

# Self-Adjusting Linear Networks

with Chen Avin\* and Stefan Schmid†

\*Ben Gurion University, †University of Vienna

Ingo van Duijn

Department of Computer Science  
Aalborg University  
Denmark



**AALBORG UNIVERSITY**  
DENMARK

# Baby Steps towards Self-Adjusting Linear Networks

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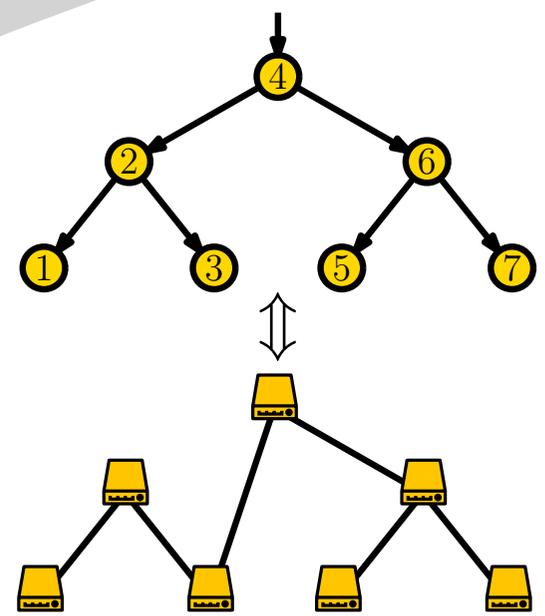
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DENMARK

# Overview

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Intro: correspondence between

- Self-Adjusting Networks
- Dynamic Data structures



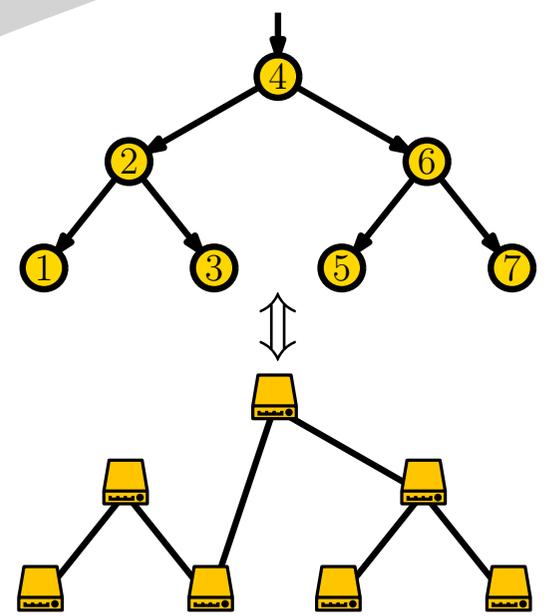
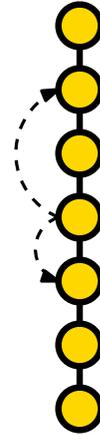
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Case study:

- List Access
- List Communication



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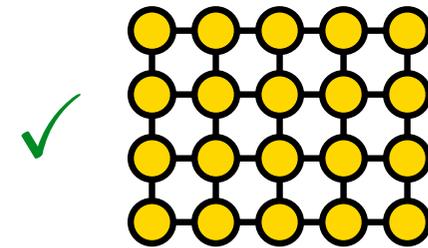
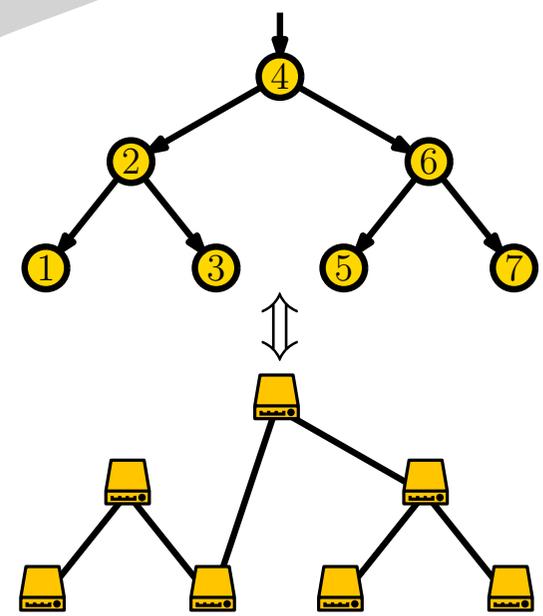
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Case study:

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Results:

- Grid Network Bounds



# Overview

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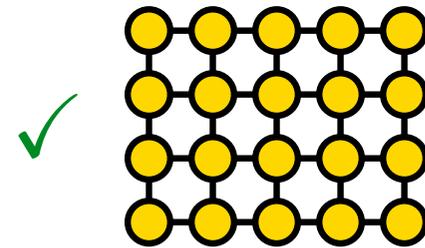
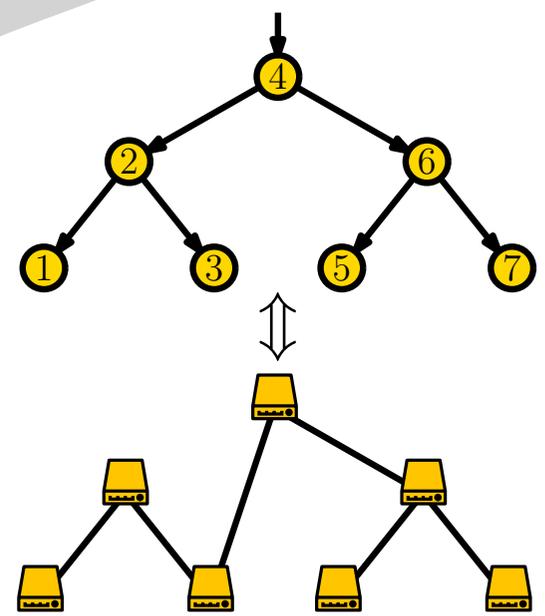
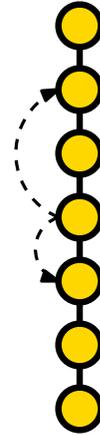
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Lower Bound Proof Sketch



