Dependable Networks: Trends and Challenges Stefan Schmid (TU Berlin & Sabbatical at Hebrew University)

INET @ TU Berlin

We aim at the investigation of future communication **networks** and future applications offered through these networks:

- Algorithms and mechanisms to design and operate communication networks
- Network architectures and protocols for future communication technologies
- Performance evaluation of networked and distributed systems
- Network security
- Wireless and cellular networks

Our vision is that networked systems should become *self-** (i.e., self-optimizing, self-repairing, self-configuring).

Accordingly, we are currently particularly interested in *automated* and *data-driven* approaches to design, optimize, and verify networked systems.



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But why?? Networks are working well today! Internet is huge success, handled all trends!



A Major Issue: Complexity

Many outages due to misconfigurations and human errors.

Entire countries disconnected...

Data Centre > Networks

Google routing blunder sent Japan's Internet dark on Friday

Another big BGP blunder

By Richard Chirgwin 27 Aug 2017 at 22:35

40 📮 SHARE 🔻

Last Friday, someone in Google fat-thumbed a border gateway protocol (BGP) advertisement and sent Japanese Internet traffic into a black hole.

The trouble began when The Chocolate Factory "leaked" a big route table to Verizon, the result of which was traffic from Japanese giants like NTT and KDDI was sent to Google on the expectation it would be treated as transit.

... 1000s passengers stranded...

British Airways' latest Total Inability To Support Upwardness of Planes* caused by Amadeus system outage

Stuck on the ground awaiting a load sheet? Here's why

By Gareth Corfield 19 Jul 2018 at 11:16 109 🖵 SHARE 🔻



RA flighte amound the world wave arounded as a result of the Amadeus outage

... even 911 services affected!

Officials: Human error to blame in Minn. 911 outage

According to a press release, CenturyLink told department of public safety that human error by an employee of a third party vendor was to blame for the outage

Aug 16, 2018

Duluth News Tribune

SAINT PAUL, Minn. — The Minnesota Department of Public Safety Emergency Communication Networks division was told by its 911 provider that an Aug. 1 outage was caused by human error.

Even Tech-Savvy Companies Struggle to Provide Reliable Networks



We discovered a misconfiguration on this pair of switches that caused what's called a "bridge loop" in the network.

> A network change was [...] executed incorrectly [...] more "stuck" volumes and added more requests to the remirroring storm





Service outage was due to a series of internal network events that corrupted router data tables

> Experienced a network connectivity issue [...] interrupted the airline's flight departures, airport processing and reservations systems



Another Issue: *Lack of Tools* Anecdote "Wall Street Bank"

- Outage of a data center of a Wall Street investment bank
- Lost revenue measured in USD 10⁶ / min
- Quickly, an emergency team was assembled with experts in compute, storage and networking:
 - **The compute team:** soon came armed with reams of logs, showing how and when the applications failed, and had already written experiments to reproduce and isolate the error, along with candidate prototype programs to workaround the failure.
 - **The storage team:** similarly equipped, showing which file system logs were affected, and already progressing with workaround programs.
 - "All the networking team had were two tools invented over 20y ago to merely test end-to-end connectivity. Neither tool could reveal problems with switches, the congestion experienced by individual packets, or provide any means to create experiments to identify, quarantine and resolve the problem. Whether or not the problem was in the network, the networking team would be blamed since they were unable to demonstrate otherwise."

Source: «The world's fastest and most programmable networks» White Paper Barefoot Networks

Roadmap

- Opportunity: emerging networking technologies
 - Automation and "self-driving networks"
 - Programmable networks for improved visibility
- Challenge: emerging network technologies
 - New threat models

It's an *exciting period*! New tools, simple abstractions, disburdening human operators, etc.



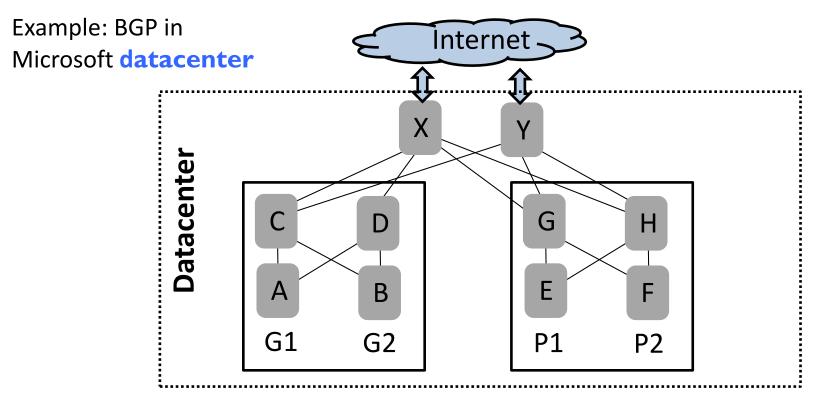
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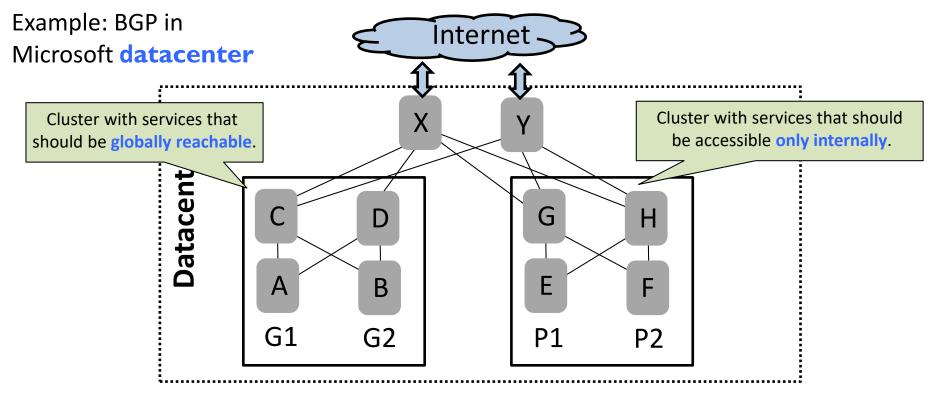
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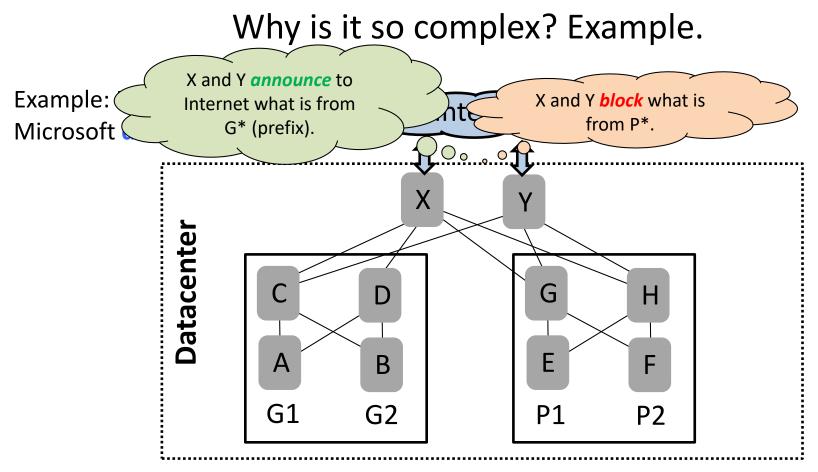


