Learning-Augmented Online Algorithms Stefan Schmid (TU Berlin)

"We cannot direct the wind, but we can adjust the sails."

(Folklore)

Acknowledgements:





It's a great time to be an algorithm researcher!



Flexibilities everywhere in networked systems!

It's a great time to be an algorithm researcher!



Flexibilities everywhere in networked systems!
Enables self-adjusting systems: adapt to demand.

This Talk: Datacenters

Datacenters ("hyper-scale")

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NETFLIX



Interconnecting networks:
a critical infrastructure
of our digital society.



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Datacenters ("hyper-scale")

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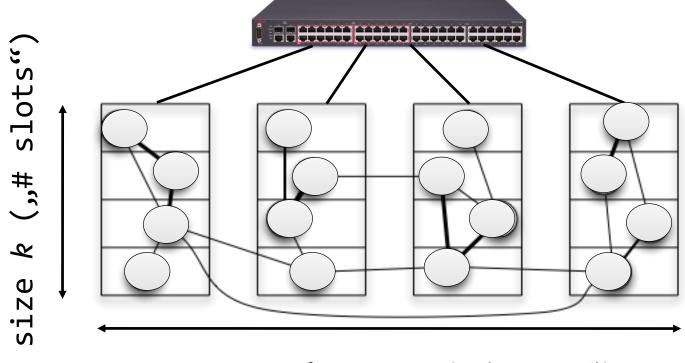


Interconnecting networks:
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Example 1: Scheduling

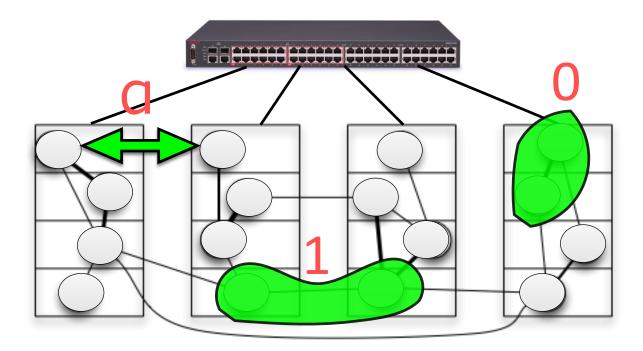
Online Re-Partitioning (Sigmetrics'19, SODA'21)



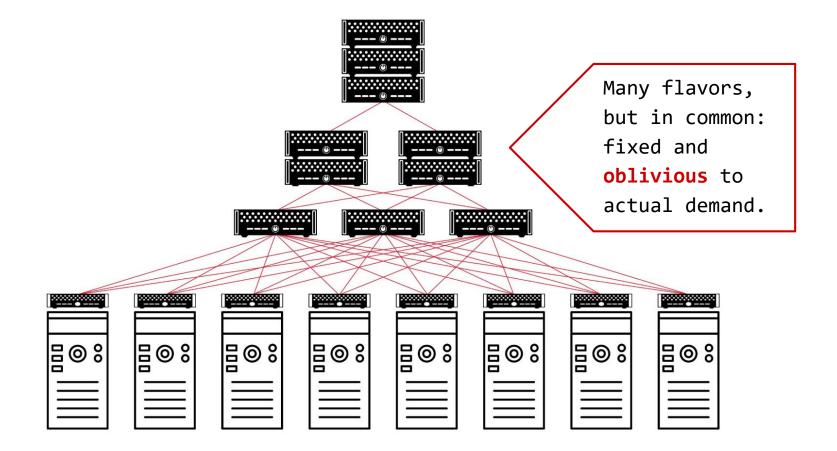
 ℓ servers ("clusters")

