How Hard Can It Be?

Understanding the Complexity of Replica Aware Virtual Cluster Embeddings

<u>Carlo Fuerst</u> (TU Berlin, Germany), Maciek Pacut (University of Wroclaw, Poland) Paolo Costa (Microsoft Research, UK), Stefan Schmid (TU Berlin & T-Labs, Germany)

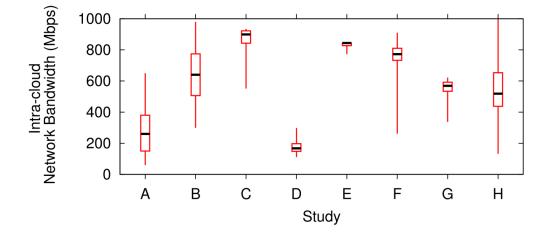
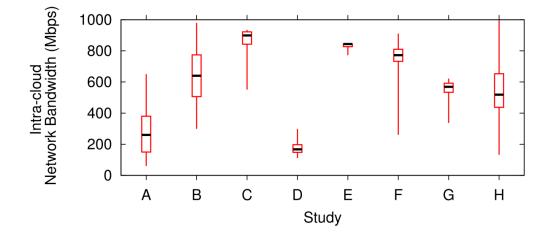


Figure 1: Percentiles $(1-25-50-75-99^{th})$ for intracloud network bandwidth observed by past studies.

Source: Ballani et al. [1] in Sigcomm'11

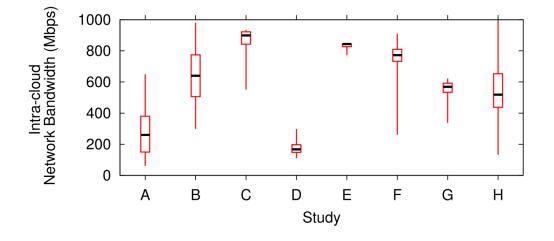


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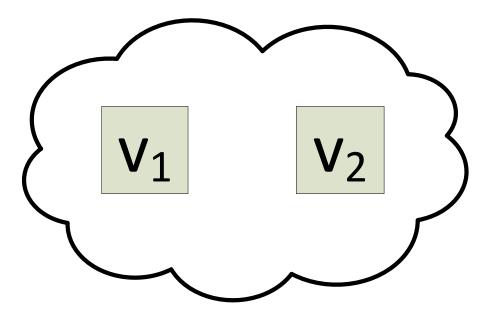
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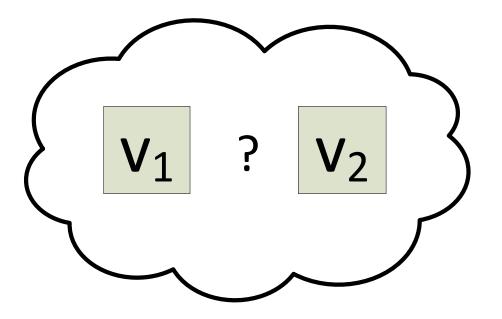
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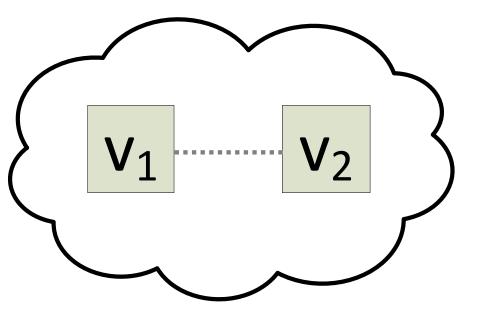
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Costs for the tenats become unpredictable



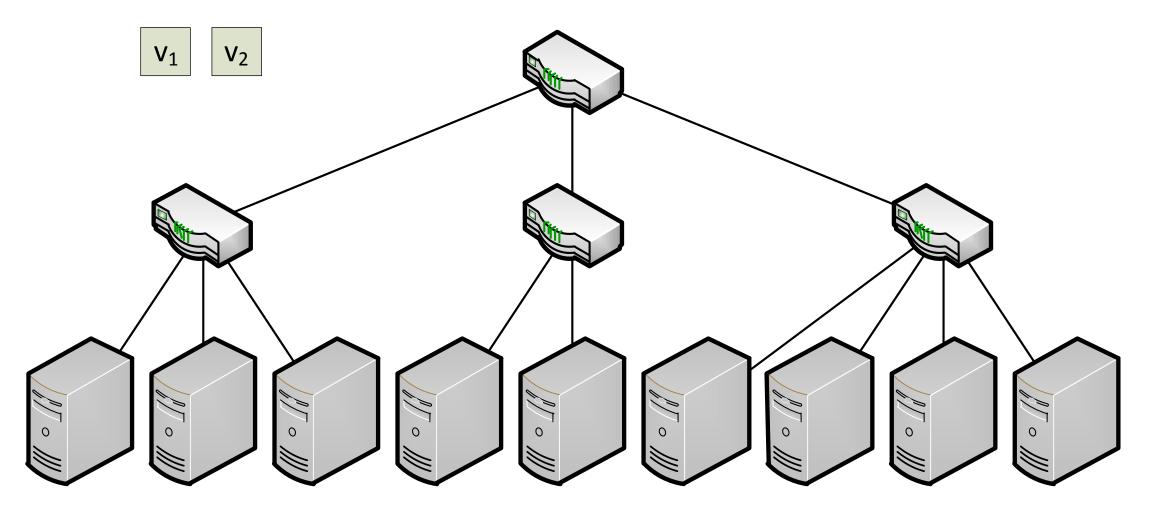




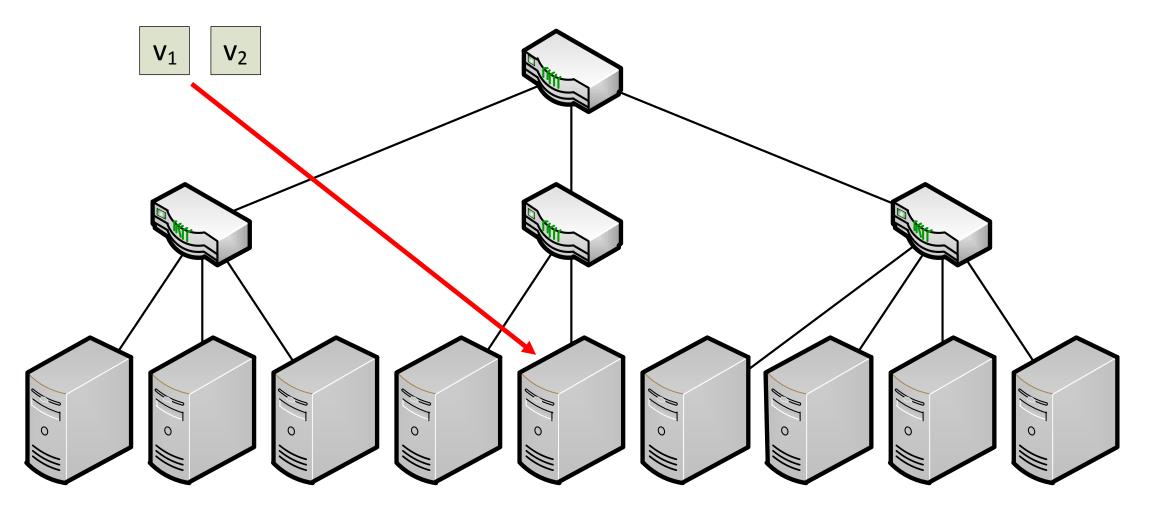
Remove the uncertainty by specifiying the bandwidth connecting the VMs

- Introduced by Ballani et al. [1]
- Provides absolute guarantees on VMs and network perfomance
- Specified by two parameters:
 - N the number of VMs
 - B the available bandwidth between VMs.

