

1. Motivation:

- Reducing amount of communication in datacenters.
- Our approach: use VM migration for co-locating VMs which communicate frequently.
- Can be used to balance load and alleviate hotspots, or even to save energy.

2. Settings:

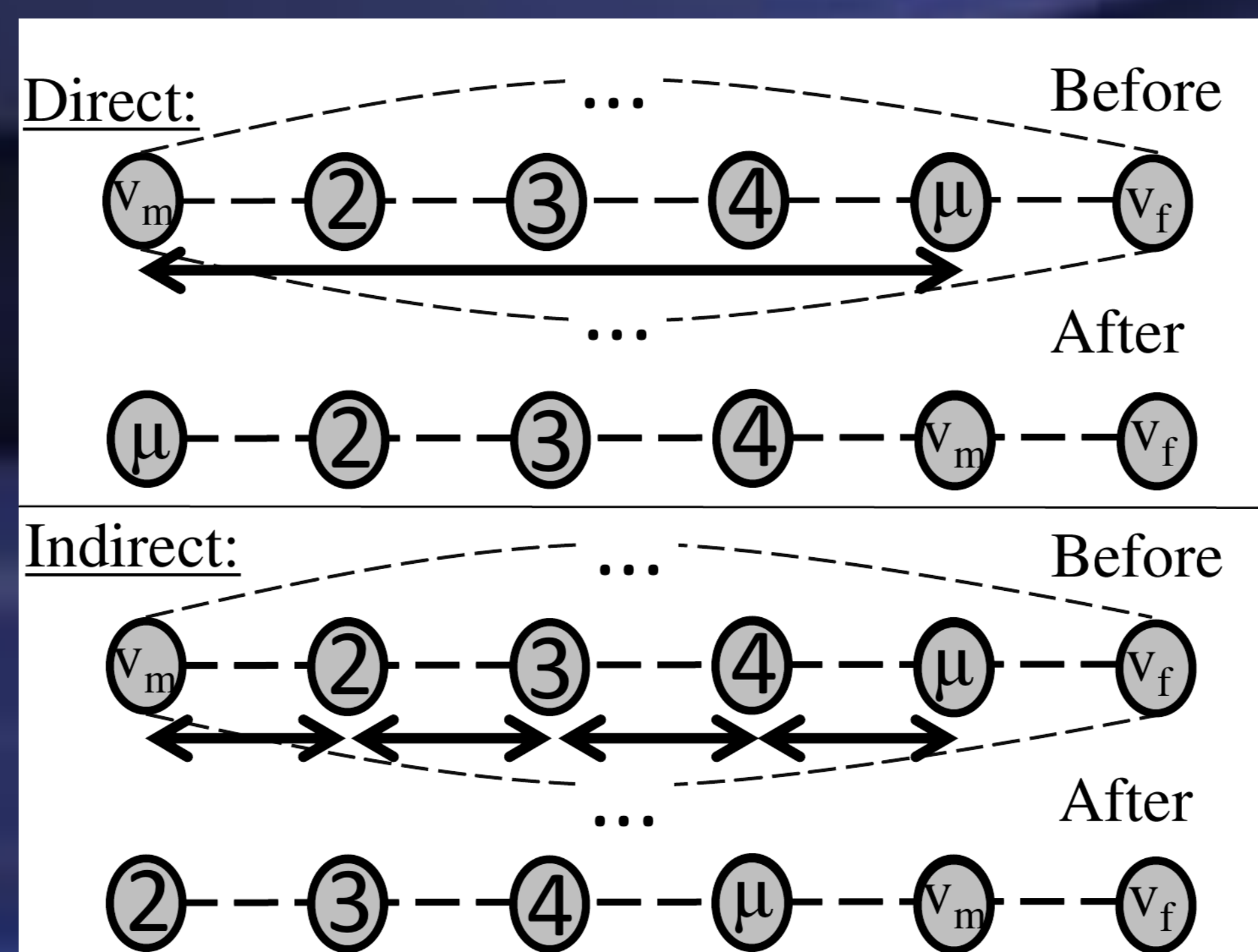
- A physical host graph H (the datacenter):
 $H = (M, L)$.
 M – physical machines.
 L – physical links.
- A guest graph G (Collection of tenants):
Application i : $G_i = (V_i, E_i, w_i)$.
 V_i – a set of VMs.
 $W_i(e_i)$ – frequency of communication
 $G = \cup G_i$