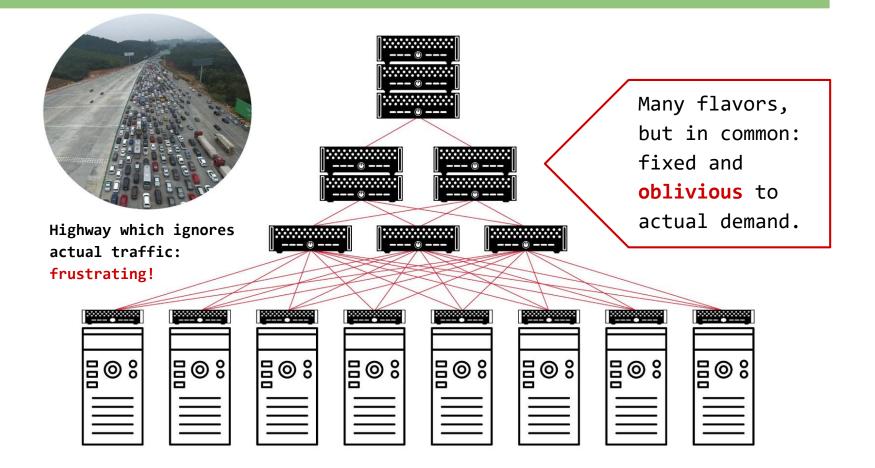
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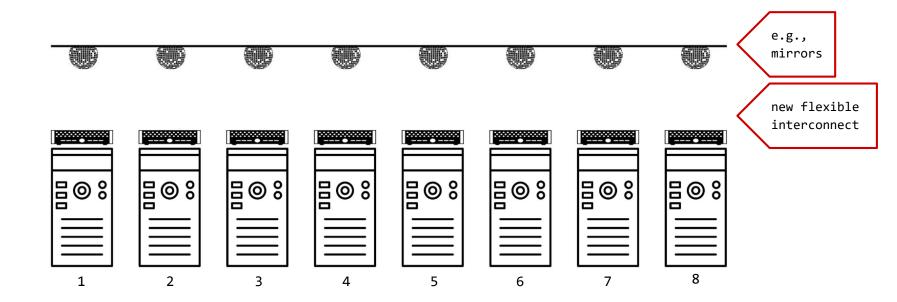


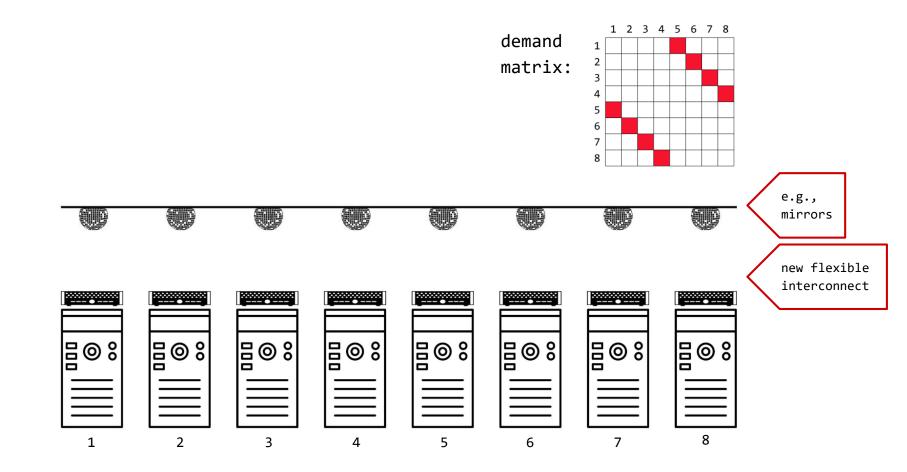
Migrate to reduce communication costs? Tradeoff!