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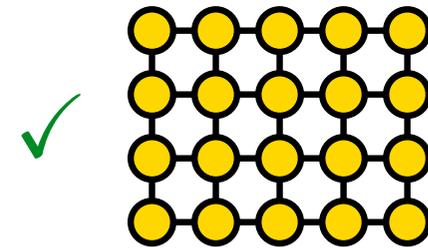
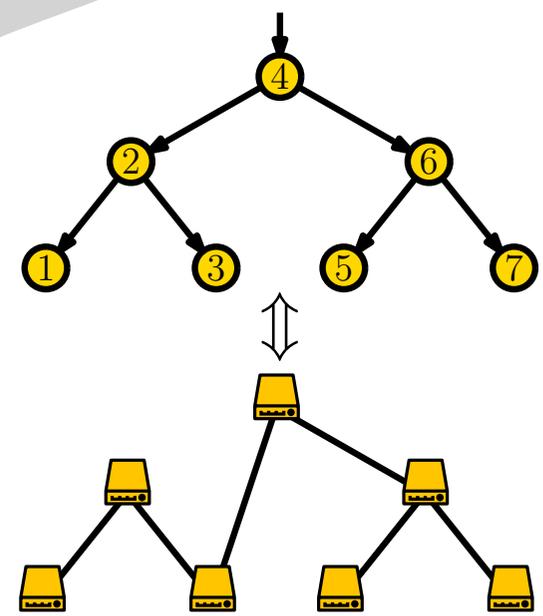
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Results:

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Lower Bound Proof Sketch

Future Work

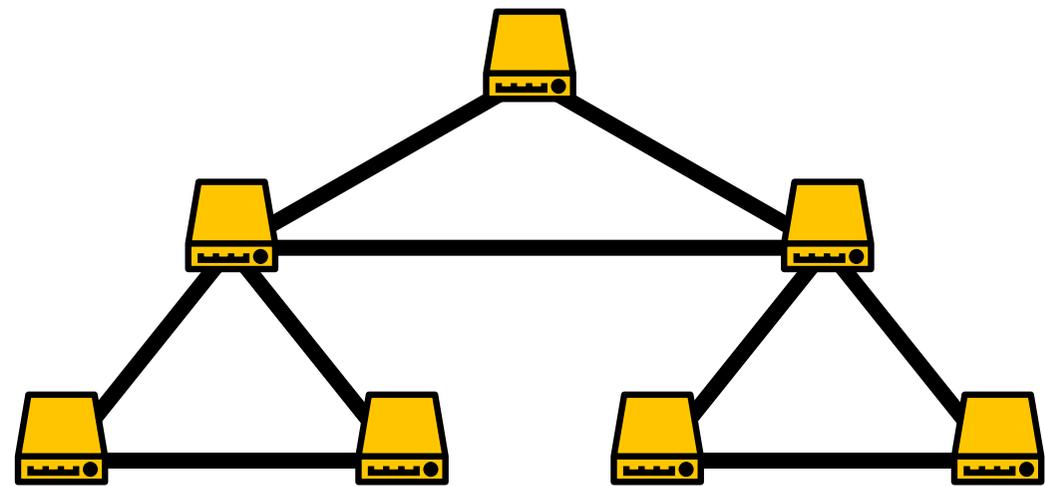


# Self-Adjusting Networks

# Self-Adjusting Networks

The Setting:

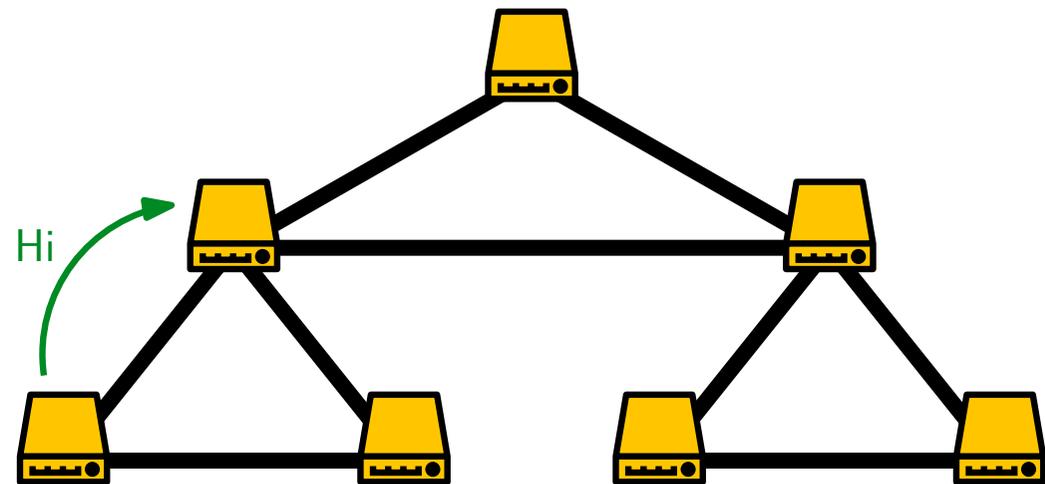
■ Network



# Self-Adjusting Networks

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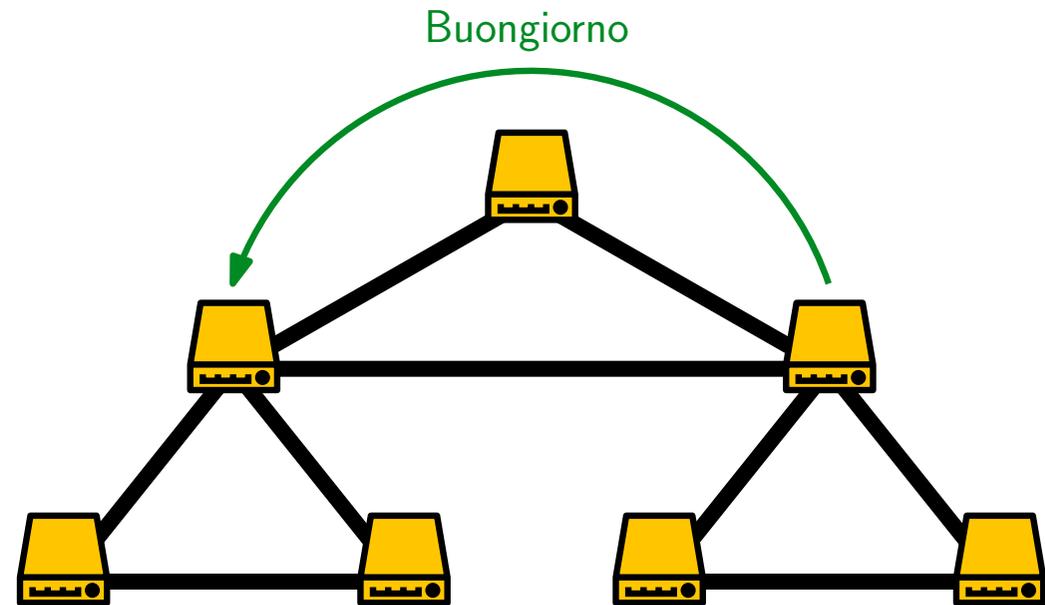
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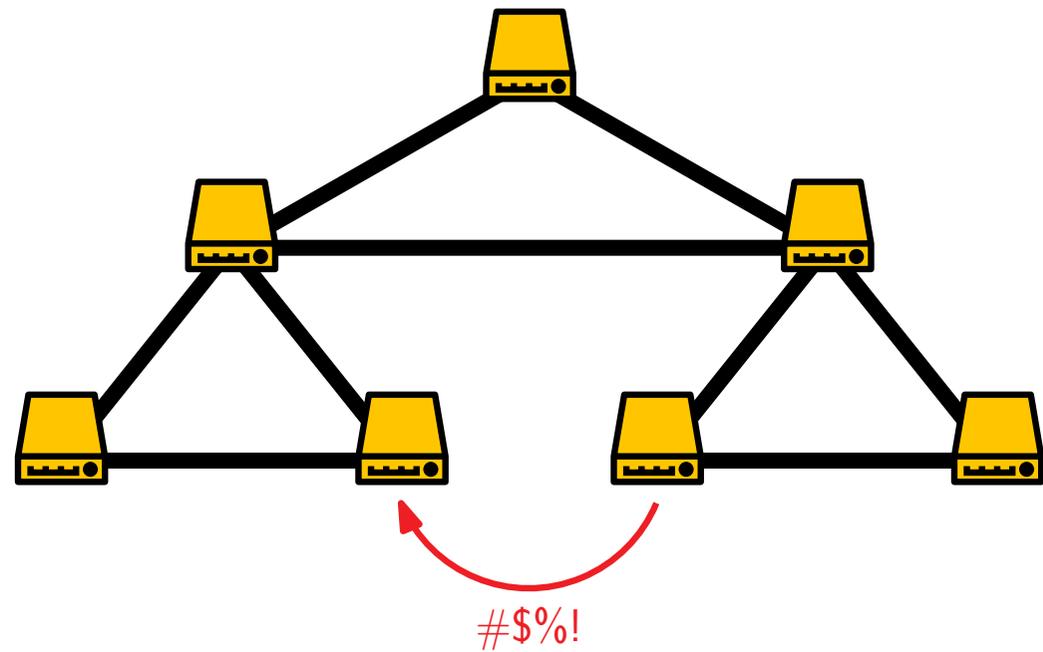
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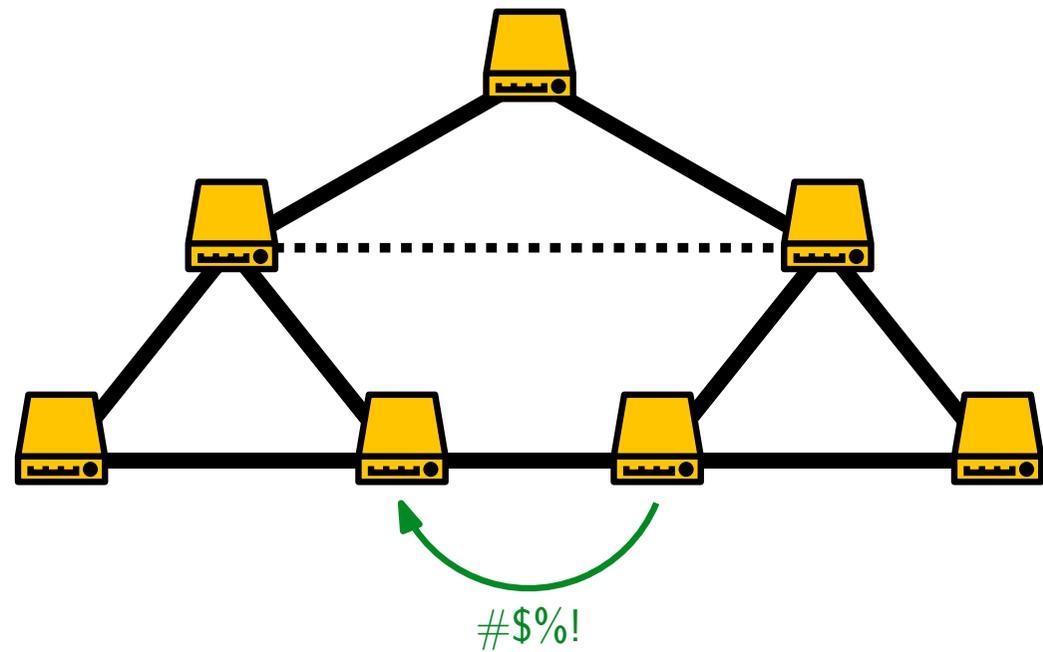
- Network
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# Self-Adjusting Networks