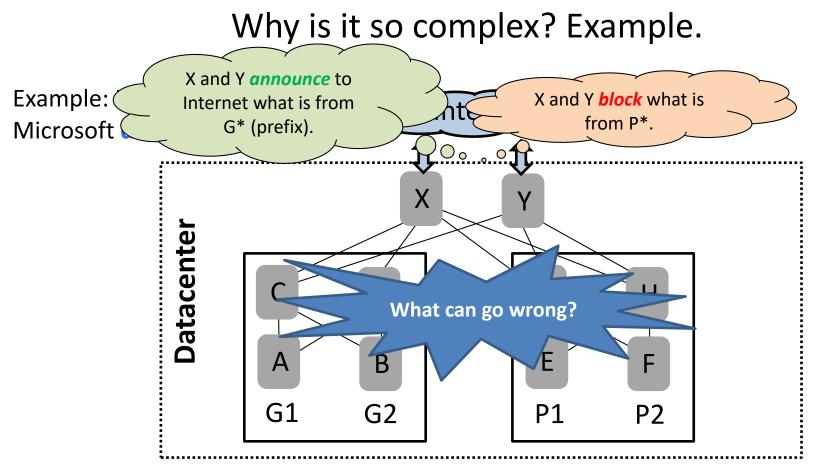
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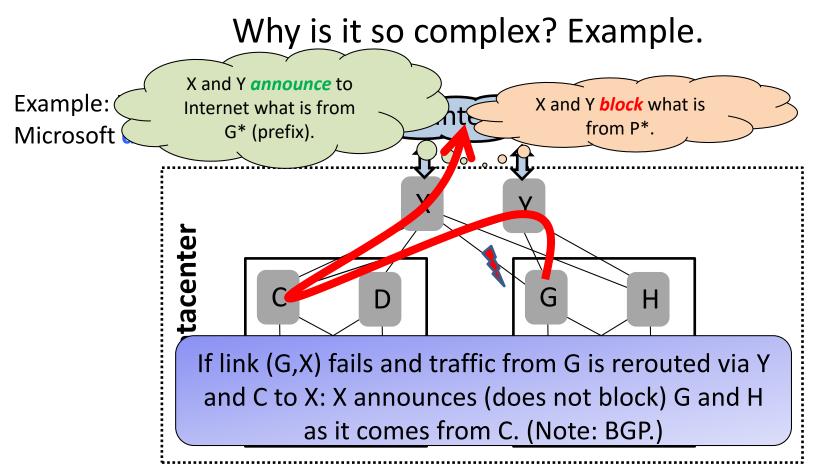


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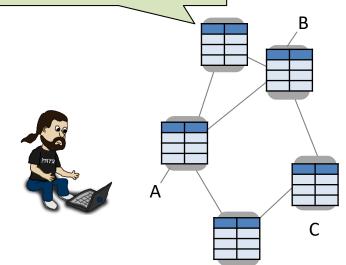


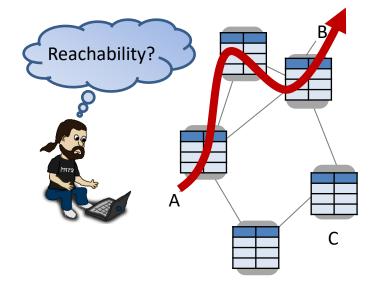






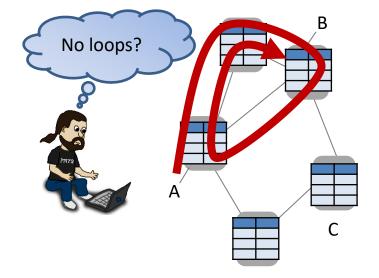
Routers and switches store list of forwarding rules, and conditional failover rules.





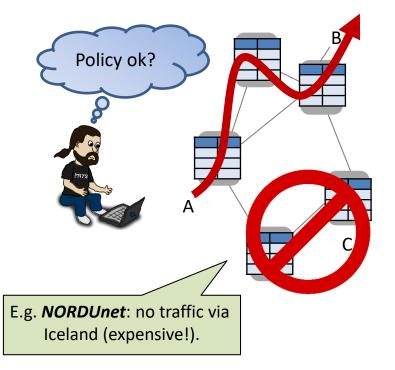
Sysadmin responsible for:

• **Reachability:** Can traffic from ingress port A reach egress port B?



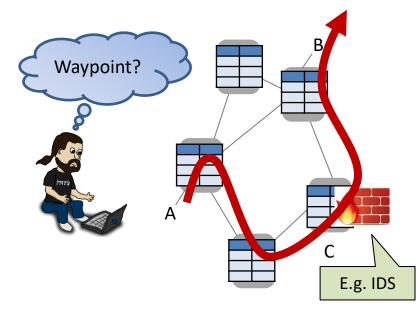
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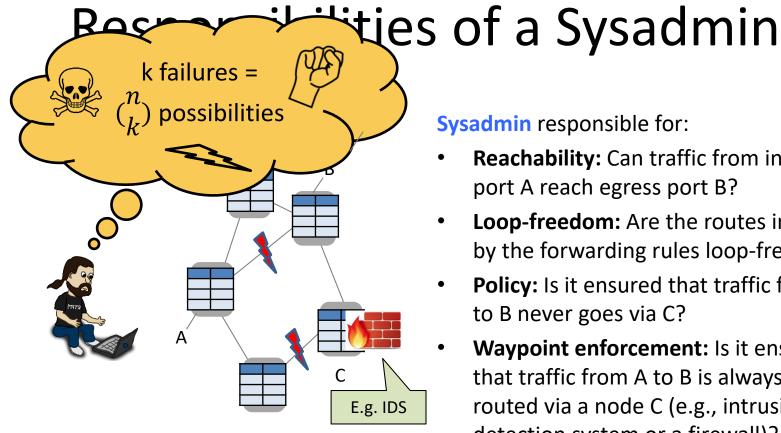
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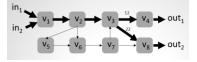
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... and everything even under multiple failures?!