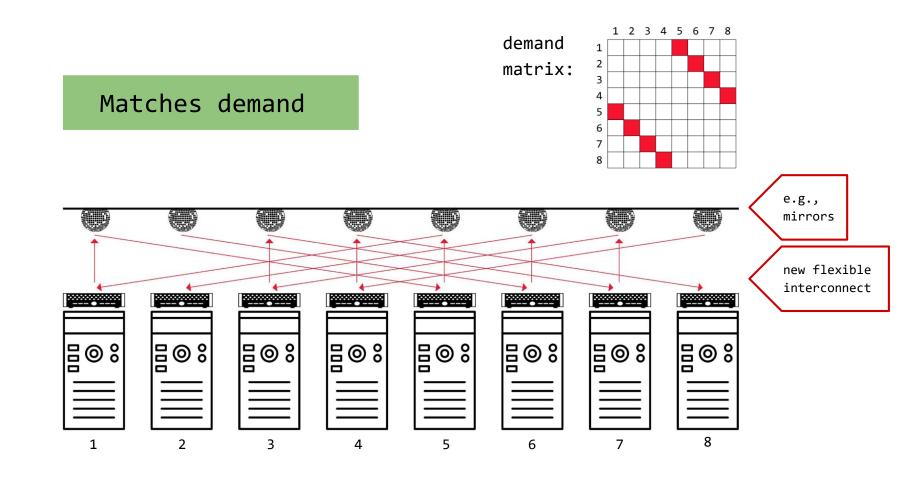


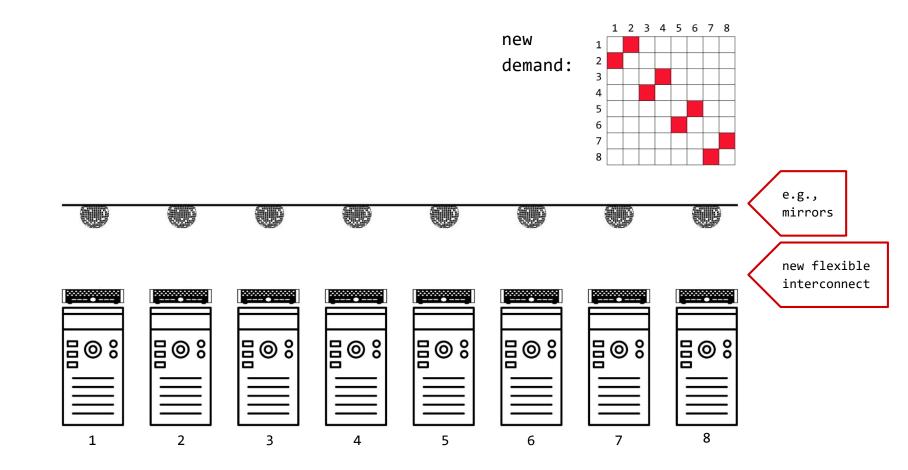


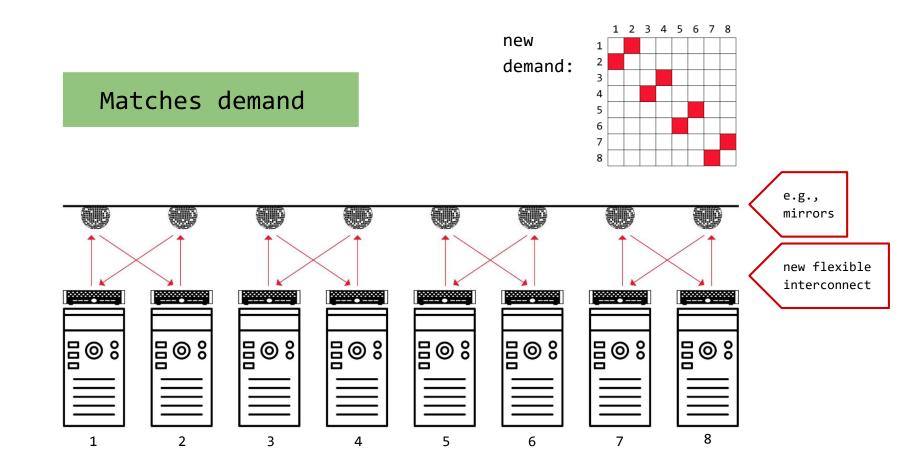
							
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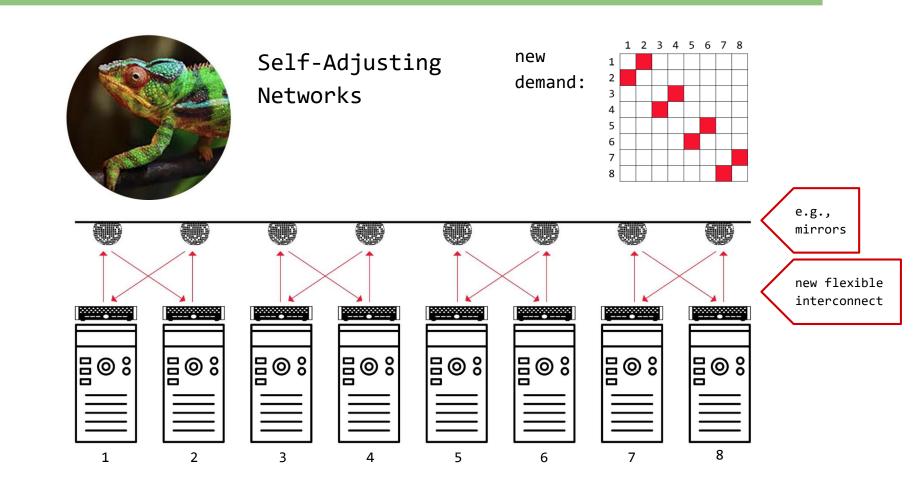