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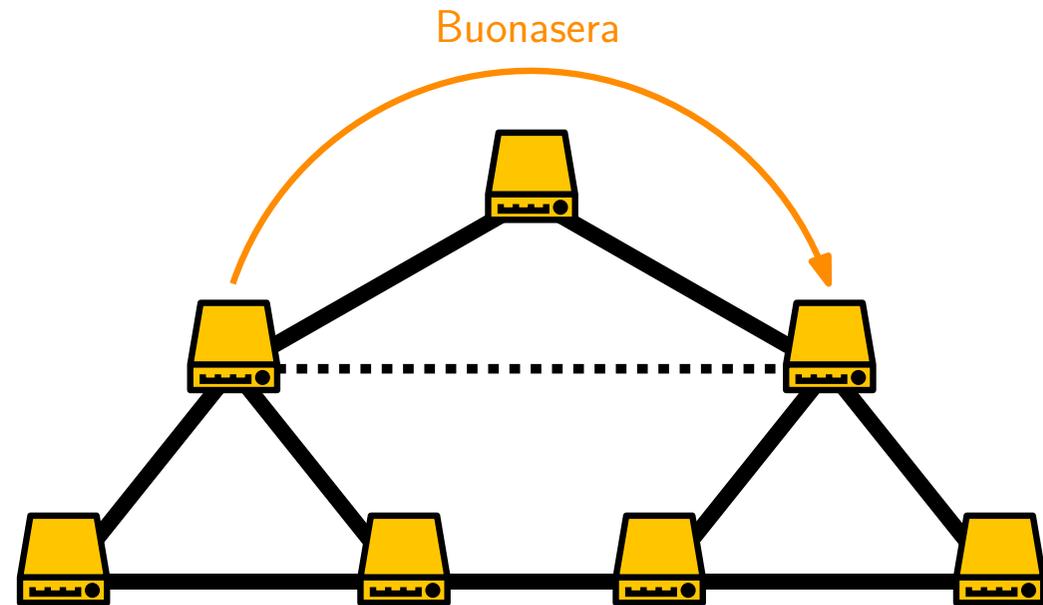
- Network
- Communication
- Adjustments



# Self-Adjusting Networks

The Setting:

- Network
- Communication
- Adjustments
- Online



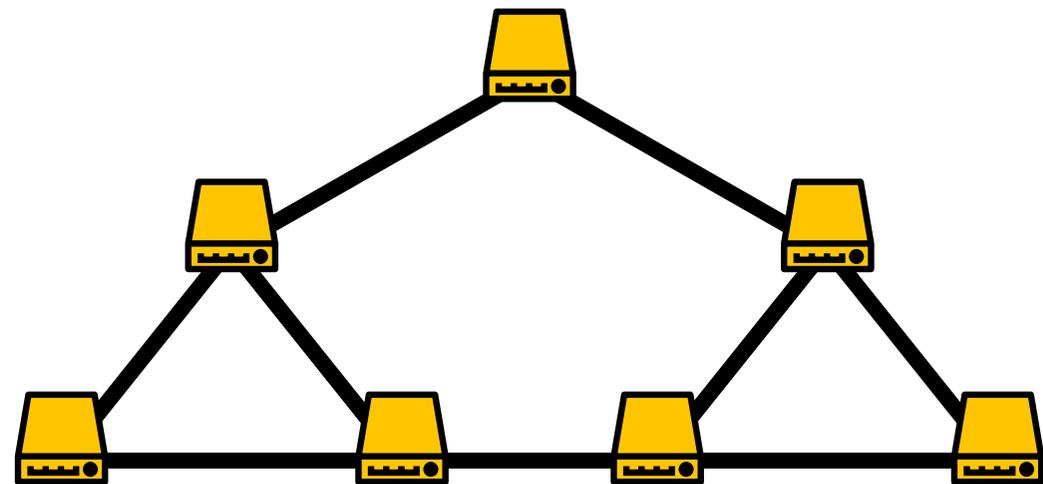
# Self-Adjusting Networks

The Setting:

- Network
  - Communication
  - Adjustments
  - Online
- } Cost!