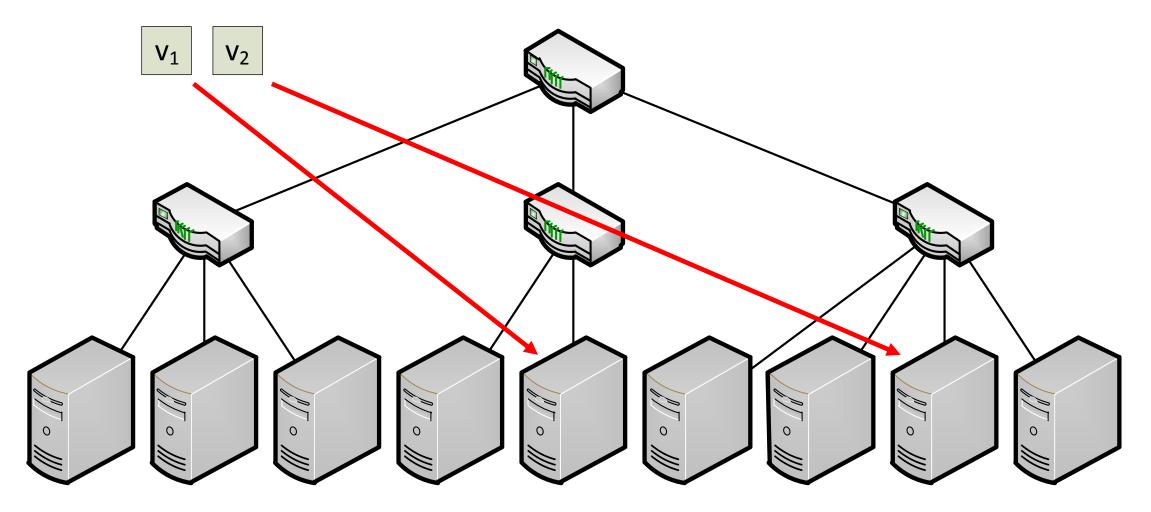
Embedding

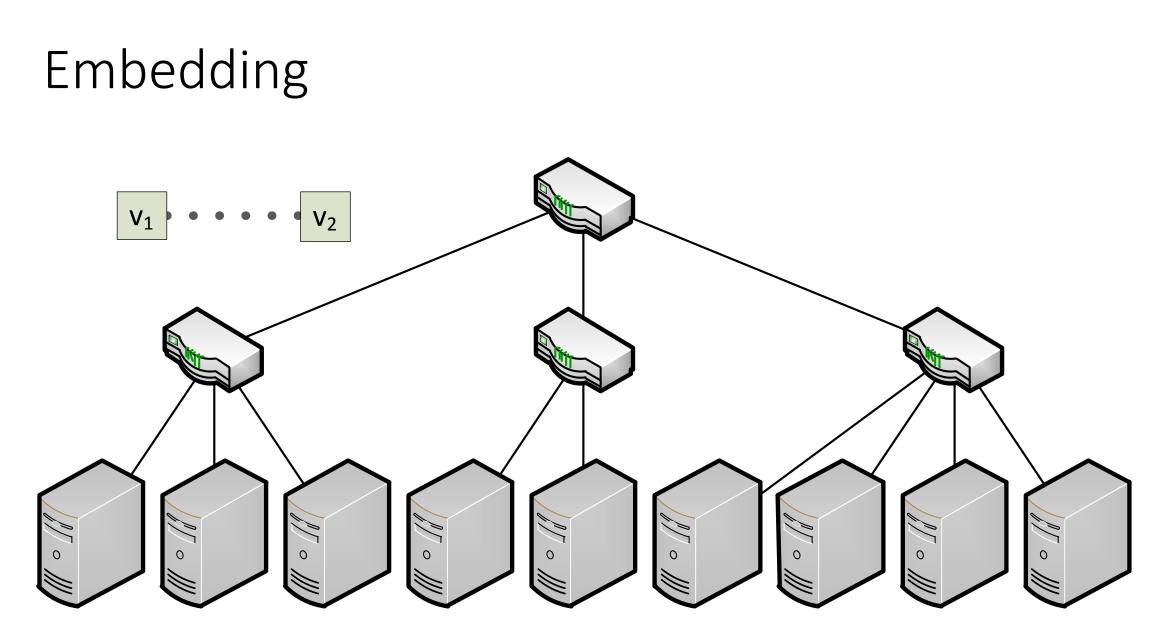


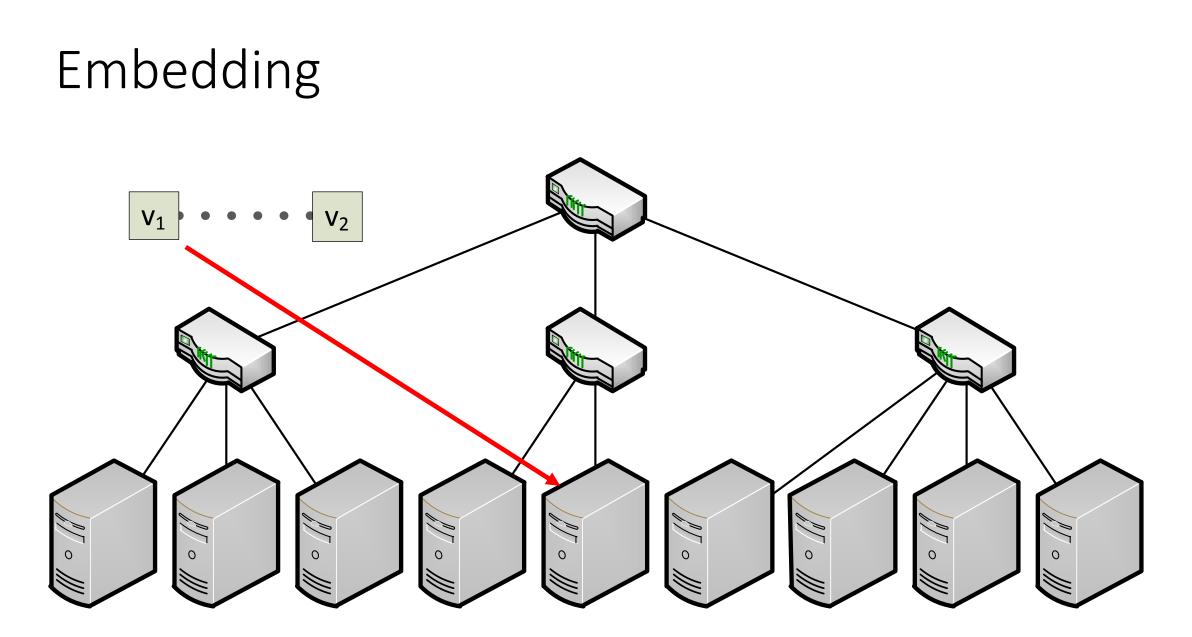
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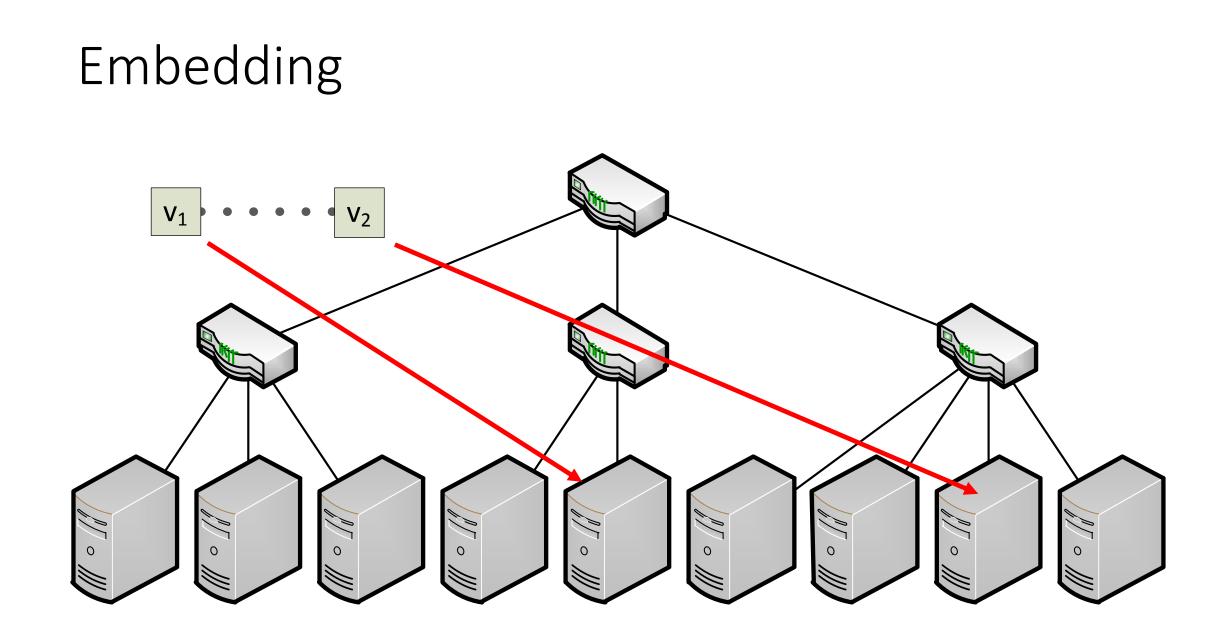