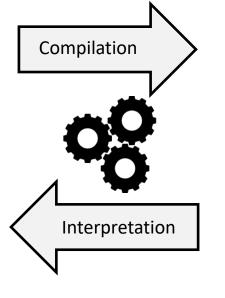
Vision: Automation and Formal Methods



FL	In-I	In-Label	Out-I	op
τ_{v_1}	in ₁	1	(v_1, v_2)	push(10)
10	in_2	1	(v_1, v_2)	push(20)
T_{W_2}	(v_1, v_2)	10	(v_2, v_3)	swap(11)
~~~	$(v_1, v_2)$	20	$(v_2, v_3)$	swap(21)
$\tau_{vs}$	$(v_2, v_3)$	11	$(v_3, v_4)$	swap(12)
	$(v_2, v_3)$	21	$(v_3, v_8)$	swap(22)
	$(v_7, v_3)$	11	$(v_3, v_4)$	swap(12)
	$(v_7, v_3)$	21	$(v_3, v_8)$	swap(22)
$\tau_{v_4}$	$(v_3, v_4)$	12	out1	pop
$\tau_{v_{5}}$	$(v_2, v_5)$	40	$(v_5, v_6)$	pop
$\tau_{v_6}$	$(v_2, v_6)$	30	$(v_6, v_7)$	swap(31)
	$(v_5, v_6)$	30	$(v_6, v_7)$	swap(31)
	$(v_5, v_6)$	61	$(v_6, v_7)$	swap(62)
	$(v_5, v_6)$	71	$(v_6, v_7)$	swap(72)
$\tau_{v_7}$	$(v_6, v_7)$	31	$(v_7, v_3)$	pop
	$(v_6, v_7)$	62	$(v_7, v_3)$	swap(11)
	$(v_6, v_7)$	72	$(v_7, v_8)$	swap(22)
$\tau_{v_{\theta}}$	$(v_3, v_8)$	22	out ₂	pop
	$(v_7, v_8)$	22	out ₂	pop

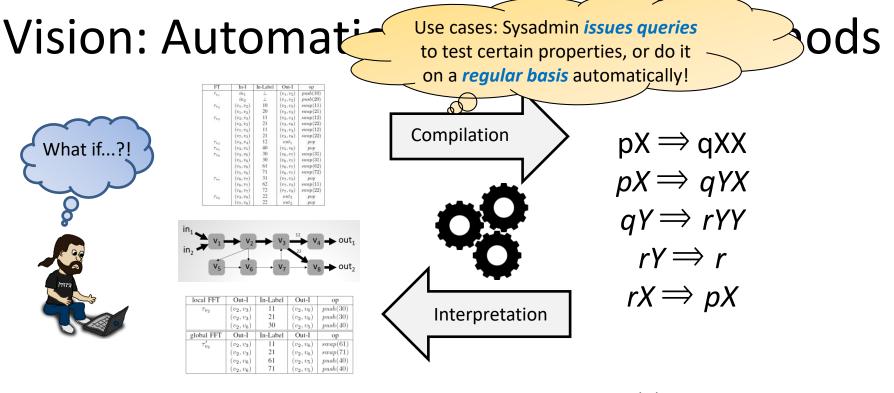


local FFT	Out-I	In-Label	Out-I	op
$\tau_{v_2}$	$(v_2, v_3)$	- 11	$(v_2, v_6)$	push(30)
	$(v_2, v_3)$	21	$(v_2, v_6)$	push(30)
	$(v_2, v_6)$	30	$(v_2, v_5)$	push(40)
global FFT	Out-I	In-Label	Out-I	op
$\tau'_{v_2}$	$(v_2, v_3)$	11	$(v_2, v_6)$	swap(61)
-	$(v_2, v_3)$	21	$(v_2, v_6)$	swap(71)
	$(v_2, v_6)$	61	$(v_2, v_5)$	push(40)
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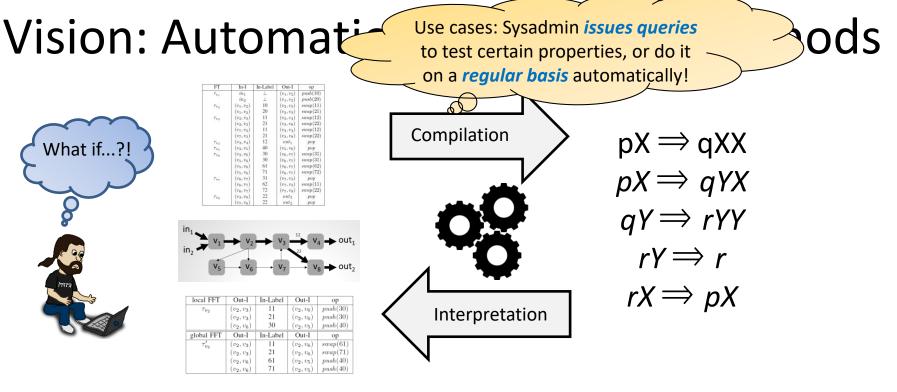


$$pX \Rightarrow qXX$$
$$pX \Rightarrow qYX$$
$$qY \Rightarrow rYY$$
$$rY \Rightarrow r$$
$$rX \Rightarrow pX$$

Router **configurations**, Segment Routing etc. Pushdown Automaton and Prefix Rewriting Systems Theory

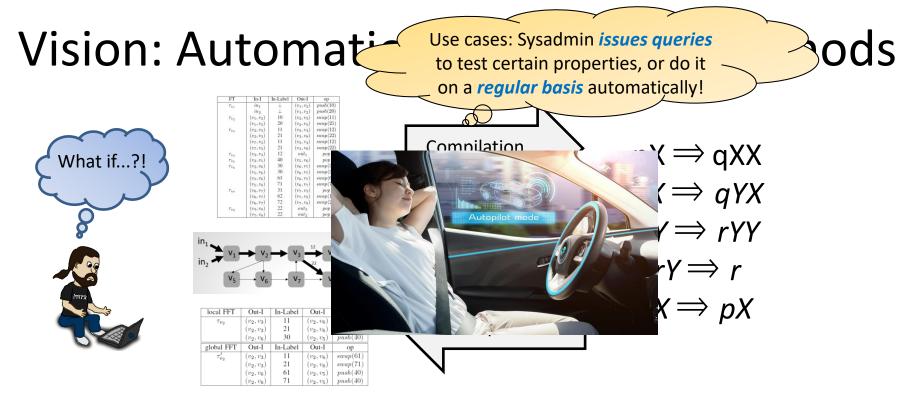


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P-Rex: Fast Verification of MPLS Networks with Multiple Link Failures. Jensen et al., ACM CoNEXT, 2018.



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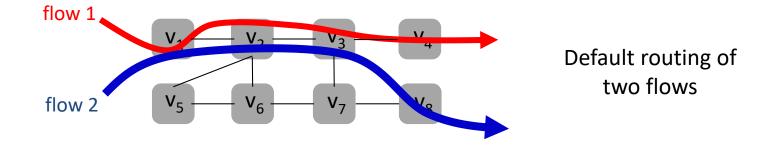
# Case Study: MPLS Networks

- Widely deployed networks by Internet Service Providers (ISPs)
- Often used for traffic engineering
  - Avoid congestion by going non-shortest paths
- Allows for *header re-writing* upon failures
  - Header based on stack of labels

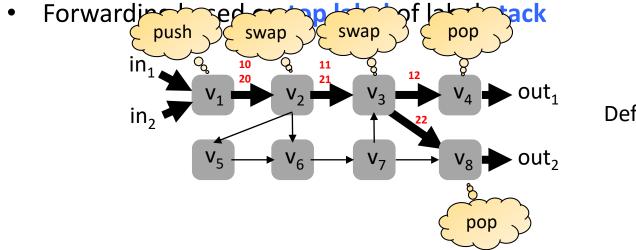


# How (MPLS) Networks Work

• Forwarding based on top label of label stack



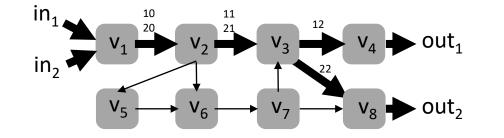