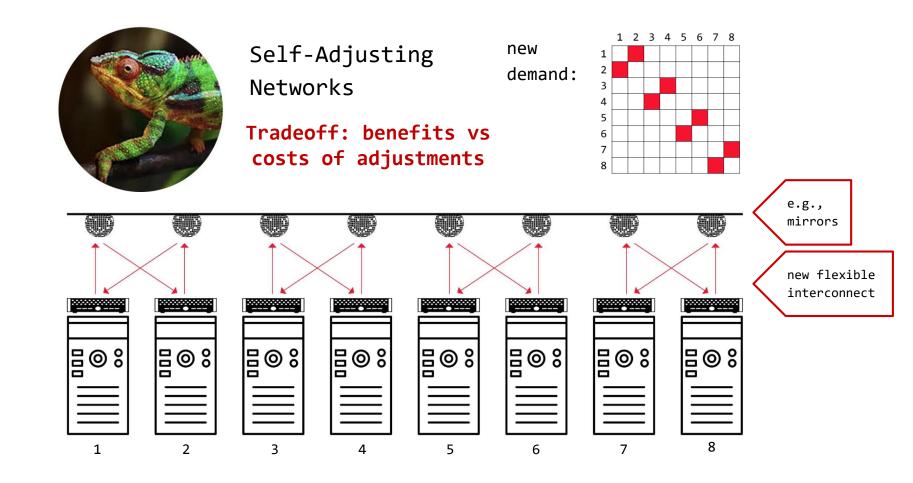








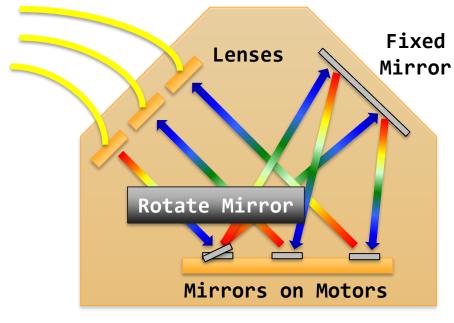




Underlying Technology

Optical Circuit Switch

---> Optical Circuit Switch rapid adaption of physical layer



\rightharpoonup Based on rotating mirrors

Optical Circuit Switch

By Nathan Farrington, SIGCOMM 2010

Competitive Ratio

Metric for Evaluating Self-Adjusting Systems

$\rho = \max_{\sigma} \text{Cost}_ON(\sigma)/\text{Cost}_OFF(\sigma)$

Competitive Ratio

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Competitive Ratio

Metric for Evaluating Self-Adjusting Systems

$$\rho = m_{\sigma} \cos (\sigma) / \cos (\sigma)$$



Too conservative? Demand often not "worst case".

Too Conservative?

Much Structure in the Demand

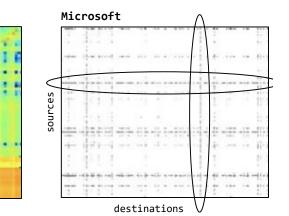
Empirical studies:

destinations

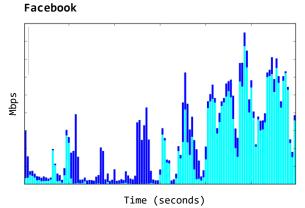
Facebook

sources

traffic matrices sparse and skewed



traffic bursty over time

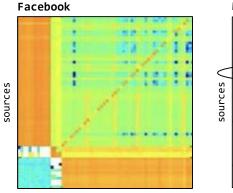


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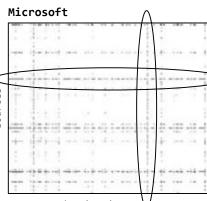
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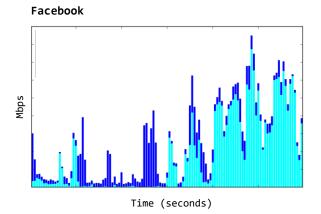


destinations



destinations

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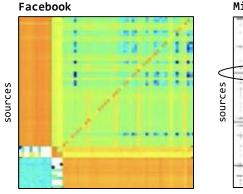
The hypothesis: can be learned/predicted/exploited.

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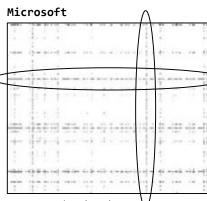
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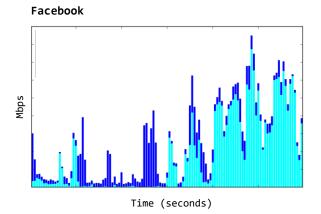


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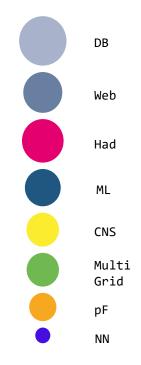
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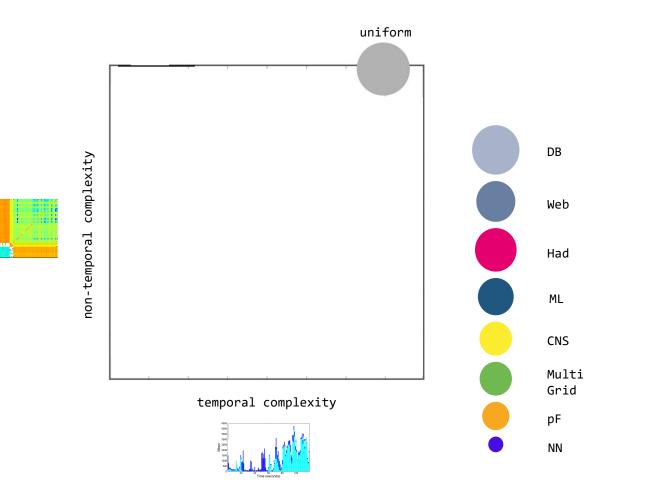


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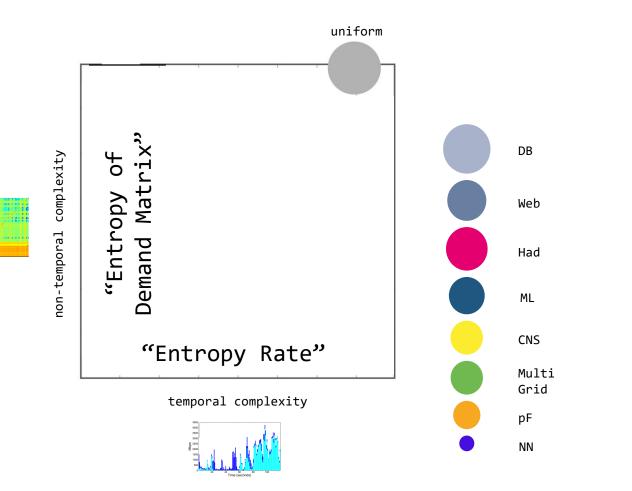
Check out trace complexity website!