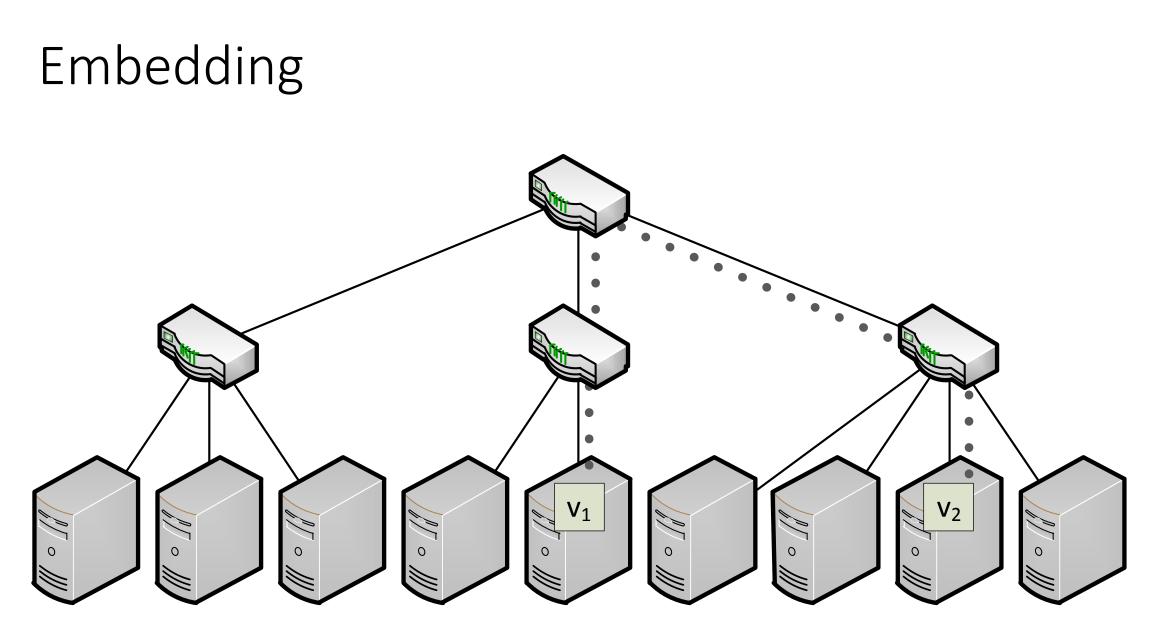
Embedding











Virtual Cluster Embedding Problem

- Subproblem of the NP-hard virtual network embeddding problem
- Good heuristics available
 - Ballani et al. [1] in Sigcomm'11
 - Xie et al. [3] in Sigcomm'12

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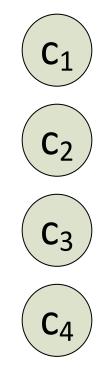
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The virtual cluster embedding problem is *not* NP-hard.[4]

Can the problem be solved efficiently with additional properties?

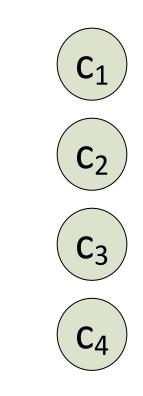
Example: MapReduce

1. Input is given by a set of atomic chunks



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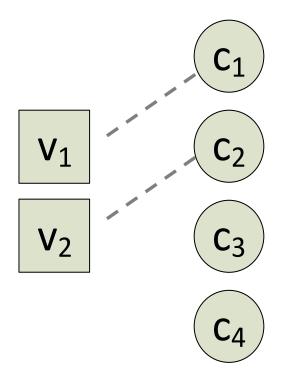
- 1. Input is given by a set of atomic chunks
- 2. Every chunk is processed by a map task



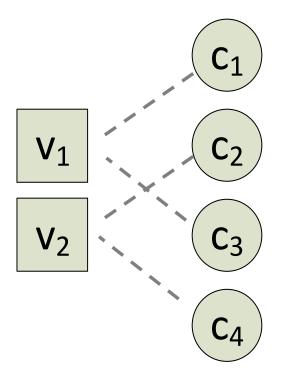
 V_1

 V_2

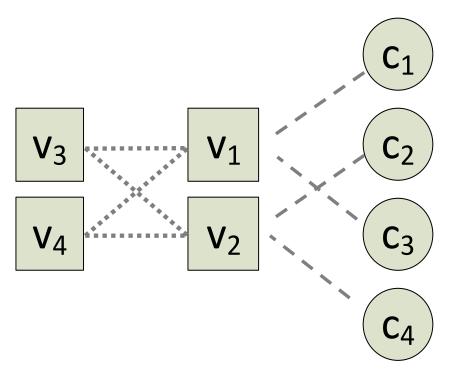
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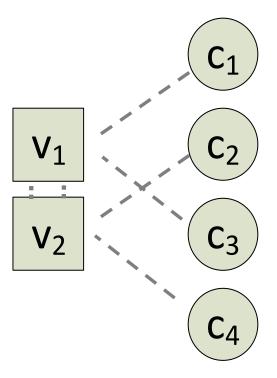
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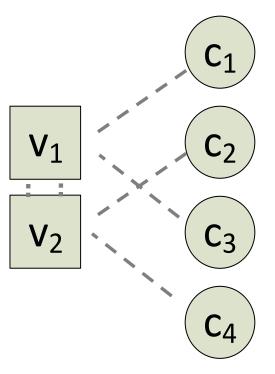
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- 4. Reduce tasks are executed