## How (MPLS) Networks Work



Default routing of two flows

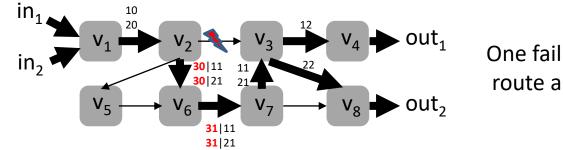
### Fast Reroute Around 1 Failure

• Forwarding based on top label of label stack (in packet header)



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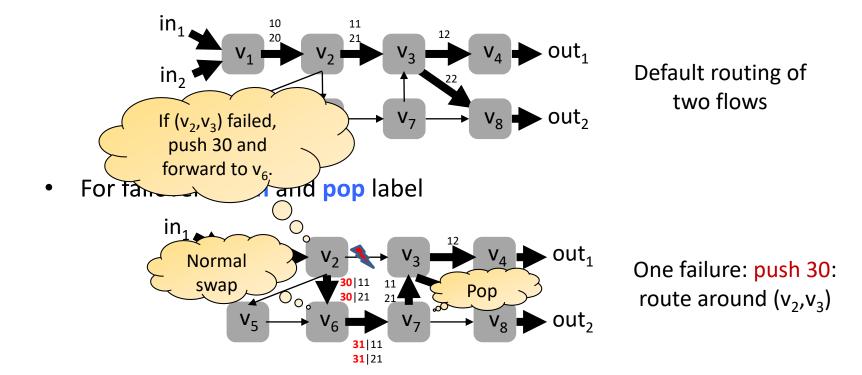
• For failover: push and pop label



One failure: push 30: route around  $(v_2, v_3)$ 

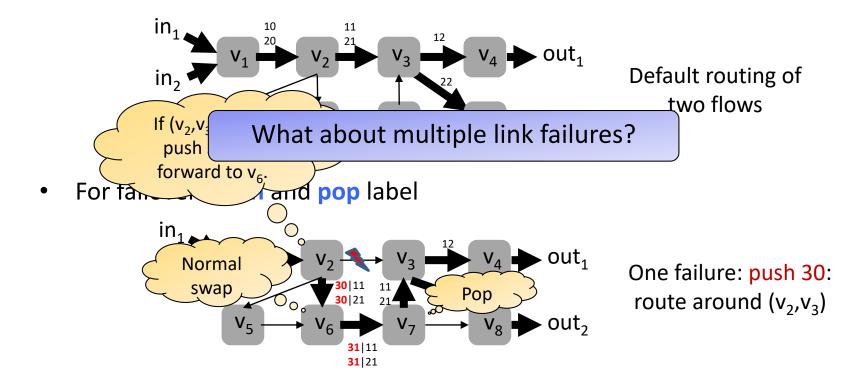
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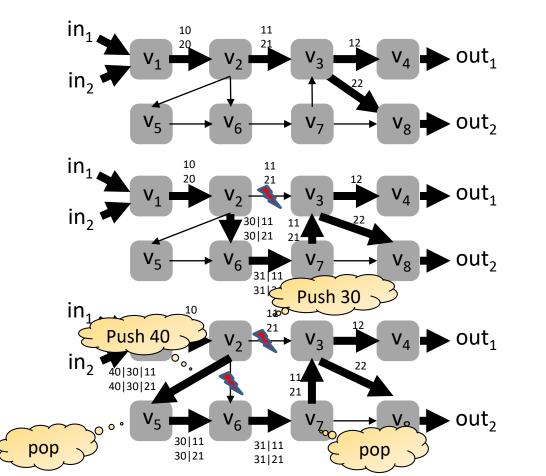


### Fast Reroute Around 1 Failure

• Forwarding based on top label of label stack (in packet header)



### 2 Failures: Push *Recursively*



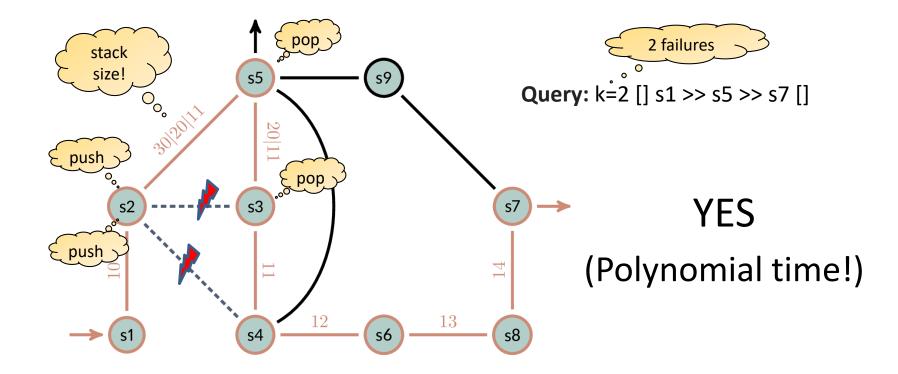
**Original** Routing

**One failure**: push 30: route around  $(v_2, v_3)$ 

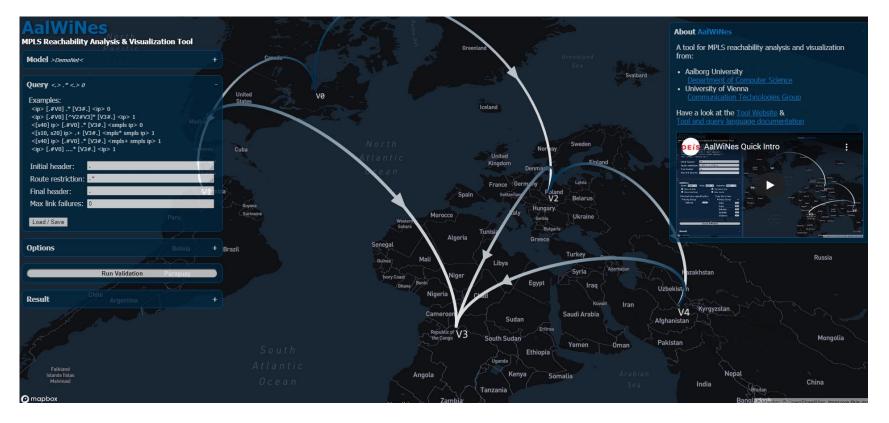
Two failures: first push 30: route around (v₂,v₃) *Push recursively* 40: route around (v₂,v₆)

### Example: P-Rex for MPLS Networks

Can traffic starting with [] go through s5, under up to k=2 failures?



### Demo of P-Rex / AalWiNes Tool



Tool: <u>https://demo.aalwines.cs.aau.dk/</u>, Youtube: <u>https://www.youtube.com/watch?v=mvXAn9i7_Q0</u>

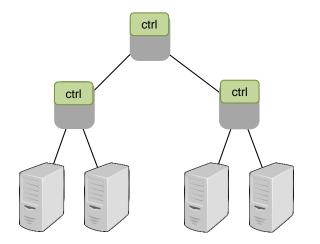
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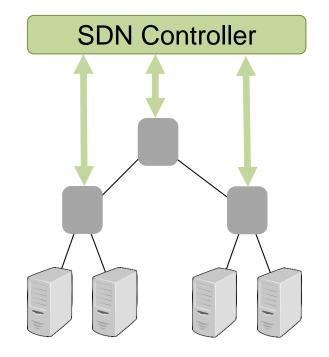


#### Software-Defined Networks



Traditionally:

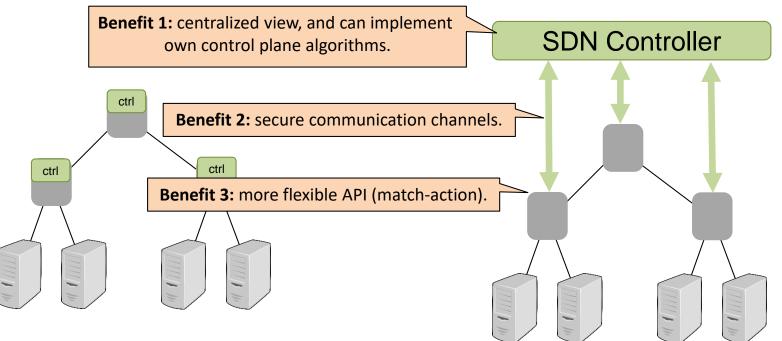
- Distributed control plane
- Blackbox, not programmable



Software-defined Networs (SDN):

- Logically centralized control
- Programmable, match-action

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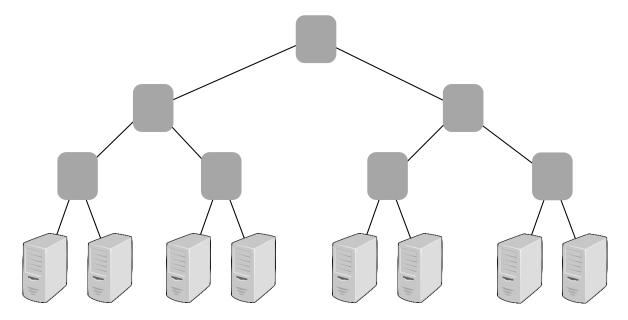
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# Example Application for SDN: Detecting Misbehavior

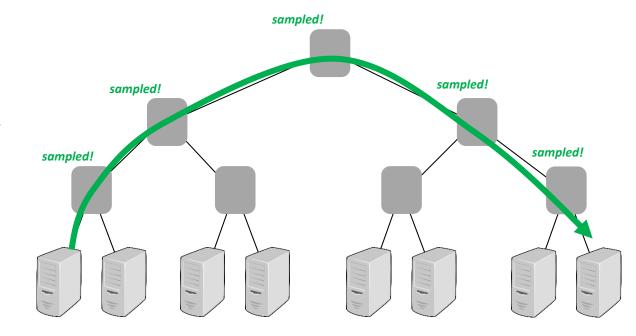
Monitor packets, traditionally: trajectory sampling

- Globally sample packets with hash(imm. header)∈[x,y]
- See full routes *of some packets*



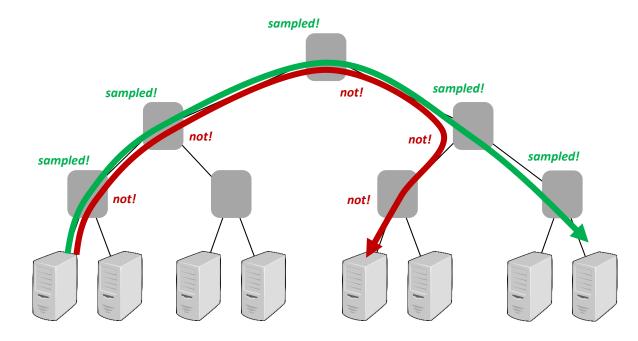
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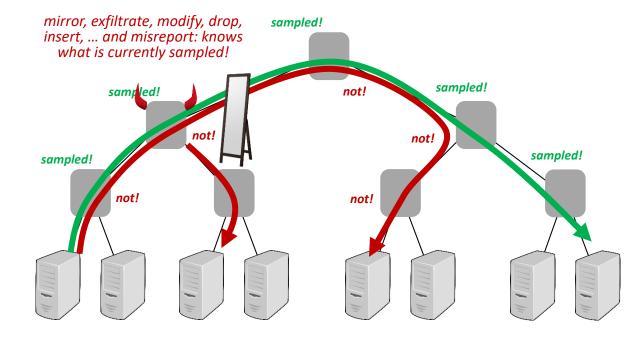
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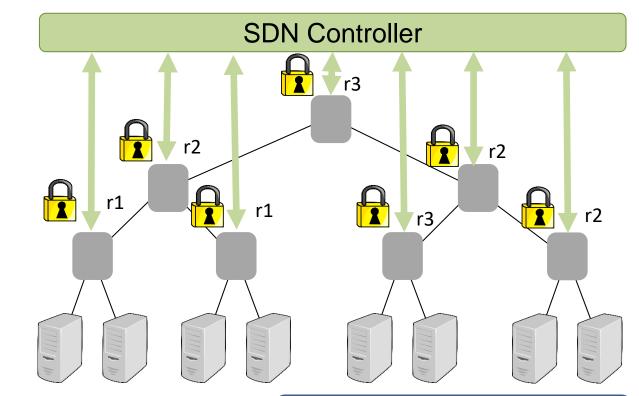
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### Solution: Use SDN for Secure Trajectory Sampling

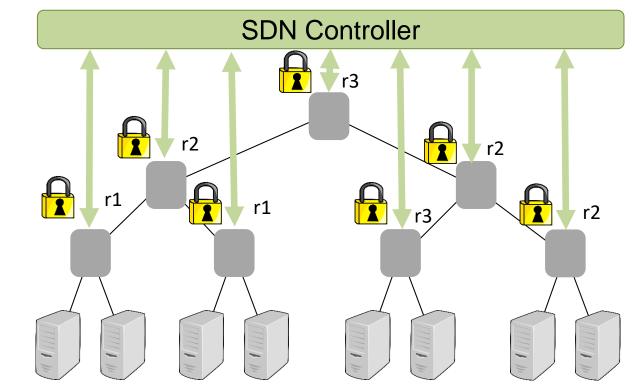
