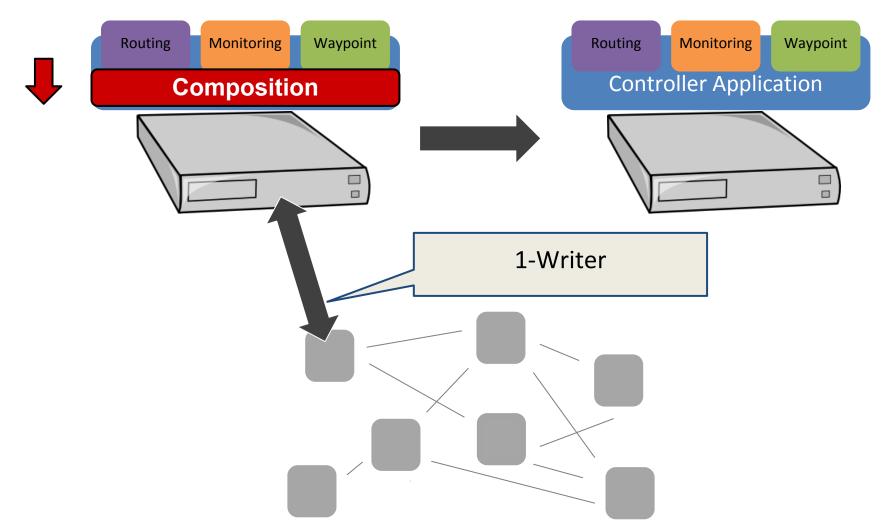
How does the design of the **distributed control plane** affect policy composition with respect to:

- 1. The consistency of the composition
- 2. The semantics of consistent network update

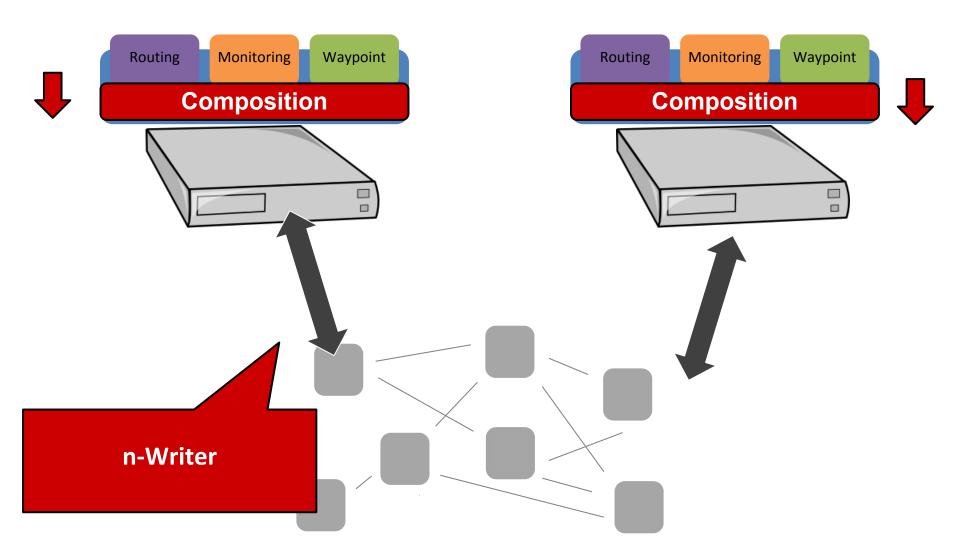
Distributed Control Models 1: Hot Standby Replication