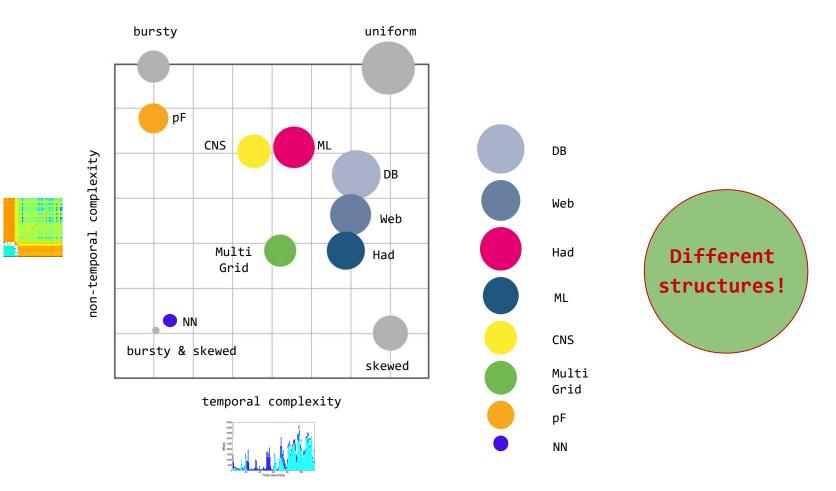


Griner et al., SIGMETRICS 2020

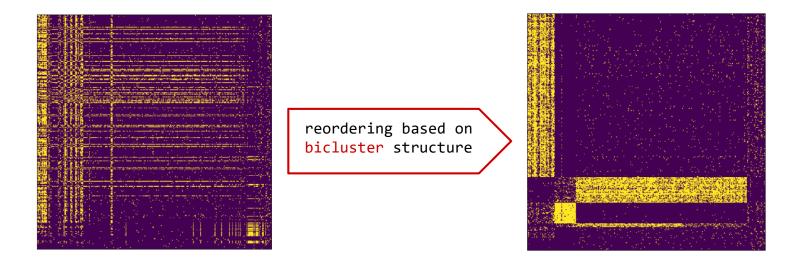


Griner et al., SIGMETRICS 2020





Traffic is also clustered (WWW'23): Small Stable Clusters



Opportunity: exploit with little reconfigurations!

Förster et al., Analyzing the Communication Clusters in Datacenters. WWW 2023

Beyond Worst-Case

Many approaches, also for online algorithms!



- Restricting inputs: random arrival order, locality of reference, access graph, smoothed analysis, independent sampling, diffused adversaries, distributional analysis, ...
- Deviating from competitive analysis: resource augmentation, loose competitiveness, and competitiveness with high probability
- ---> Advice: Next slide ©

Advice Models

Advice from an oracle: *side-loaded information* about the future helps online algorithms to make better decisions.

Model 1: perfect advice

- ---> Assumes a powerful, fully *trustworthy oracle*
- \dashrightarrow Provides algorithm with any information about the future
- → Question how many bits of the advice an online algorithm needs to achieve a certain competitive ratio c

Model 2: predictions (untrusted advice)

- Introduced by Mitzenmacher and Vassilvitskii: predictor may be faulty, and the competitive ratio depends on its error
- For small error, algorithm should perform close to the offline optimum (*consistency*), for large error, not worse than non-augmented online algorithms (*robustness*)

Model 3: lookahead

---> Related to local algorithms of distributed computing*

Limitations



All these models require new algorithmic
features (e.g., a designated advice tape)

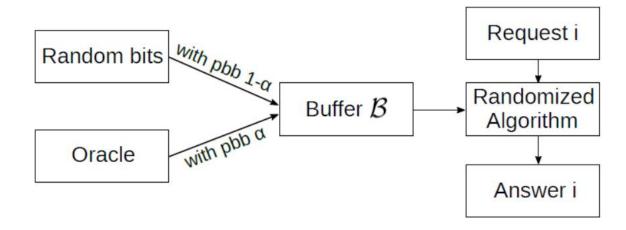
Hence not applicable to existing online algorithms

→ Bounds also depend on choice of error function

Infused Advice Model

An alternative

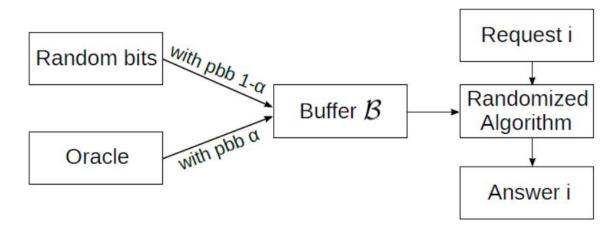
Idea: for randomized online algorithms, we may feed advice "non-intrusively" via the random bits tape



Infused Advice Model

An alternative

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Works also for existing online algorithms