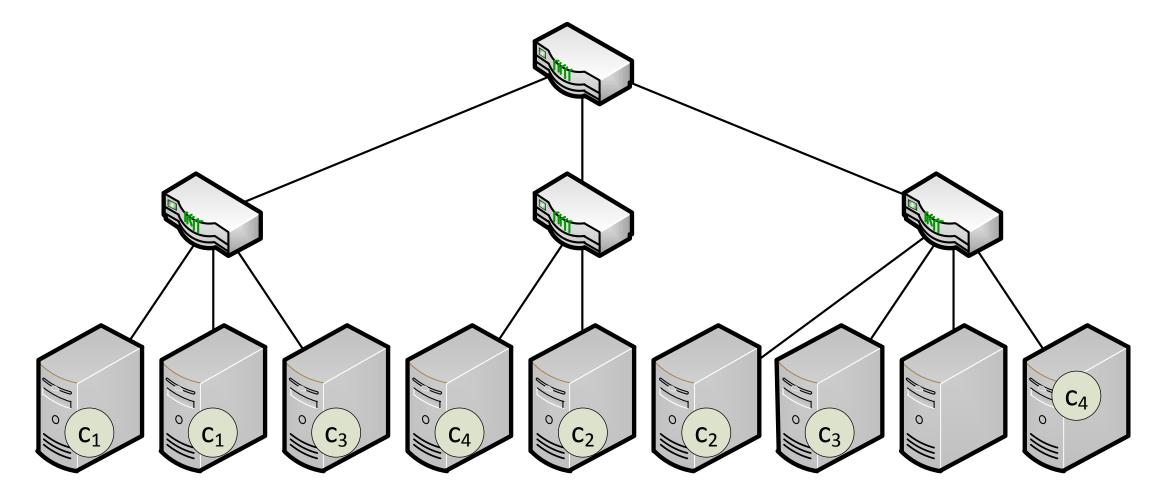


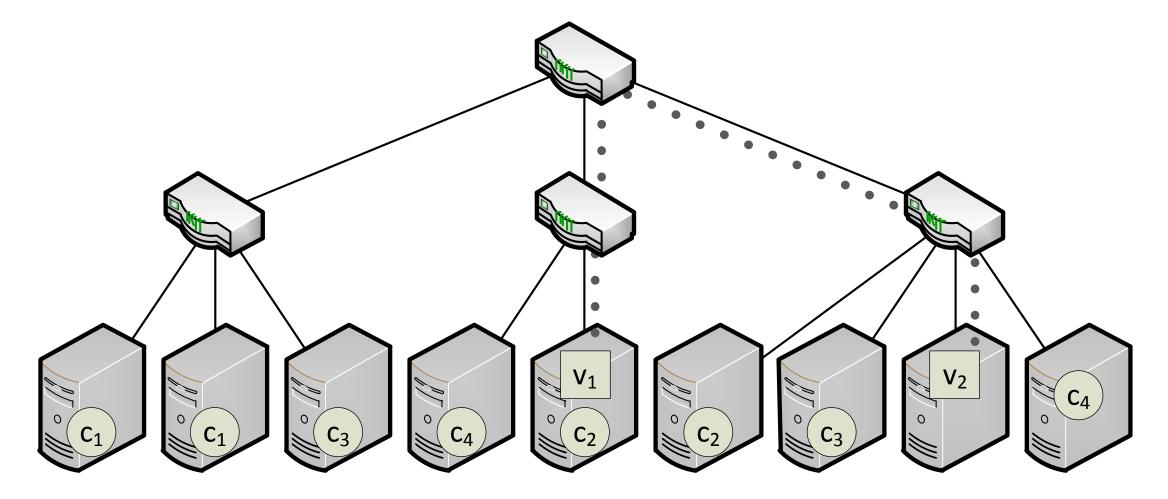
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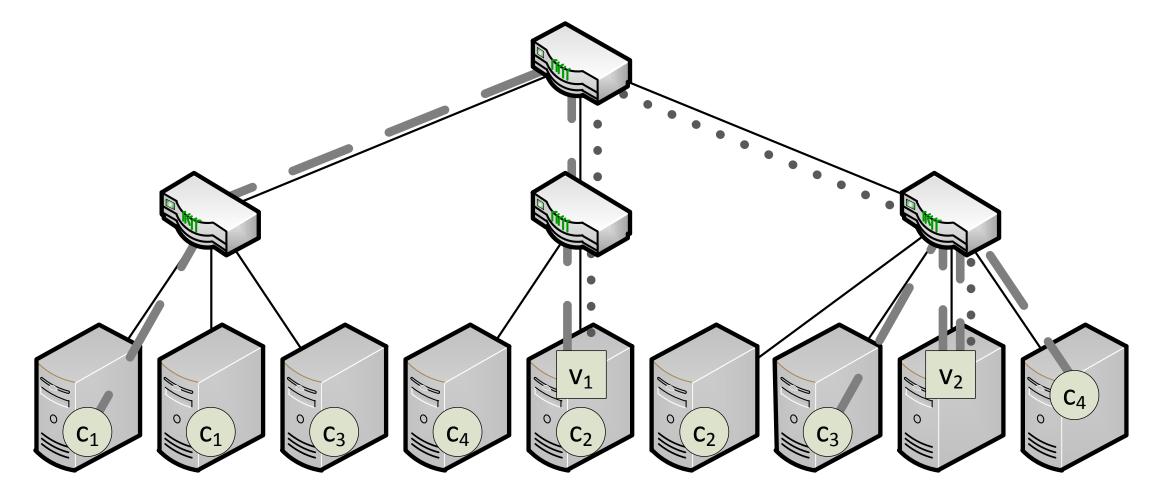
	C ₁
V ₁	C ₂
V ₂	C ₃
	C ₄