2. Model:

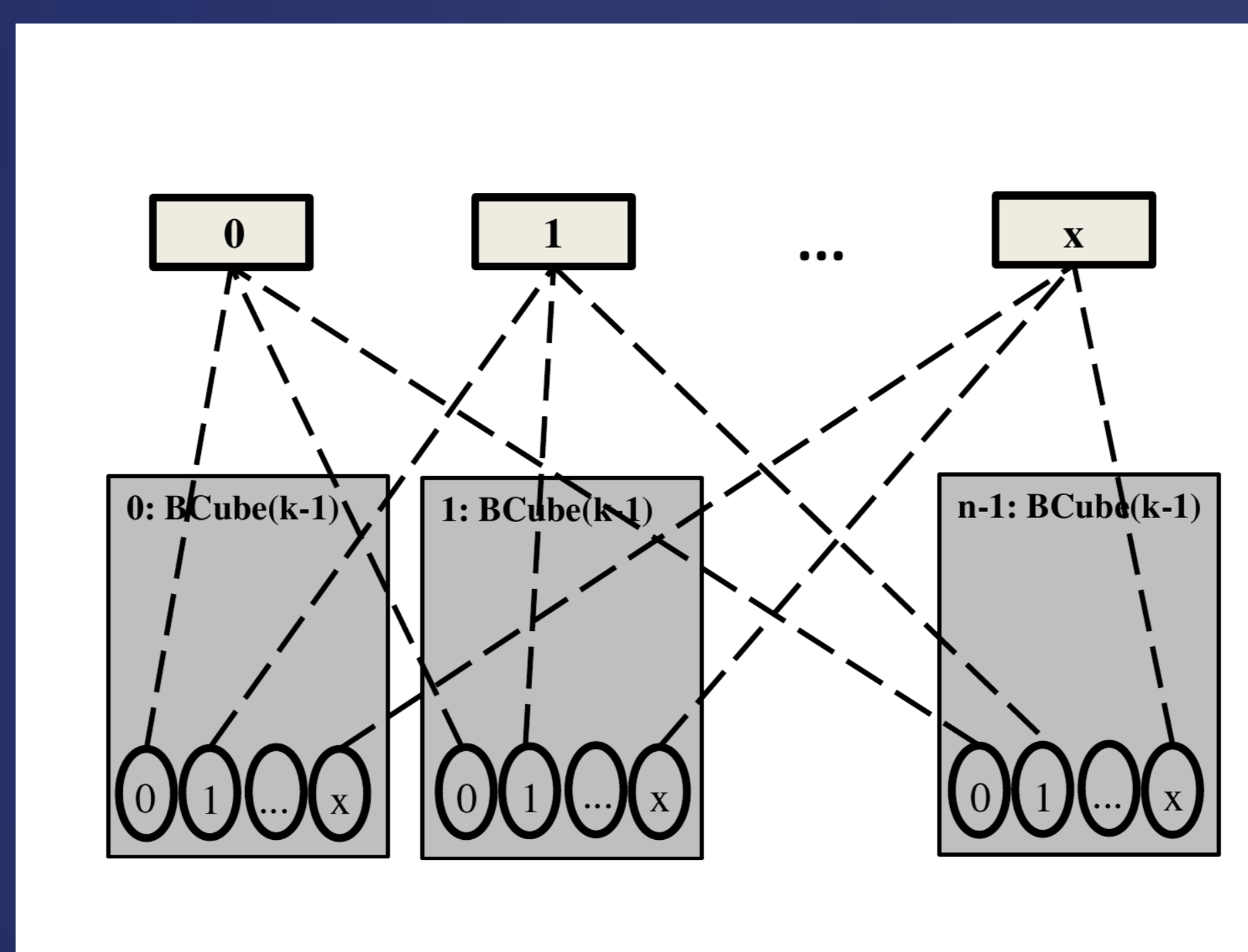
- $\sigma = \{\sigma_1 \dots \sigma_n\}$ a sequence of communication requests.
- Mapping function:
For $u \in V$, $\lambda_t(u)$ – the physical machine that hosts u at time t .
- ALG – migration algorithm can change the location of the VM.
- For $u, v \in V$, $d(u, v)$ – shortest path between $\lambda(u)$ and $\lambda(v)$.
- Amortized Cost: (H, ALG, σ)

$$= \frac{1}{|\sigma|} \sum_{t=1, (u,v) \in \sigma_t} d_t(u, v)$$

Direct Vs Indirect:



BCube Architecture:



4. Destination-Swap Algorithms:

- After communication request (u, v) – make u, v neighbors.
- Select “migrated VM” according to amortized cost.
- Select a “Destination” for migration.
- Destination Methods:
 - Random.
 - BestNeighbor.
 - MeetMiddle.
- Swapping Methods:
 - Direct.
 - Indirect.