The Goal:

- Minimise Cost



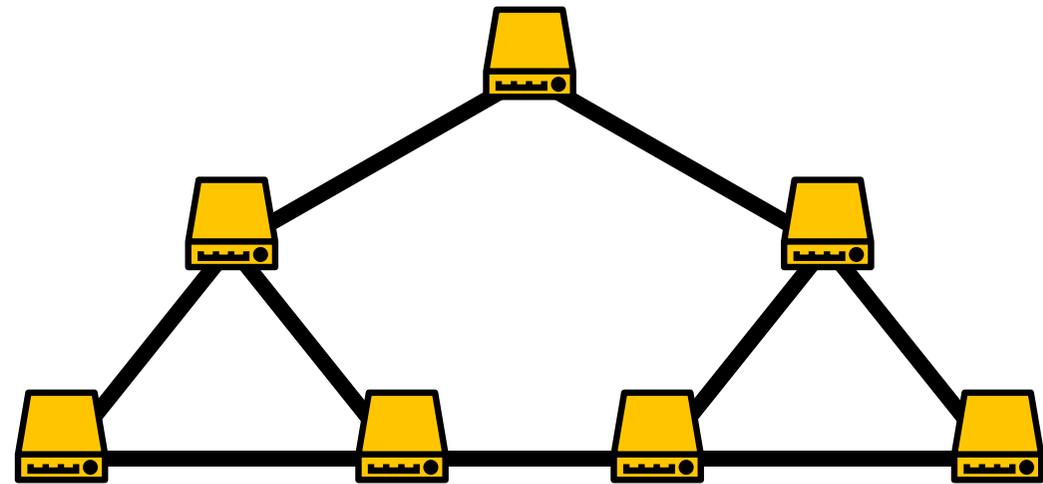
# Self-Adjusting Networks

The Setting:

- Network
  - Communication
  - Adjustments
  - Online
- } Cost!

The Goal:

- Minimise Cost
- (Distributed Algorithm)



# Self-Adjusting Networks

The Setting:

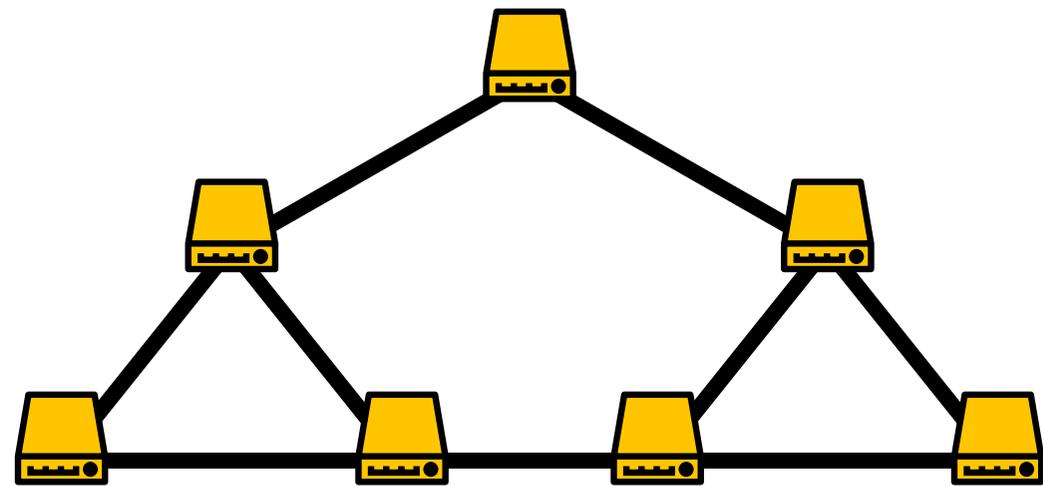
- Network
  - Communication
  - Adjustments
  - Online
- } Cost!

The Goal:

- Minimise Cost
- (Distributed Algorithm)

The Questions:

- Which Model?



# Self-Adjusting Networks

The Setting:

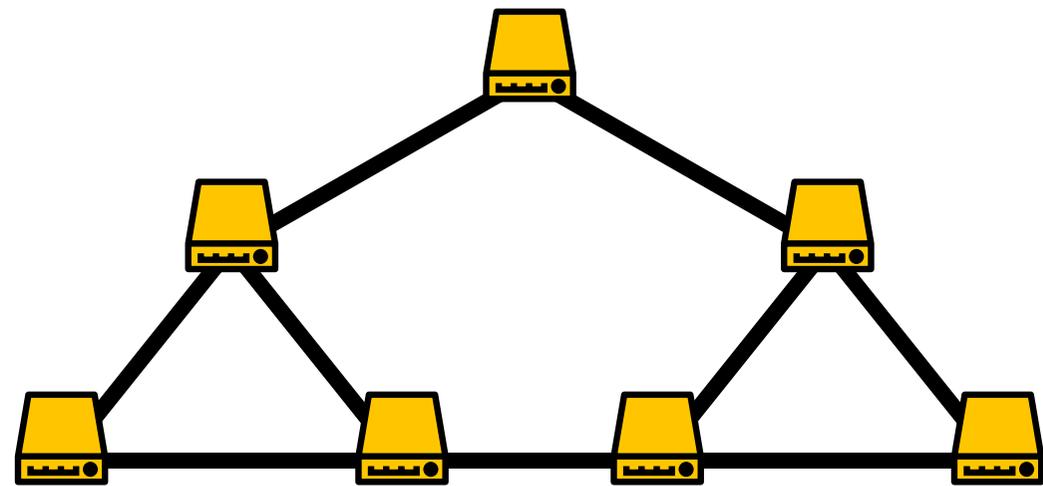
- Network
  - Communication
  - Adjustments
  - Online
- } Cost!