- Idea:
  - Use secure channels between controller and switches to distribute hash ranges
  - Give different hash ranges hash ranges to different switches, but add some redundancy: risk of being caught!



Network Policy Checker for Adversarial Environments. Kashyap Thimmaraju, Liron Schiff, and S. SRDS 2019.

### Solution: Use SDN for Secure Trajectory Sampling

- Idea:
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  - Give different hash ranges hash ranges to different switches, but add some redundancy: risk of being caught!
- In general: obtaining live data from the network *becomes easier!*



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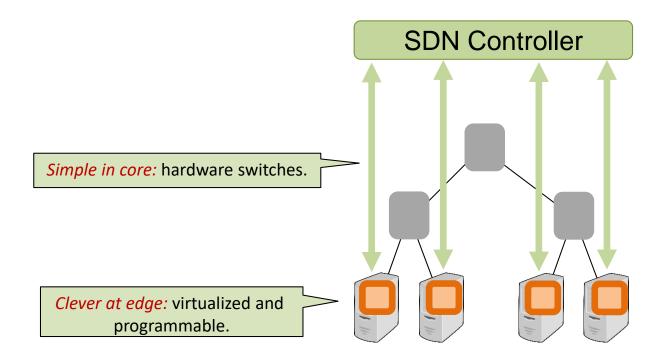
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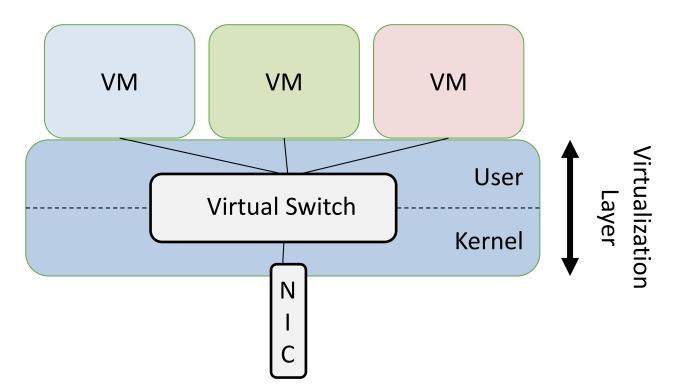
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#### SDN in Datacenters: Virtual Switches

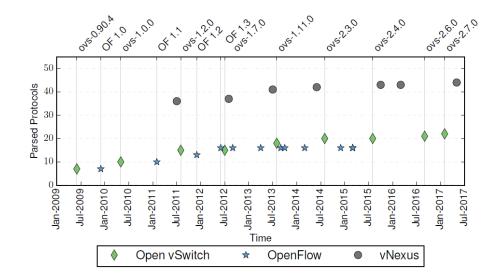


#### The Virtual Switch



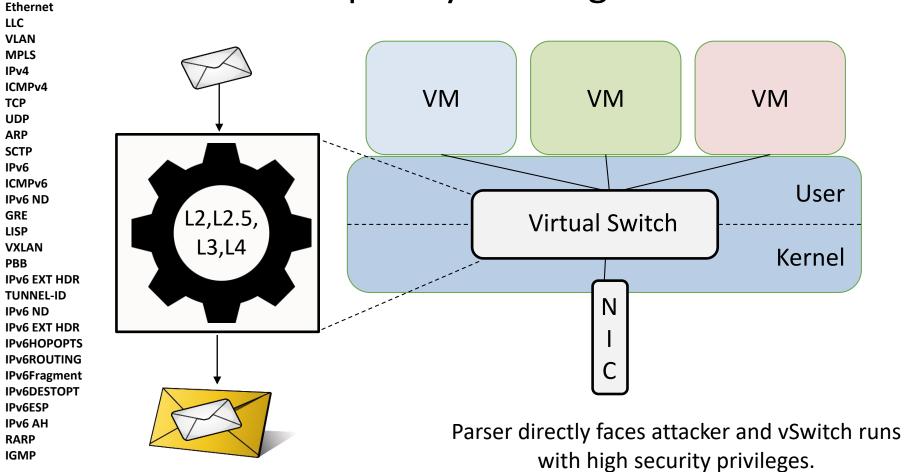
Virtual switches reside in the **server's virtualization layer** (e.g., Xen's Dom0). Goal: provide connectivity and isolation.

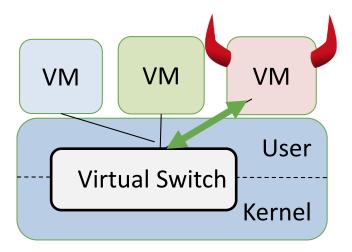
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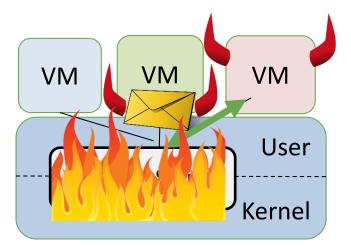


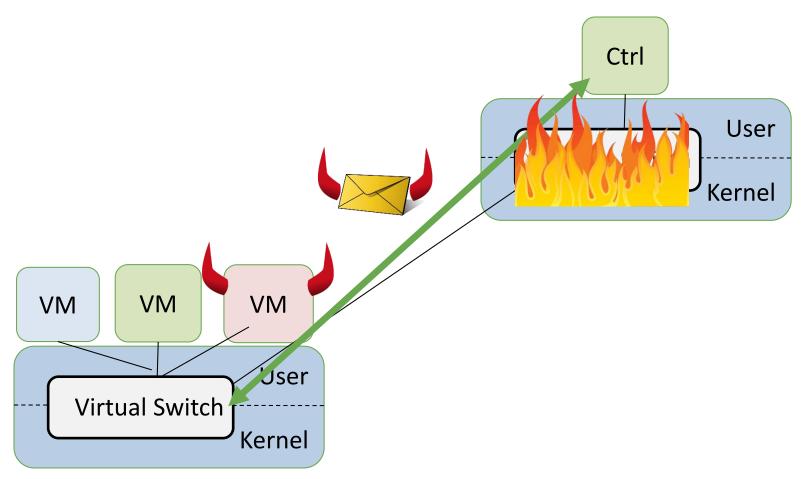
Number of parsed high-level protocols constantly increases...

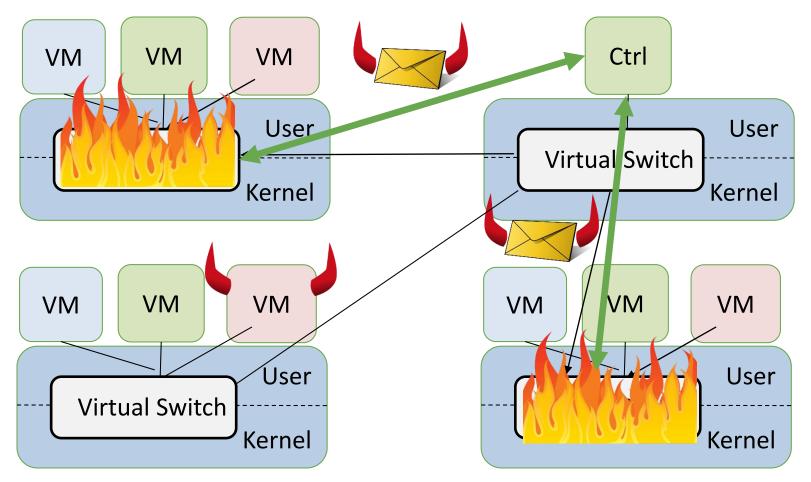
#### **Complexity: Parsing**











#### Challenge: How to provide better isolation efficiently?

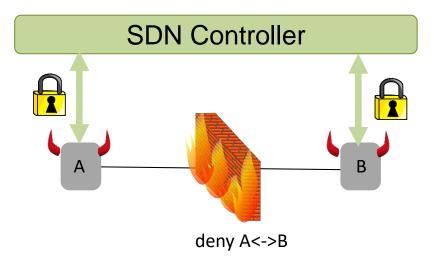
- Idea for better *isolation*: put vSwitch in a VM
- But what about *performance*?
- Or container?

VM
Virtual Switch

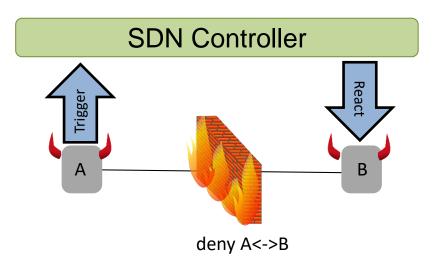
MTS: Bringing Multi-Tenancy to Virtual Switches Kashyap Thimmaraju, Saad Hermak, Gabor Retvari, and S. USENIX ATC, 2019.

# Also A Challenge: Covert and Side Channels

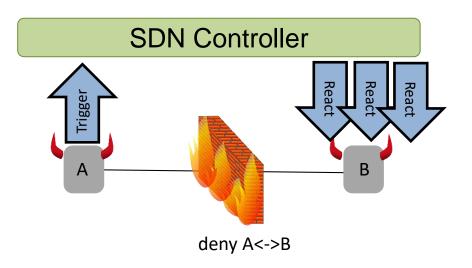
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- Controller may be attacked or exploited
  - By design, *reacts* to switch events, e.g., by packet-outs

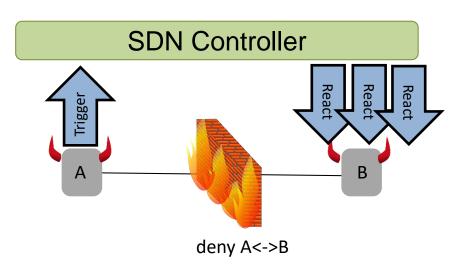


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  - Or even *multicast*: pave-path technique more efficient than hop-by-hop