Master

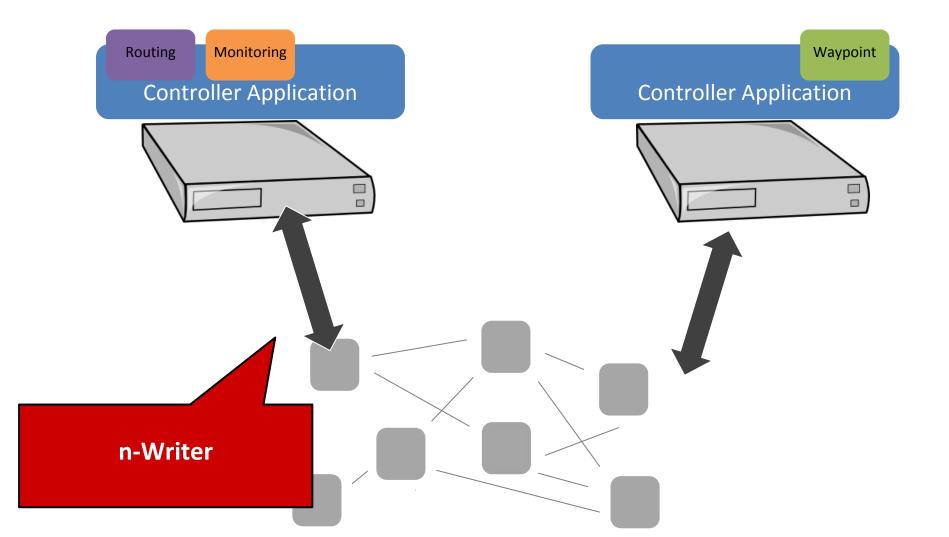
Slave

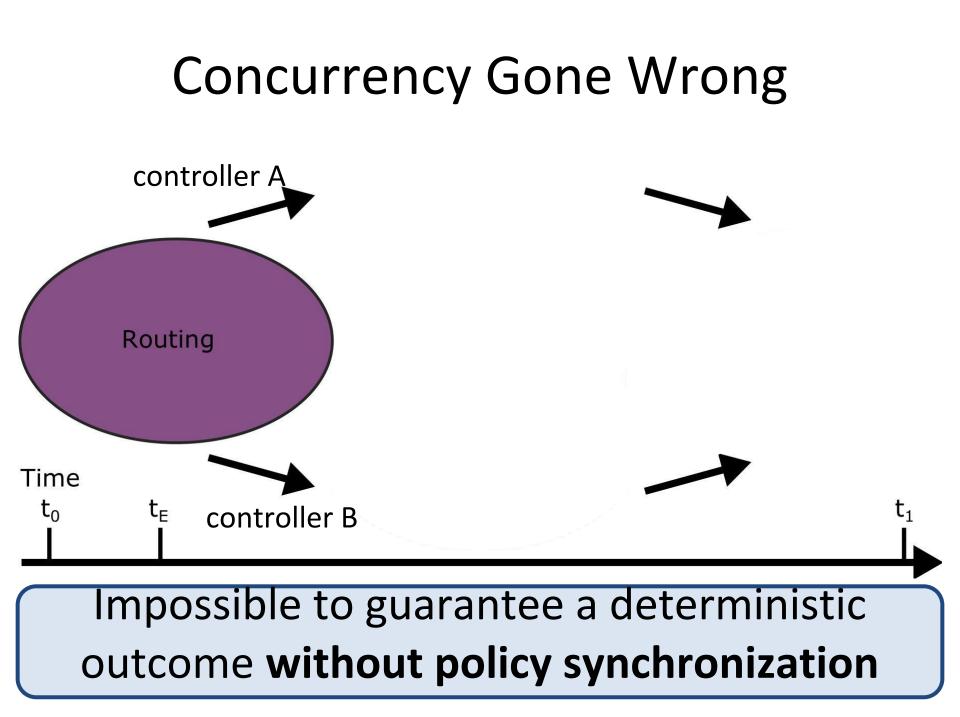


Distributed Control Models 2: Sharding by Disjoint Flow-Space

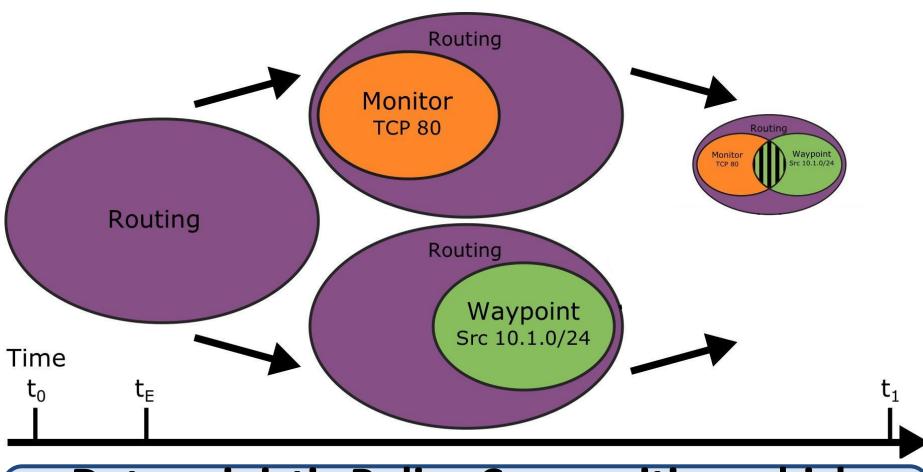


Distributed Control Models 3: Sharding by Policy





Consistent Policy Composition



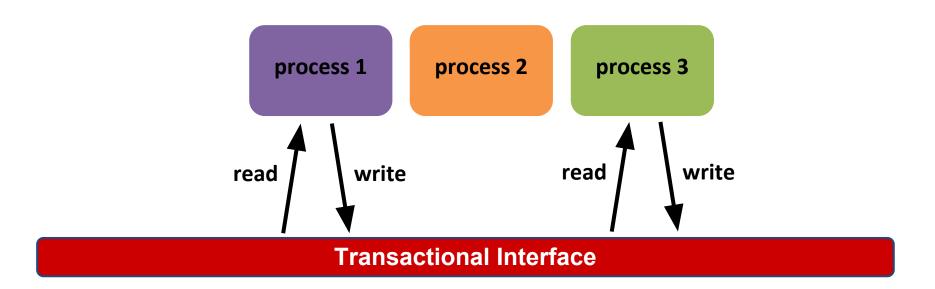
Deterministic Policy Composition which respects precedence of multi-authorship

Research Question #2

Can we realize a distributed policy composition interface that...

is agnostic to:	guarantees:		
Control Distribution model	Consitent update semantics		
Switch reader-writer model	Consitent policy composition		