The Goal:

- Minimise Cost
- (Distributed Algorithm)

The Questions:

- Which Model?
- Formal Guarantees?

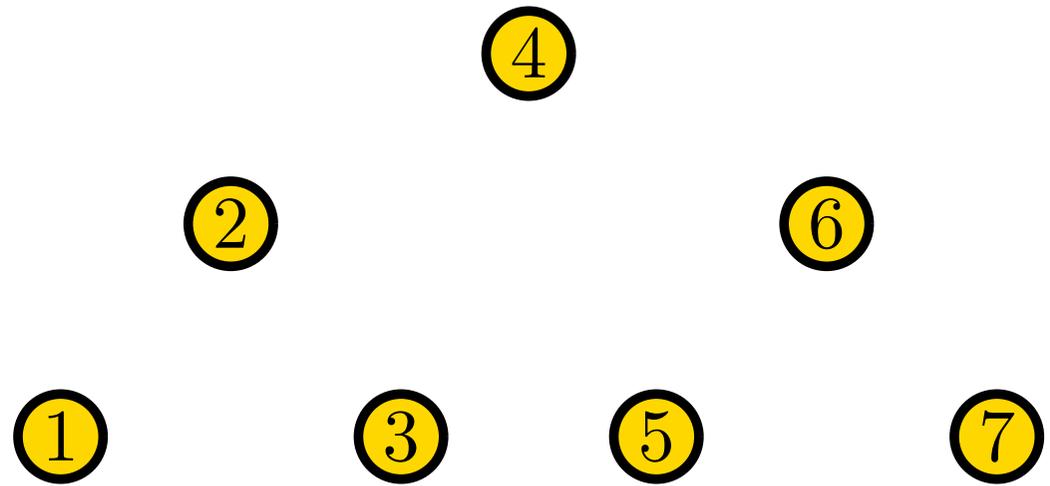


# Dynamic Dictionaries

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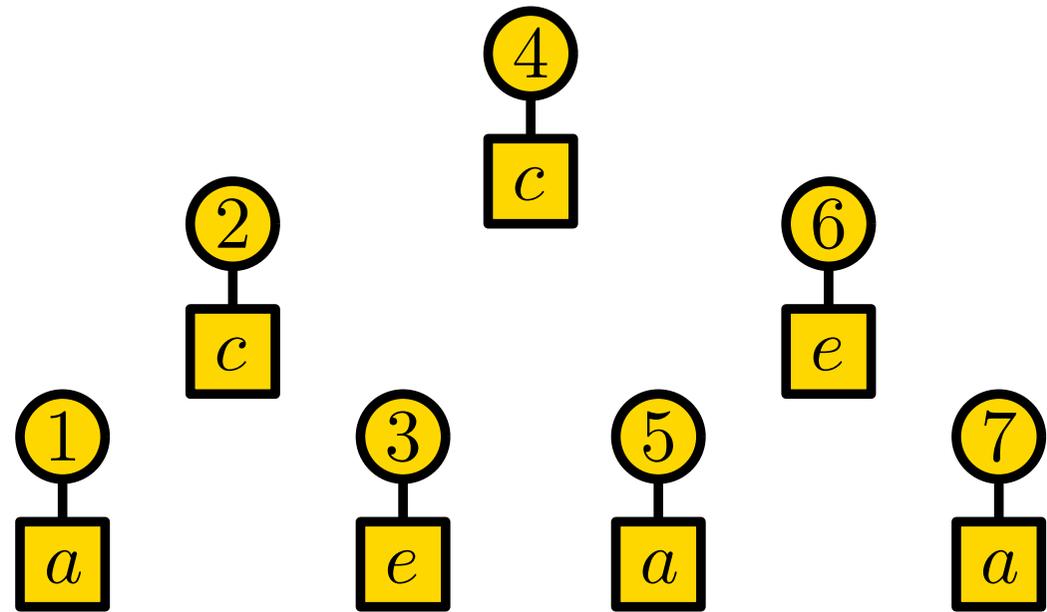
■ Keys