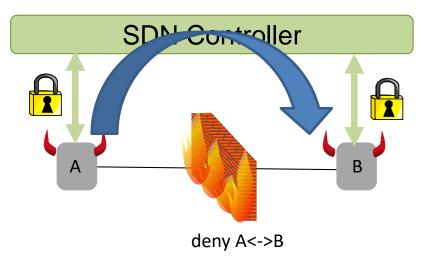


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May introduce *new communication paths* which can be used in unintendend ways!

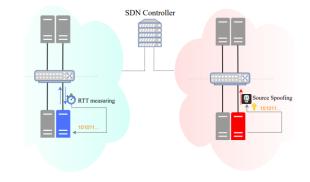


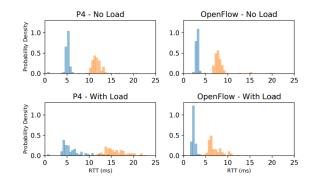
- In particular: new covert communication channels
  - E.g., exploit MAC learning (use codeword "0xBADDAD") or modulate information with timing
- May *bypass security-critical elements*: e.g., firewall in the dataplane
- *Hard to catch*: along "normal communication paths" and encrypted



#### **Example: Macchiato Side Channels**

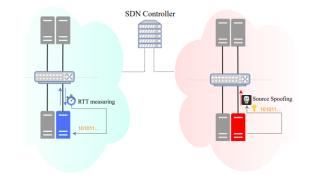
- Side channel *without malicious switches*: just between hosts
- Also leverages MAC learning for *mobility applications*: the fact that old rules associated with a MAC address are cleared when they marked in a new location
- The switch connected to the red host contacts the controller to report a new MAC address discovery. The controller interprets this that the blue host migrated to a new location and performs flow reconfigurations. Introduces *measurable differences* to the blue host that can be used repeatedly by the red host to modulate confidential data.
- Throughput can be improved with deep neural networks

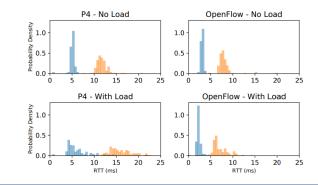




#### **Example: Macchiato Side Channels**

- Side channel *without malicious switches*: just between hosts
- Also leverages MAC learning for *mobility applications*: the fact that old rules associated with a MAC address are cleared when they marked in a new location
- The switch connected to the red host contacts the controller to report a new MAC address discovery. The controller interprets this that the blue host migrated to a new location and performs flow reconfigurations. Introduces *measurable differences* to the blue host that can be used repeatedly by the red host to modulate confidential data.
- Throughput can be improved with deep neural networks



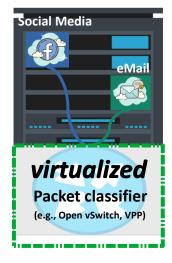


Macchiato: Importing Cache Side Channels to SDNs (Best Paper Award). Amir Sabzi, Liron Schiff, Kashyap Thimmaraju, Andreas Blenk, and Stefan Schmid. ANCS 2021.

# Another Challenge: Algorithmic Complexity Attacks

### Algorithmic Complexity Attacks

- Network dataplane runs many complex algorithms: may perform poorly under specific or *adversarial inputs*