5. Main Results:

- Random-Direct Converges to Best solution under a “matching” guest graph.
- Simulation shows great improvement for using BestNeighbor algorithm over simple random placement.
- Quick convergence time.

6. Simulation Results:

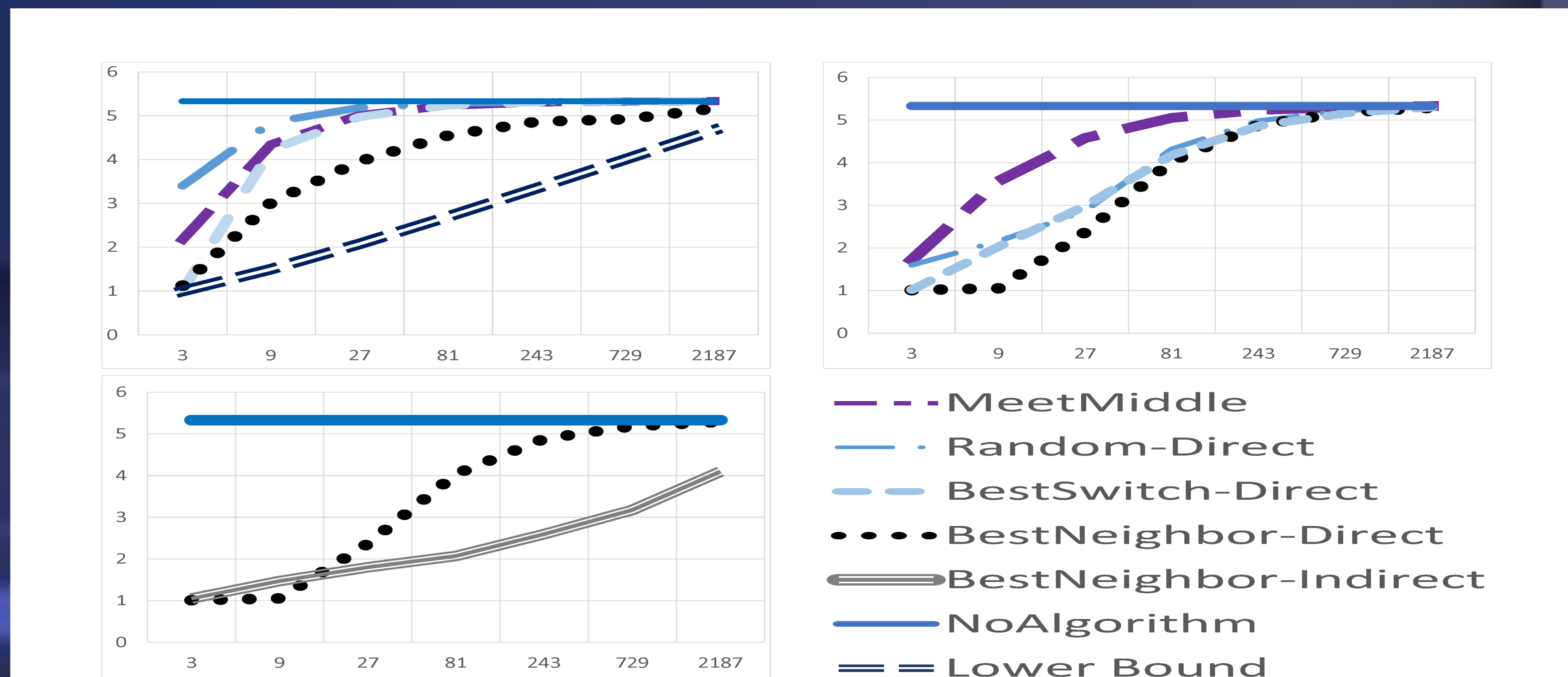


Fig. 2. Amortized communication cost as a function of the (tenant) guest graph size after 3m requests for host graph BCube(3,7) and for upper left: all-to-all communication $G^{(K_x)}$, upper right: one-to-all communication $G^{(S_x)}$, bottom left: $G^{(S_x)}$, Direct VS Indirect swaps for BESTNEIGHBOR