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	C ₁
v ₁	C ₂
V ₂	C_3
	C ₄

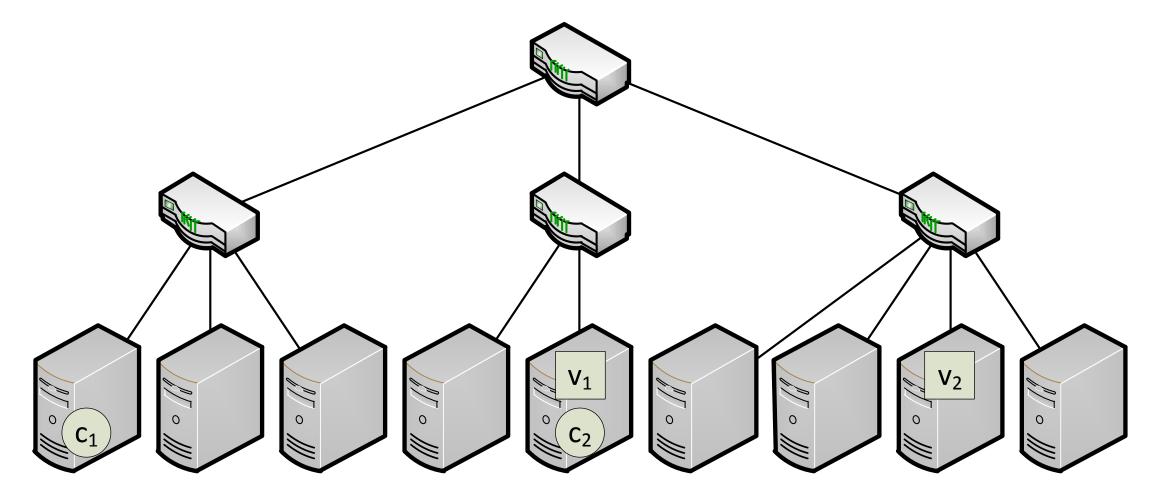




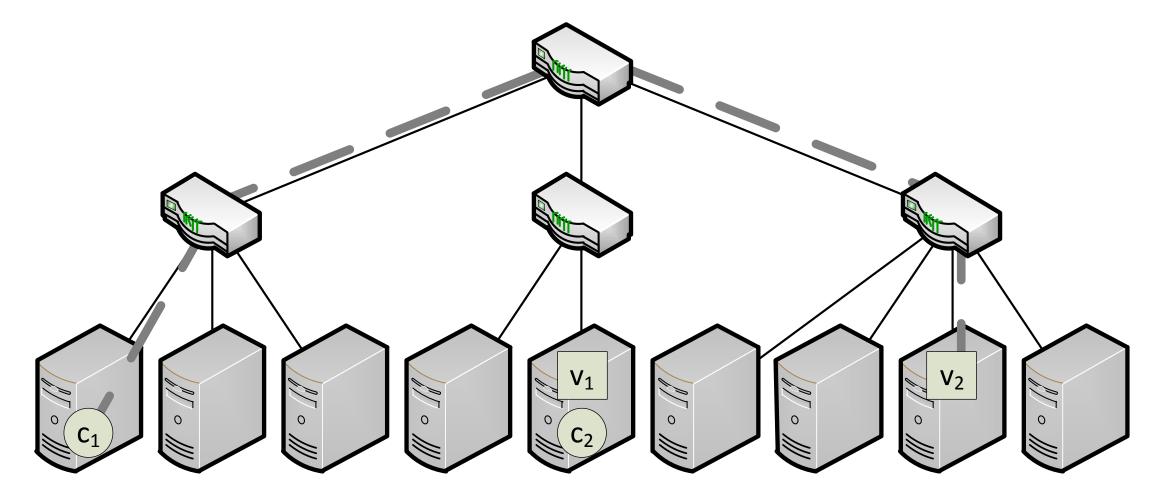




Basic Problem



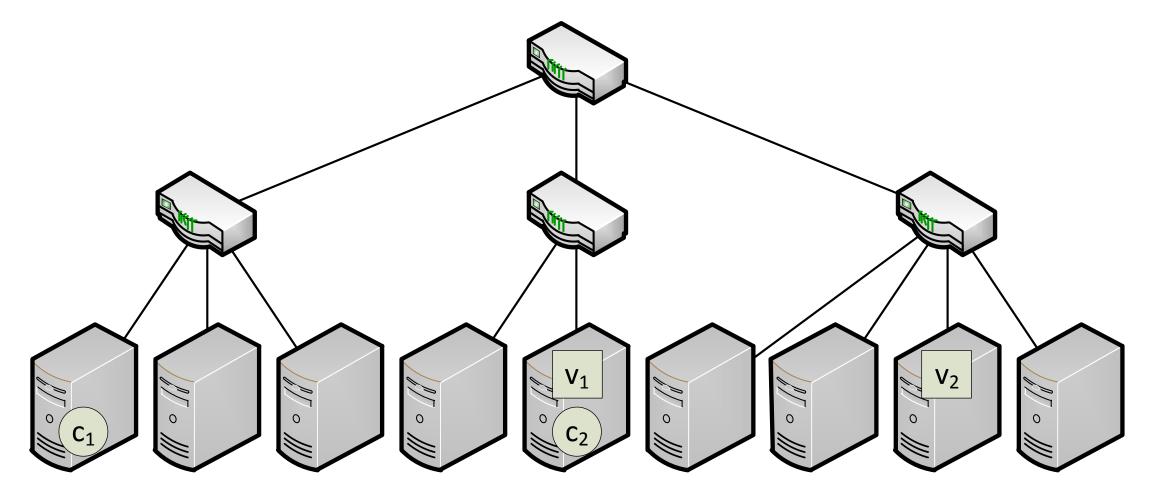
Basic Solution

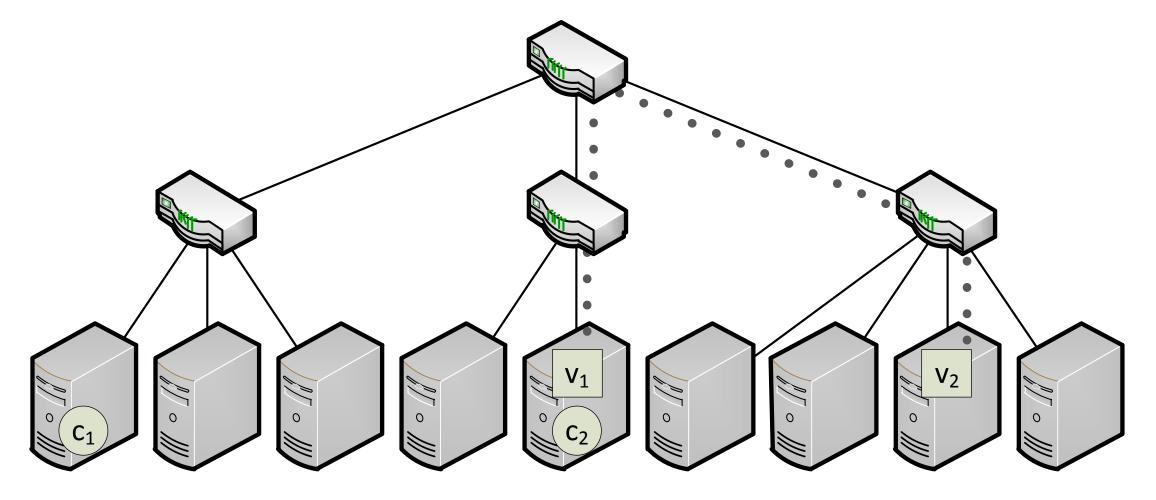


Problem Decomposition

The basic problem can be extended with:

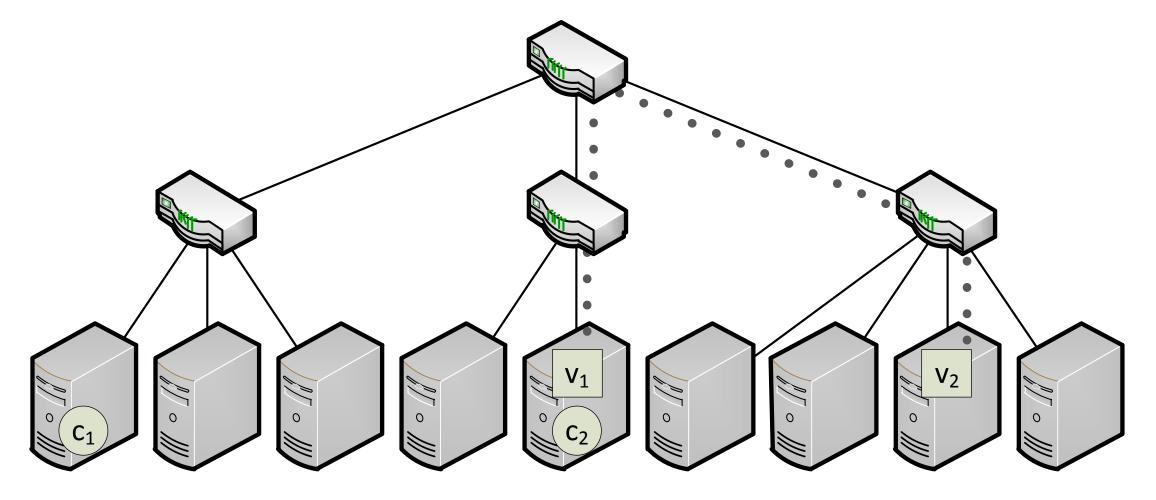
• VM interconnect (NI)

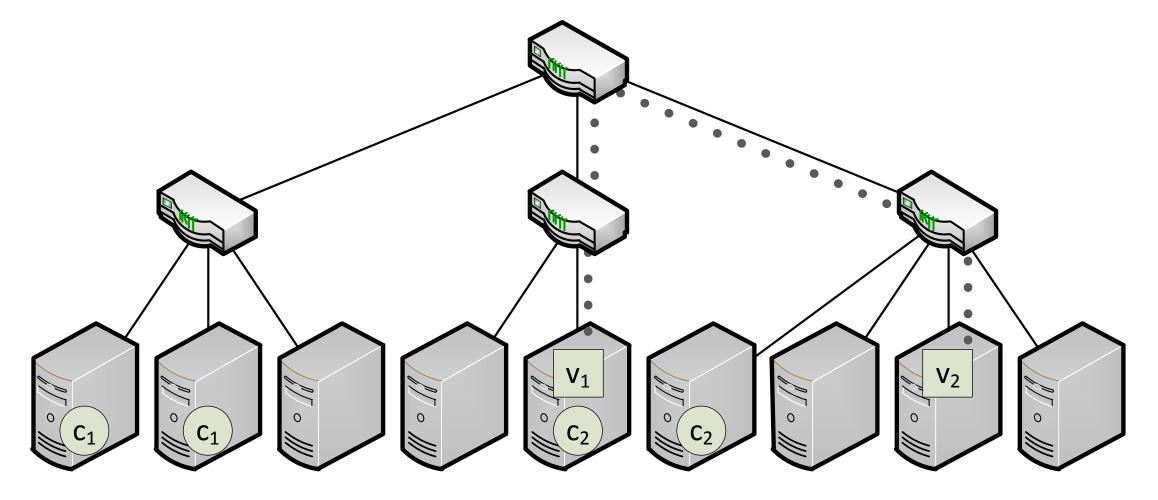




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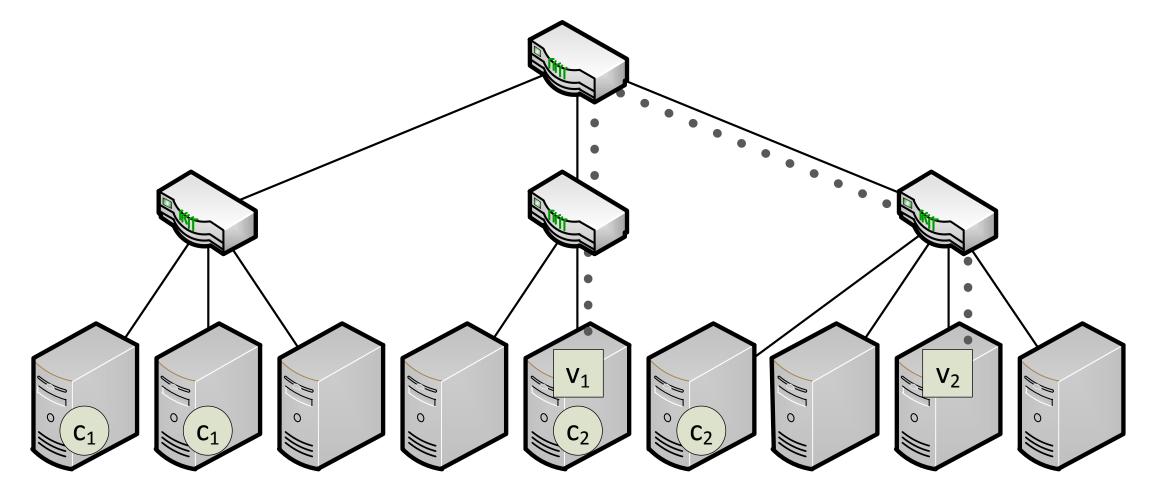
- VM interconnect (NI)
- Replica Selection (RS)