- E.g., packet classifier: runs Tuple Space Search algorithm (e.g., in OVS)
- Can be exploited: adversary can *degrade performance* to ~10% of the baseline (10 Gbps) with only <1 Mbps (!) attack traffic</li>
- Idea:
  - Tenants can use the Cloud Management System (CMS) to set up their ACLs to access-control, redirect, log, etc.
  - Attacker's goal: send some *packet towards the virtual switch* that when subjected to the ACLs will *exhaust resources*



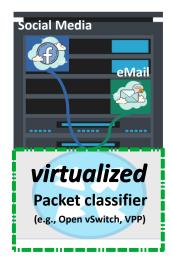
Tuple Space Explosion: A Denial-of-Service Attack Against a Software Packet Classifier. Levente Csikor et al. ACM CoNEXT, 2019.

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#### Use AI to find such attacks?!

Tuple Space Explosion: A Denial-of-Service Attack Against a Software Packet Classifier. Levente Csikor et al. ACM CoNEXT, 2019.



# Conclusion

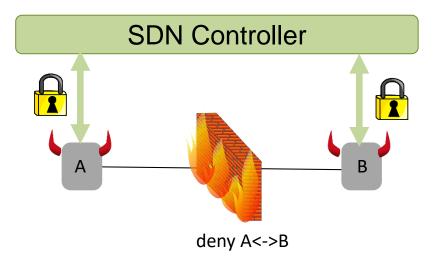
- Can we trust our networks today? Challenges, due to complexity, security assumptions and lack of tools
- Opportunities of emerging network technologies
  - Automation and programmability: new tools and improved network monitoring
- Challenges of emerging network technologies
  - New threat models: e.g., *propagate* worm in datacenter
  - Algorithmic complexity attacks: e.g., make virtual switch crawl



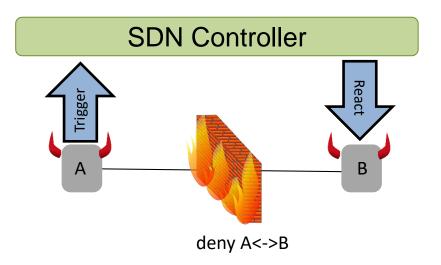
#### P-Rex: Fast Verification of MPLS Networks with Multiple Link Failures Jesper Stenbjerg Jensen, Troels Beck Krogh, Jonas Sand Madsen, Stefan Schmid, Jiri Srba, and Marc Tom Thorgersen. 14th International Conference on emerging Networking EXperiments and Technologies (CONEXT), Heraklion, Greece, December 2018. Macchiato: Importing Cache Side Channels to SDNs (Best Paper Award) Amir Sabzi, Liron Schiff, Kashyap Thimmaraju, Andreas Blenk, and Stefan Schmid. 16th ACM/IEEE Symposium on Architectures for Networking and Communications Systems (ANCS), Virtual Conference, December 2021. **NetBOA: Self-Driving Network Benchmarking** Johannes Zerwas, Patrick Kalmbach, Laurenz Henkel, Gabor Retvari, Wolfgang Kellerer, Andreas Blenk, and Stefan Schmid. ACM SIGCOMM Workshop on Network Meets AI & ML (NetAI), Beijing, China, August 2019. MTS: Bringing Multi-Tenancy to Virtual Switches