First Results (ESA'23)

- Applies to: paging, uniform metrical task systems,
 online set cover, etc.
- \rightarrow New *upper bounds* which improve when the infusion parameter α increases
- ---> Often tight lower bounds (assuming algorithm cannot access buffer of previous rounds)

Problem	upper bound	lower bound
Caching	$\min\{2H_k, \frac{2}{\alpha}\}$	$\min\{H_k, \frac{1}{\alpha}\}$
Uniform MTS	$\min\{2H_n, \frac{2}{\alpha}+2\}$	$\min\{H_{n-1},\frac{1}{\alpha}\}$
Online set cover	$O(\min\{\log d \log n, \frac{\log n}{\alpha}\})$	$\min\{\frac{1}{2}\log d, \frac{1}{2\alpha}\}$

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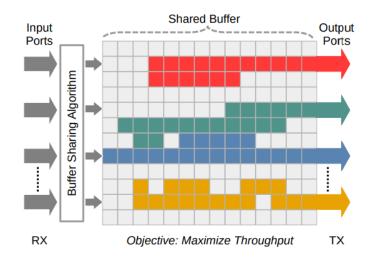
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Application domains for augmentation?

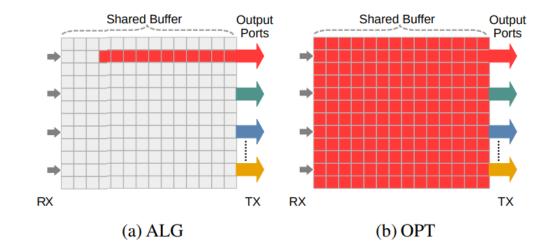
Example: Programmable Network Switches

- → Classic buffer sharing algorithm: which packets to accept/drop?
- Typically no pushout: once packet admitted to buffer, cannot be dropped later
- Future arrivals unknown:
 how to maximize throughput?
- Programmable switches enable new applications: running ML models in the data plane



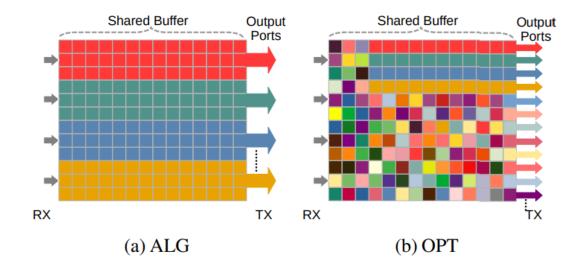
Example: Programmable Network Switches

- → Challenge: drop tail algorithms may drop red packets in order to keep buffer for other ports (and transmit in parallel)
- ---> Not competitive: if not differently colored packets arrive



Example: Programmable Network Switches

Alternatively: if drop tail algorithm *absorbs burst*, it may lead to reactive drops in future and low throughput



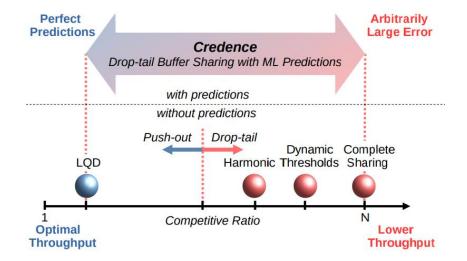
Example: Programmable Network Switches

- ---> Augment switches with predictions
- Simple random forest approach significantly improves competitive ratio
- ---> Depending on prediction error
- ---> Can be implemented on switch
 hardware...

Algorithm	Competitive Ratio
Complete Sharing [26]	N+1
Dynamic Thresholds [20, 26]	O(N)
Harmonic [33]	$\ln(N)+2$
LQD (push-out) [9, 26]	1.707
LateQD (clairvoyant) [14]	1
CREDENCE	$\min(1.707\eta,N)$

Example: Programmable Network Switches

- ---> Further reading: NSDI'24
- Wideo: https://www.youtube.com/watch?v=sAPe78RFsz0





- ---> Augmenting online algorithms with learning can benefit competitive ratio *in theory but also in practice* as input is skewed and predictable
- ---> Network equipment supports such augmentations to some extent
- ---> **Infused advice** approach allows to augment existing algorithms
- ---> Interesting use cases in networked systems, much to explore

Data Available



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To reference this website, please use:	bibtex Source Information High Performance	Туре	Lines	Size	Download		

https://trace-collection.net/ Trace collection website

Selected References

On the Complexity of Traffic Traces and Implications

Chen Avin, Manya Ghobadi, Chen Griner, and Stefan Schmid. ACM SIGMETRICS, Boston, Massachusetts, USA, June 2020.