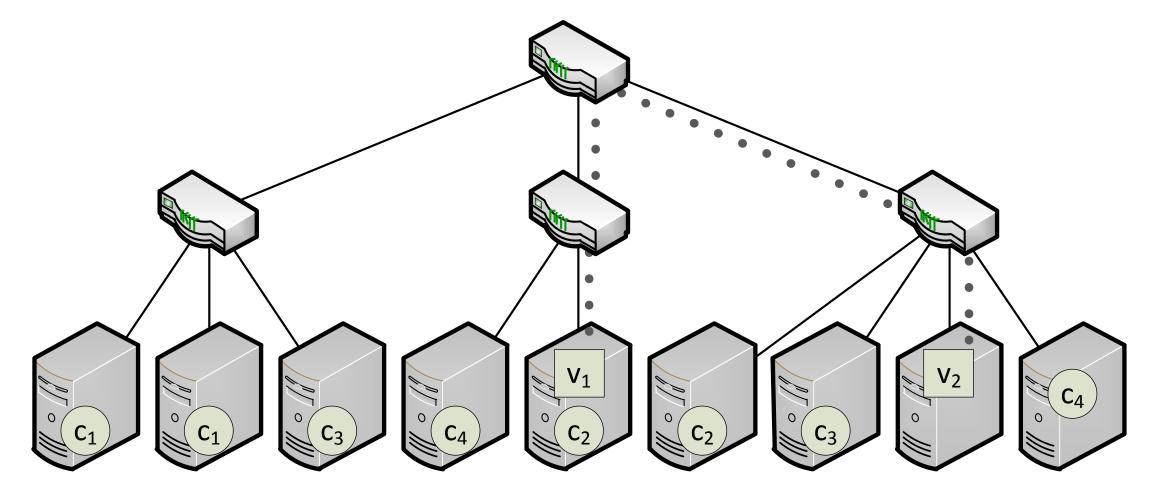




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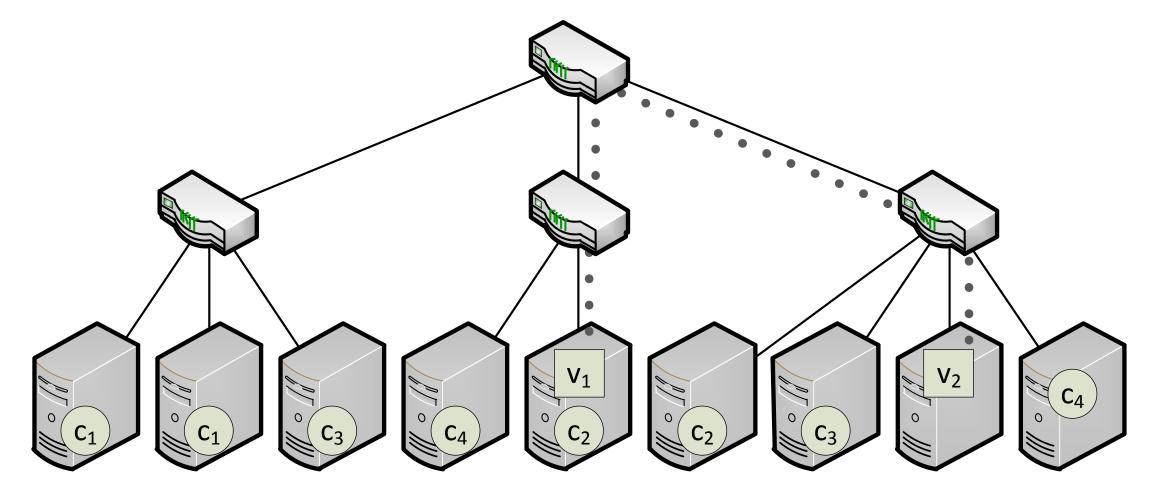
- VM interconnect (NI)
- Replica Selection (RS)
- Multiple Assignment (MA)

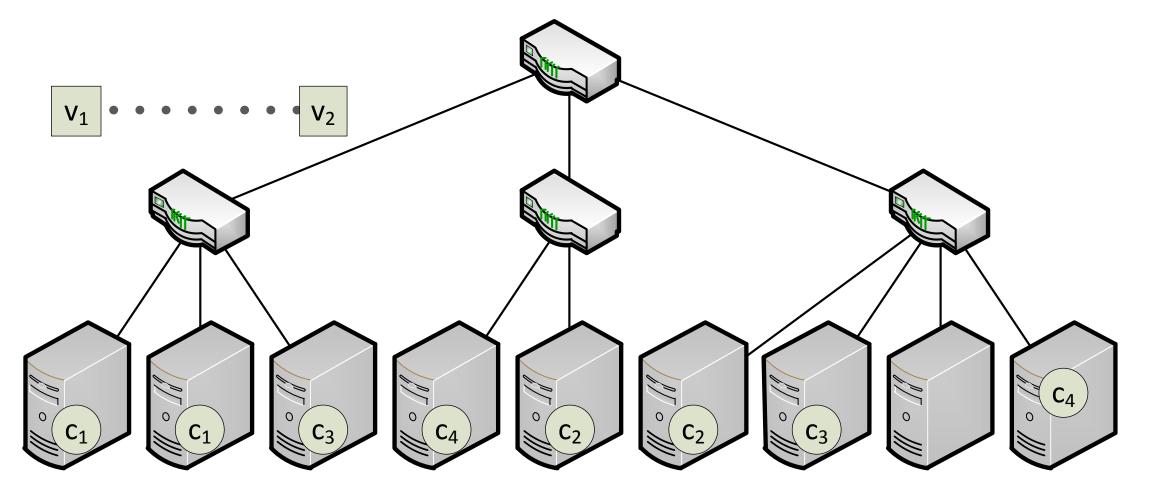




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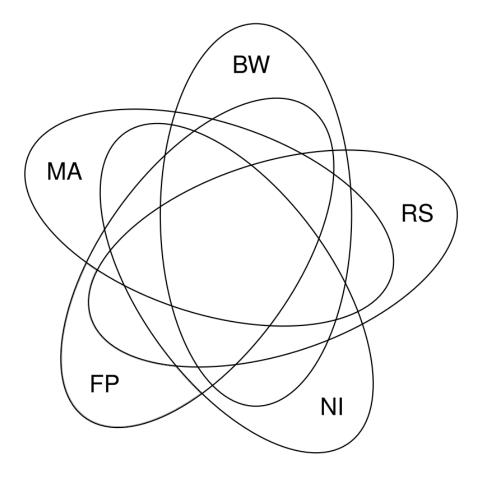
- VM interconnect (NI)
- Replica Selection (RS)
- Multiple Assignment (MA)
- Free placement of VMs (FP)



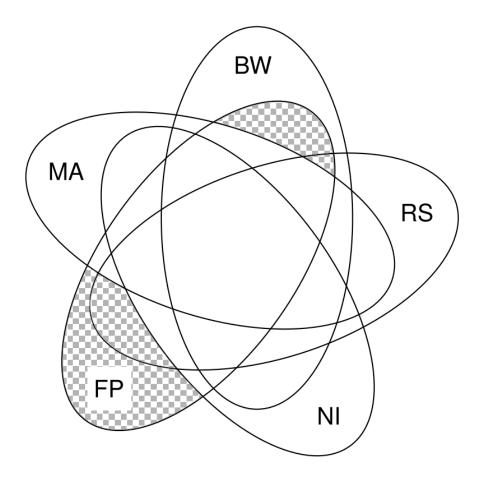


The basic problem can be extended with:

- VM interconnect (NI)
- Replica Selection (RS)
- Multiple Assignment (MA)
- Free placement of VMs (FP)
- Bandwidth Constraints (BW)