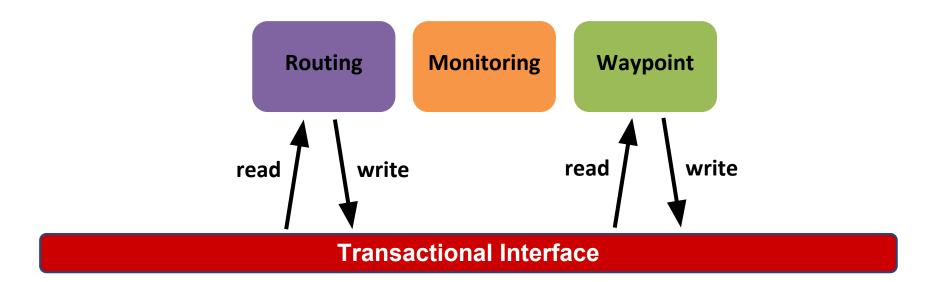
Software Transactional Memory

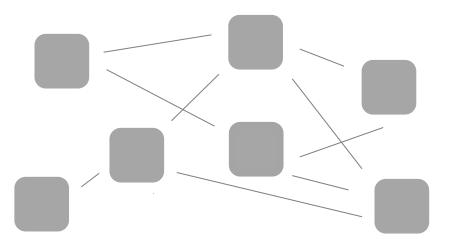


Shared Datastructure

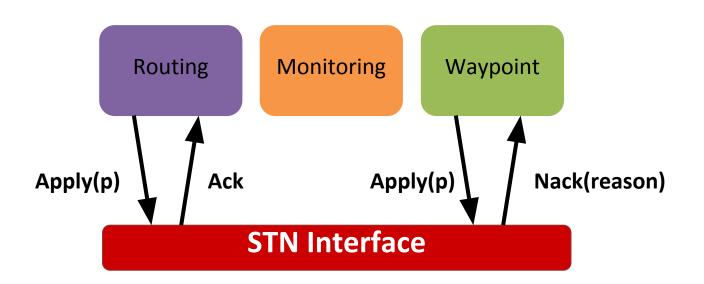


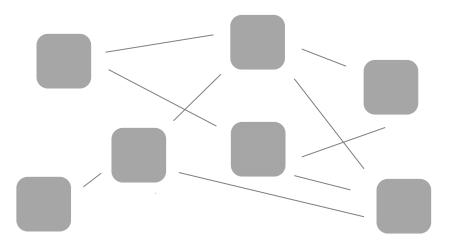
Software Transactional Networking



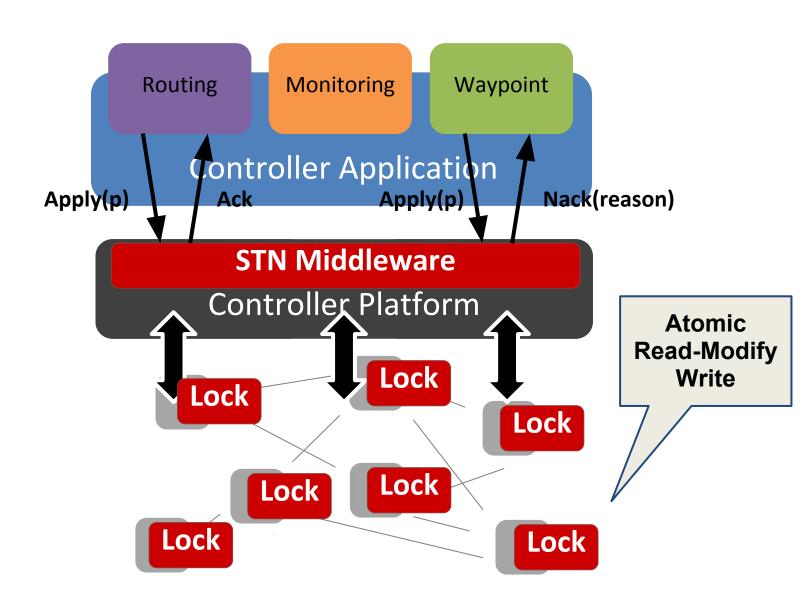


The STN Interface

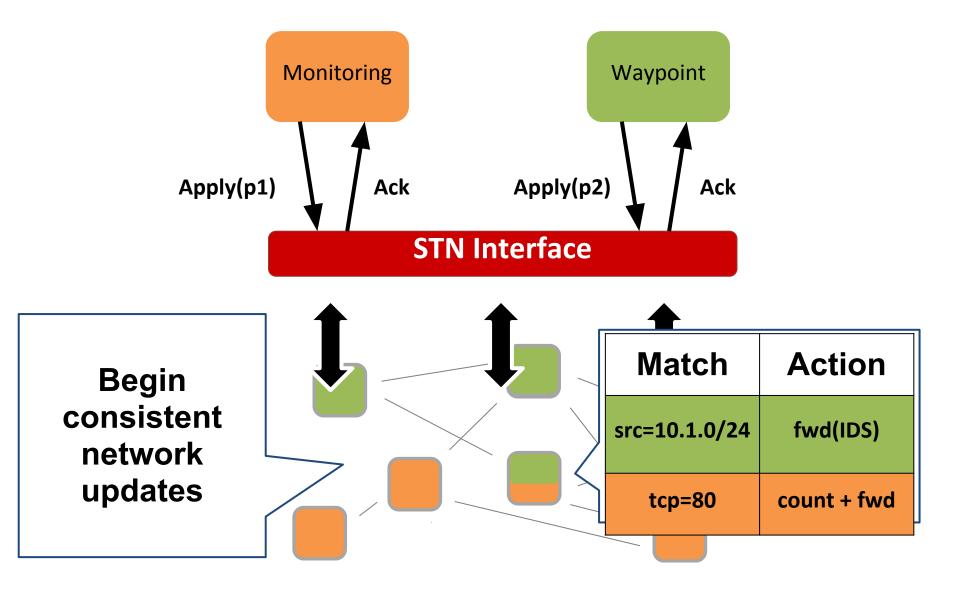




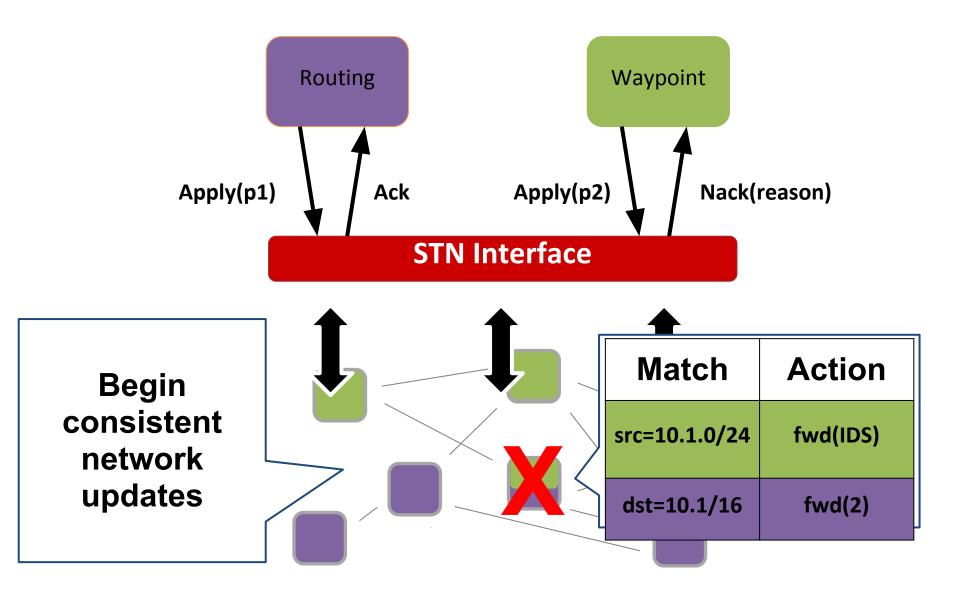
Conceptualizing STN



STN in Action (Ack Case)



STN in Action (Nack Case)



STN fine-grained locking algorithm

- 1 apply(*policy* p, *policy tree* f):
- 2 tag = newTag(p)

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- 3 for s in switches():
- 4 if s.rules.doCompose(p, f):
 - tag' = concurrentTag(s, tag)
 - s.addTagRules(p, tag')
- 7 else *remove* tag' *rules, nack(reason)*
- 8 for s in ingressSwitches():
- 9 s.addRule(match=*, action=push(tag'))

Open Issues

- For our simple algorithm, tag and forwardingrule state grows exponentially O(2ⁿ) for nconcurrency
- Different Update Consistency and Policy Composition Semantics
- Atomic read-modify-write primitive at switch
- Controller fault tolerance

Summary

Full technical report for more details http://arxiv.org/abs/1305.7429

Internship Opportunities @T-Labs

- Interface for distributed policy composition
- Framework to reason about concurrent policy composition and consistent update

Backup Slides

Policy Composition Review

Priority	Match	Action
1	dst=10.1/16, *	fwd(1)
1	dst=10.2/16, *	fwd(2)
2	dst=10.1/16,dport=80, *	fwd(1)
2	dst=10.2/16,dport=80, *	fwd(2)
3	dst=10.1/16, src=10.1.0/24, dport=80, *	fwd(IDS)
•••		

Routing:	dst=10.1/16 \rightarrow fwd(1) dst=10.2/16 \rightarrow fwd(2)	Composition is the "cross-product" of
Monitor:	tcp_port=80 \rightarrow count	rules
Waypoint:	src=10.1.0/24 \rightarrow fwd(IDS)	

Distributed Control Models 2: Sharding by Region