Survey of Reconfigurable Data Center Networks: Enablers, Algorithms, Complexity Klaus-Tycho Foerster and Stefan Schmid.

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Analyzing the Communication Clusters in Datacenters

Klaus-Tycho Foerster, Thibault Marette, Stefan Neumann, Claudia Plant, Ylli Sadikaj, Stefan Schmid, and Yllka Velaj. The Web Conference (WWW), Austin, Texas, USA, April 2023.

Tight Bounds for Online Graph Partitioning

Monika Henzinger, Stefan Neumann, Harald Räcke, and Stefan Schmid.

ACM-SIAM Symposium on Discrete Algorithms (SODA), Alexandria, Virginia, USA, January 2021.

Efficient Distributed Workload (Re-)Embedding

Monika Henzinger, Stefan Neumann, and Stefan Schmid.

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Duo: A High-Throughput Reconfigurable Datacenter Network Using Local Routing and Control

Johannes Zerwas, Csaba Györgyi, Andreas Blenk, Stefan Schmid, and Chen Avin.

ACM SIGMETRICS and ACM Performance Evaluation Review (PER), Orlando, Florida, USA, June 2023.

Online Algorithms with Randomly Infused Advice

Yuval Emek, Yuval Gil, Maciej Pacut, and Stefan Schmid.

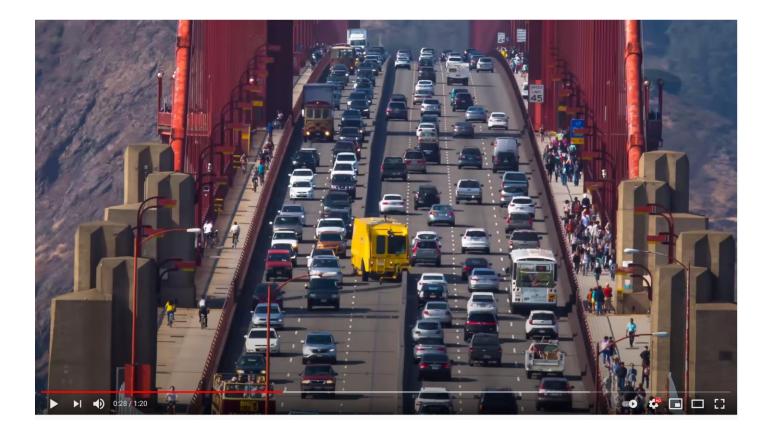
European Symposium on Algorithms (ESA), Amsterdam, Netherlands, September 2023.

Credence: Augmenting Datacenter Switch Buffer Sharing with ML Predictions

Vamsi Addanki, Maciej Pacut, and Stefan Schmid.

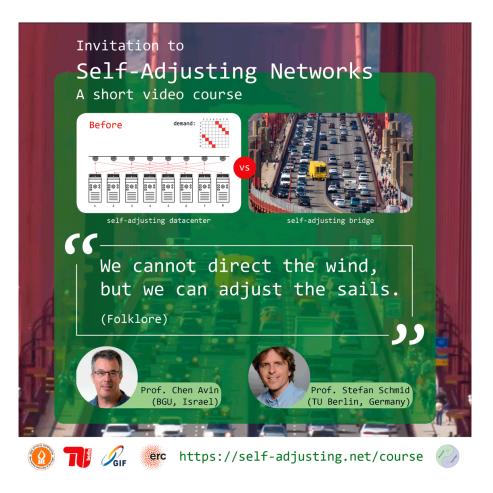
21st USENIX Symposium on Networked Systems Design and Implementation (NSDI), Santa Clara, California, USA, April 2024.

Thank you! Questions?



Golden Gate Zipper

Online Video Course



Bonus Material



Hogwarts Stair