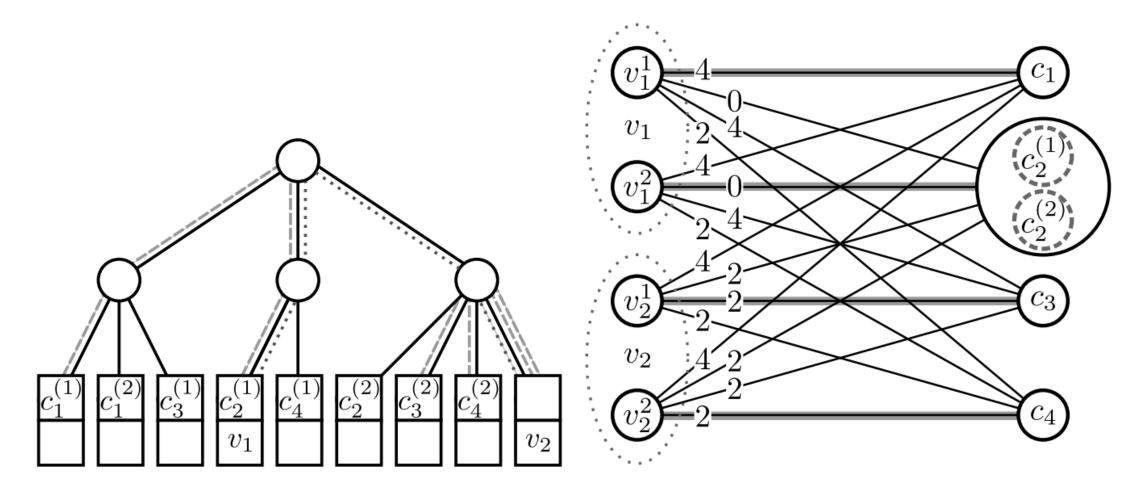


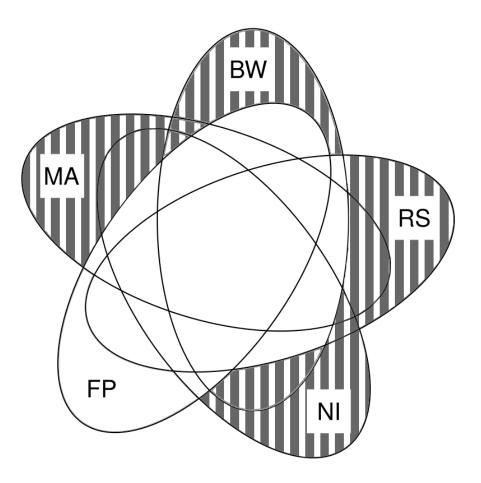
• Trivial problem identification



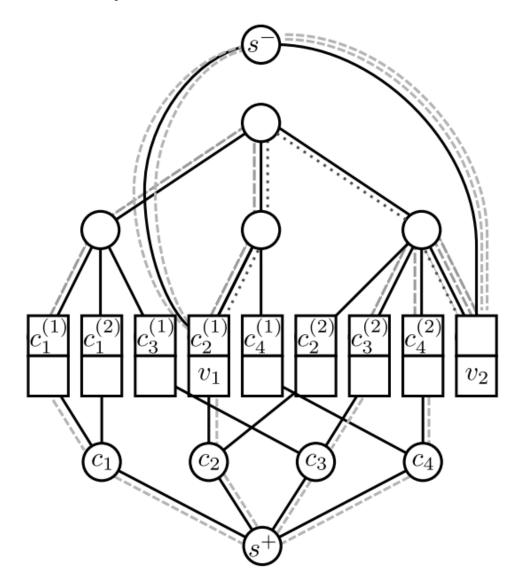
- Trivial problem identification
- Matching based algorithms

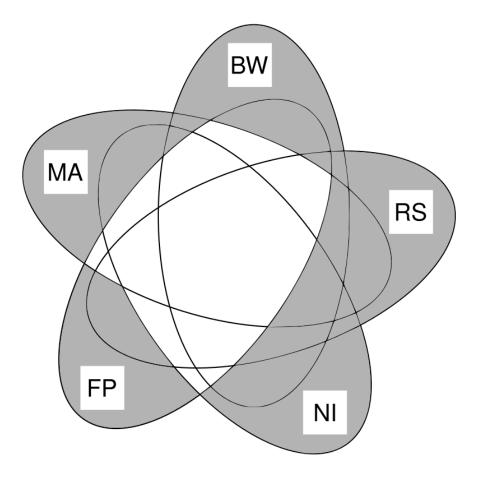
What is in the Paper?



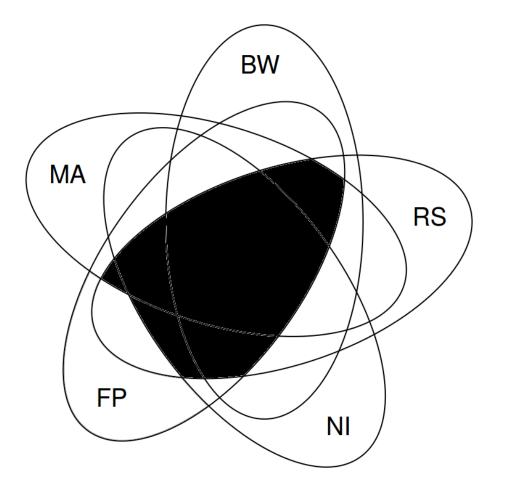


- Trivial problem identification
- Matching based algorithms
- Flow based algorithm

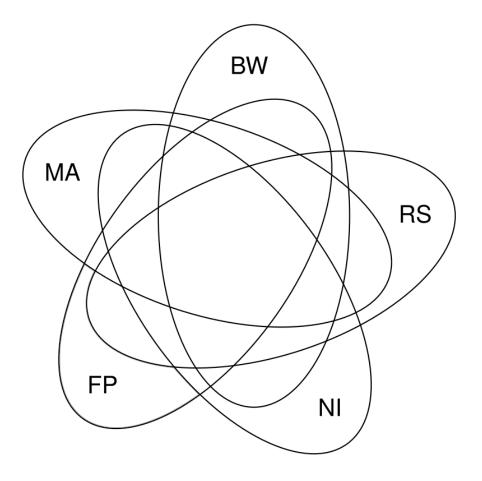




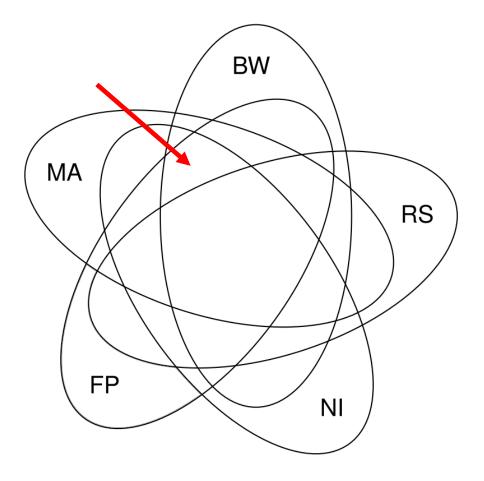
- Trivial problem identification
- Matching based algorithms
- Flow based algorithm
- Hardness results



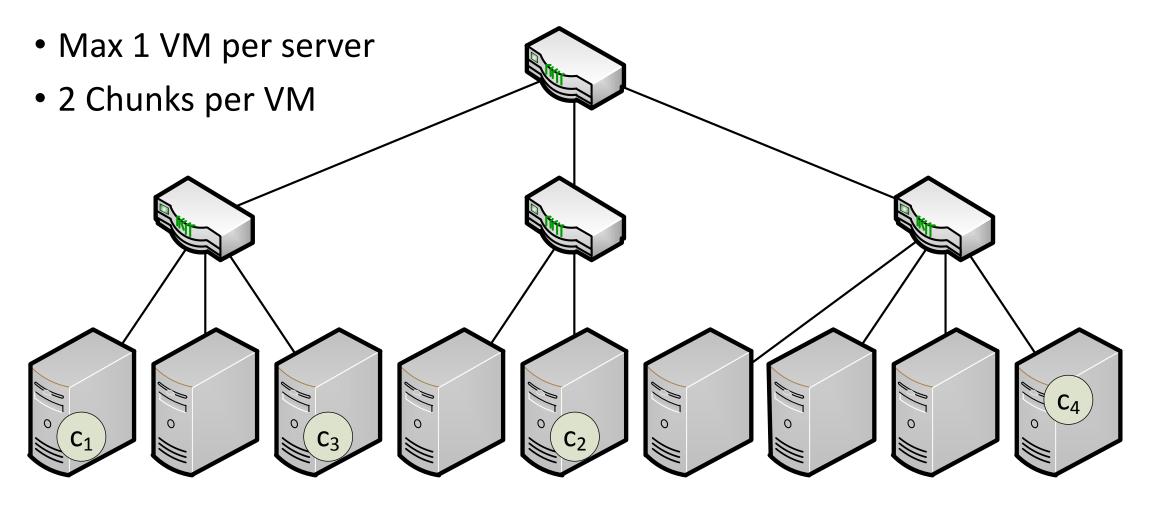
Everything but Replicas (MA + NI + FP + BW)



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- Create physical topology annotations in a bottom-up manner
- Start at the servers
- For each amount n of VMs in {0,...,N}
 - Set cost[n] to ∞ if n exceeds the servers capacity
 - Set cost[n] to the bandwidth costs of placing n VMs at the server