Kashyap Thimmaraju, Saad Hermak, Gabor Retvari, and Stefan Schmid. USENIX Annual Technical Conference (ATC), Renton, Washington, USA, July 2019. Taking Control of SDN-based Cloud Systems via the Data Plane (Best Paper Award) Kashyap Thimmaraju, Bhargava Shastry, Tobias Fiebig, Felicitas Hetzelt, Jean-Pierre Seifert, Anja Feldmann, and Stefan Schmid. ACM Symposium on SDN Research (SOSR), Los Angeles, California, USA, March 2018. **Outsmarting Network Security with SDN Teleportation** Kashyap Thimmaraju, Liron Schiff, and Stefan Schmid. 2nd IEEE European Symposium on Security and Privacy (EuroS&P), Paris, France, April 2017. Preacher: Network Policy Checker for Adversarial Environments Kashyap Thimmaraju, Liron Schiff, and Stefan Schmid.

38th International Symposium on Reliable Distributed Systems (SRDS), Lyon, France, October 2019.

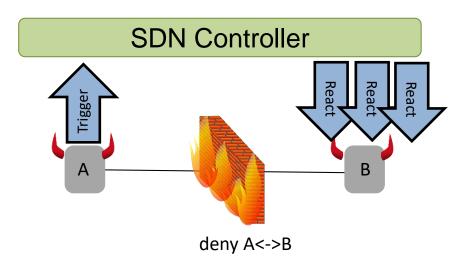
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  - By design, *reacts* to switch events, e.g., by packet-outs

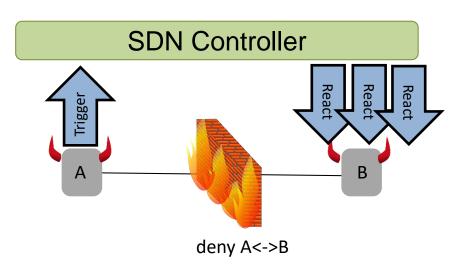


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  - Or even *multicast*: pave-path technique more efficient than hop-by-hop

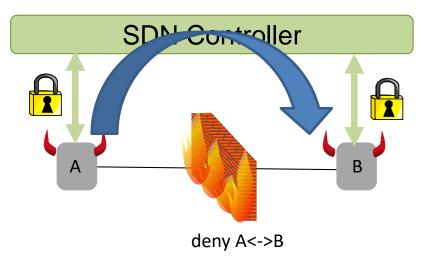


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May introduce *new communication paths* which can be used in unintendend ways!

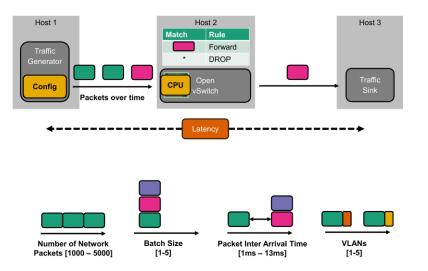


- In particular: new covert communication channels
  - E.g., exploit MAC learning (use codeword "0xBADDAD") or modulate information with timing
- May *bypass security-critical elements*: e.g., firewall in the dataplane
- *Hard to catch*: along "normal communication paths" and encrypted



#### NetBOA: Automated Performance Benchmarking

- Idea: *automate*! Generate different input, measure impact (e.g., latency)
  - Similar to *fuzzing*
- Different dimensions:
  - Packet size, inter-arrival time, packet type, etc.



NetBOA: Self-Driving Network Benchmarking Zerwas et al. ACM SIGCOMM Workshop on Network Meets AI & ML (NetAI), Beijing, China, August 2019.

#### **Baysian Optimization Approach**

- Complex systems (such as vSwitch) have complex behavior: e.g., sometimes sending less packets increases CPU load
  - Hard to find for humans

• Baysian optimization much faster than random baseline