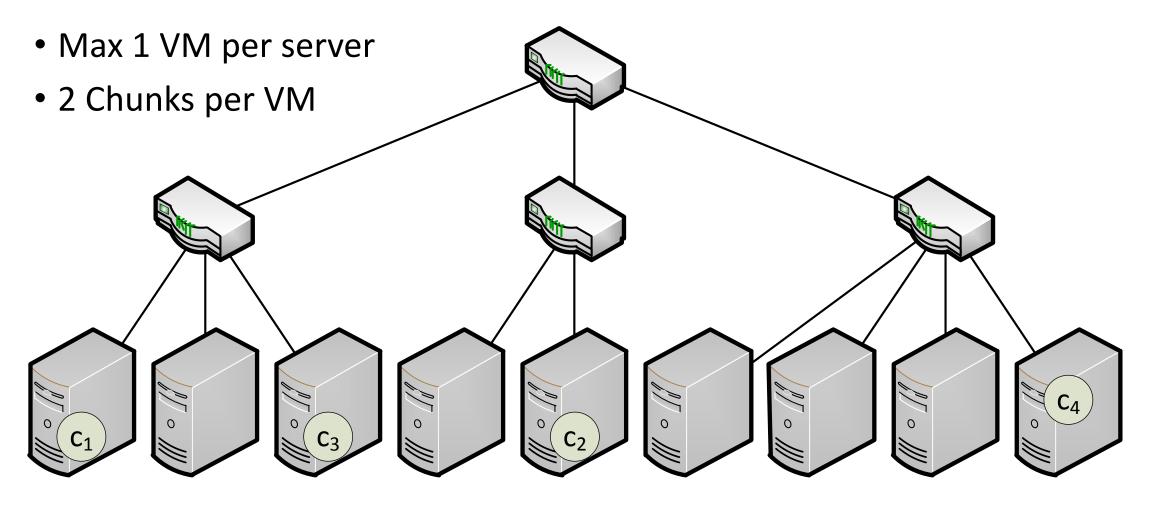
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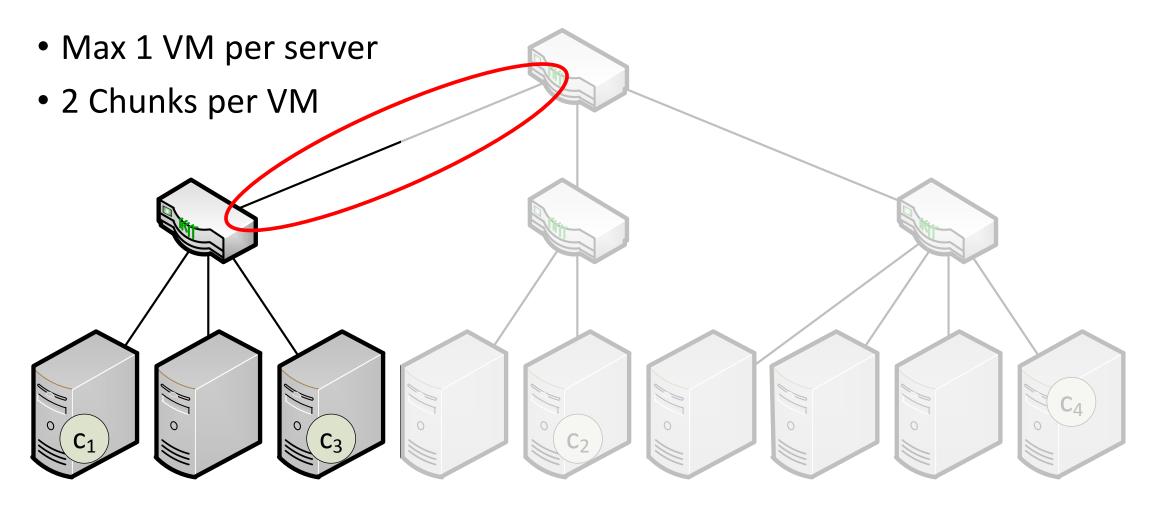
Ca

- Max 1 VM per server
- 2 Chunks per VM

 C_4

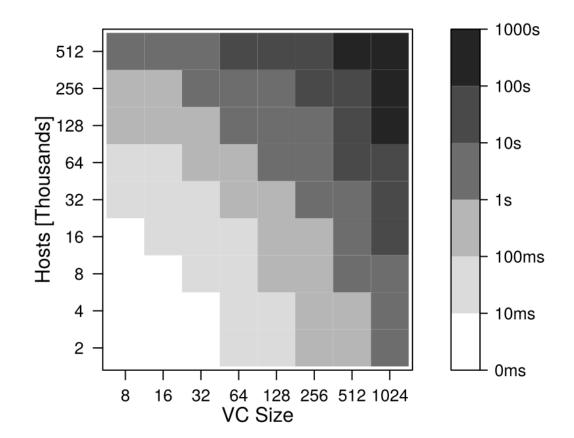
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 - Set cost[n] to the bandwidth costs of placing n VMs at the server
- For each switch and each amount of VMs in {0,...,N}
 - Set cost[n] to the sum of the cheapest combination of the children and add the costs for the bandwdith on the uplink



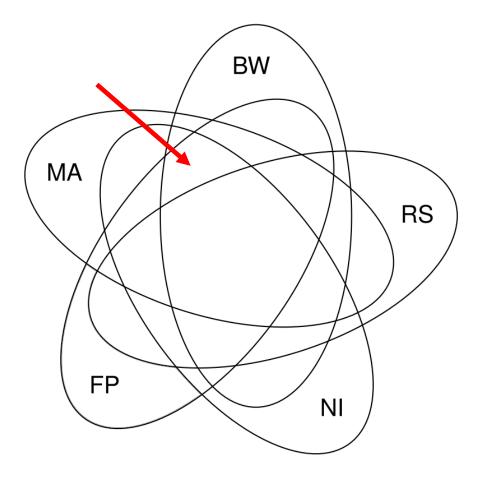


Runtimes

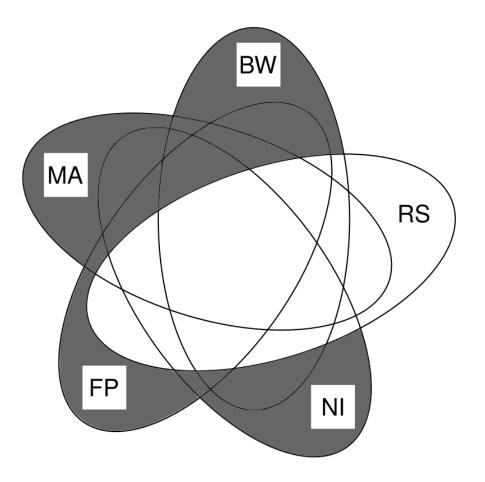
- Intel(R) Xeon(R) CPU L5420 @
 2.50GHzwith (single threaded)
- 512 MB
- openjdk-7
- Max 4 VMs per Server
- 3 Chunks per VM

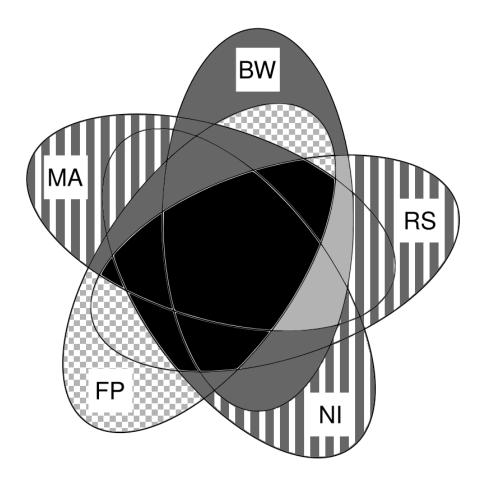


Which problems can be solved like this?



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Summary

- Virtual clusters provide dedicated resource guarantees
- Datalocality can be incorporated into the virtual cluster abstraction
- Problem decomposition into five properties
 - NP-hardness proofs for some property combinations
 - Algorithms for all other property combinations

References

[1] Ballani et al. **"Towards Predictable Datacenter Networks"** The ACM SIGCOMM Conference on Data Communication (SIGCOMM'11), Toronto, Canada, August 2011

[2] Chowdhury et al. "Managing Data Transfers in Computer Clusters with Orchestra." The ACM SIGCOMMConference on Data Communication (SIGCOMM'11)

[3] D. Xie, et al. "**The only constant is change: incorporating time-varying network reservations in data centers.**" The ACM SIGCOMMConference on Data Communication (SIGCOMM'12)

[4] M. Rost, et al. "**Beyond the stars: Revisiting virtual cluster embeddings**." ACM SIGCOMM Computer Communication Review 45.3(2015)