# Dynamic Dictionaries

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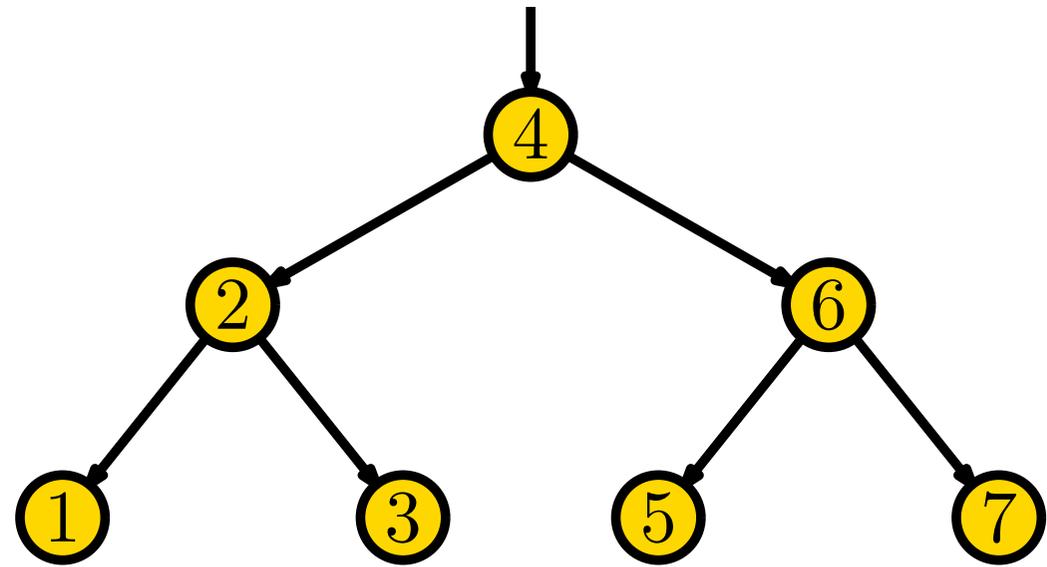
- Keys (with data)



# Dynamic Dictionaries

The Setting:

- Keys
- Pointer Structure

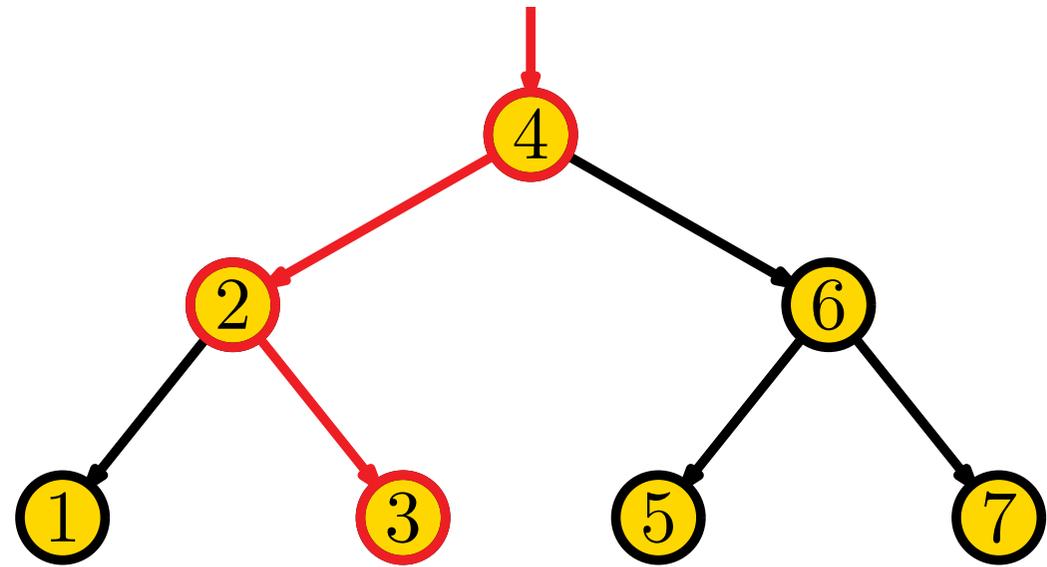


# Dynamic Dictionaries

The Setting:

- Keys
- Pointer Structure
- Key Access

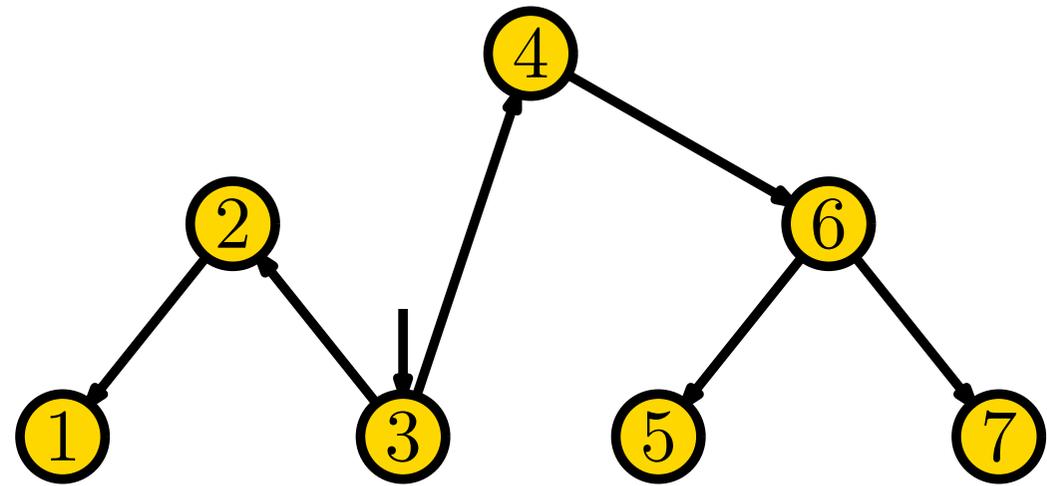
Access 3



# Dynamic Dictionaries

The Setting:

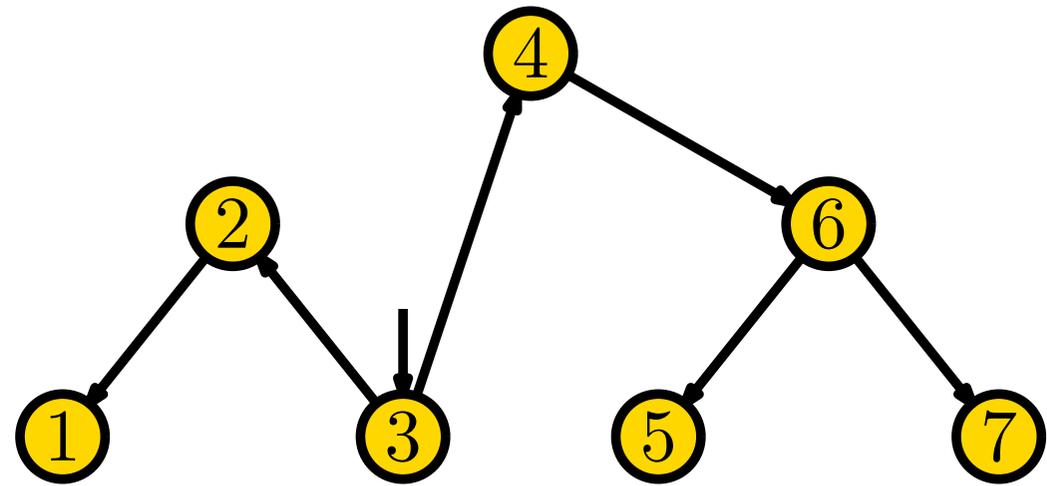
- Keys
- Pointer Structure
- Key Access
- Adjustments



# Dynamic Dictionaries

The Setting:

- Keys
- Pointer Structure
- Key Access
- Adjustments
- Online



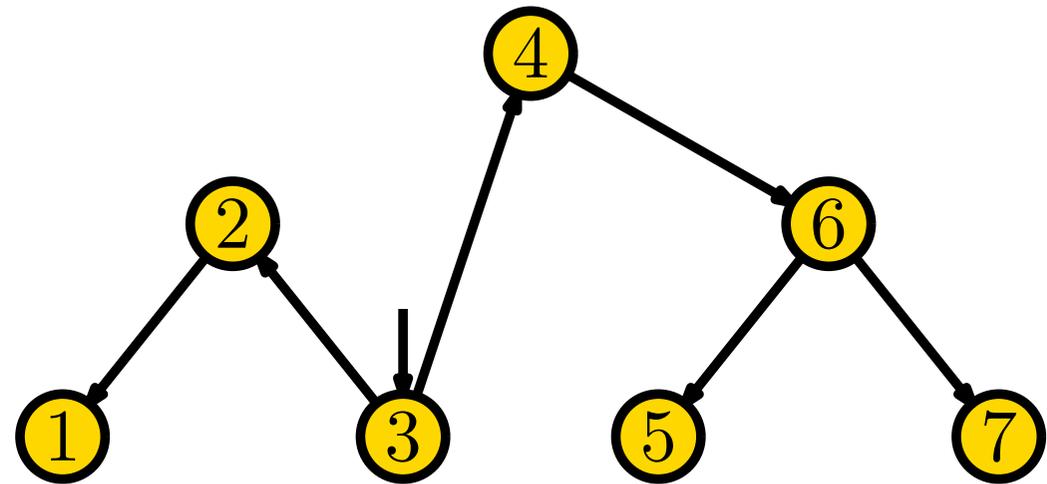
# Dynamic Dictionaries

The Setting:

- Keys
- Pointer Structure
- Key Access
- Adjustments
- Online

The Goal:

- Dynamic Optimality



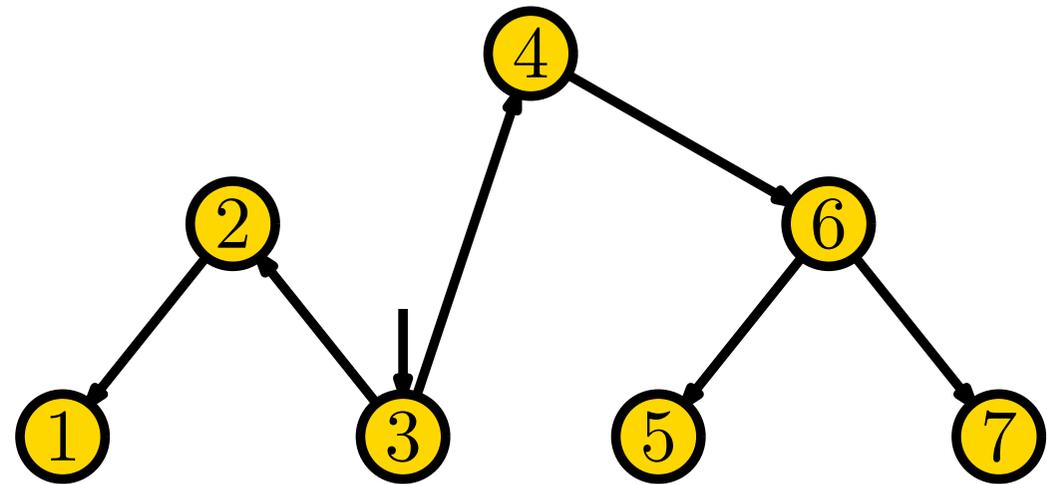
# Dynamic Dictionaries

The Setting:

- Keys
- Pointer Structure
- Key Access
- Adjustments
- Online

The Goal:

- Dynamic Optimality
- Minimise:  $\frac{\text{Best Online Cost}}{\text{Best Offline Cost}}$



# Dynamic Dictionaries

The Setting:

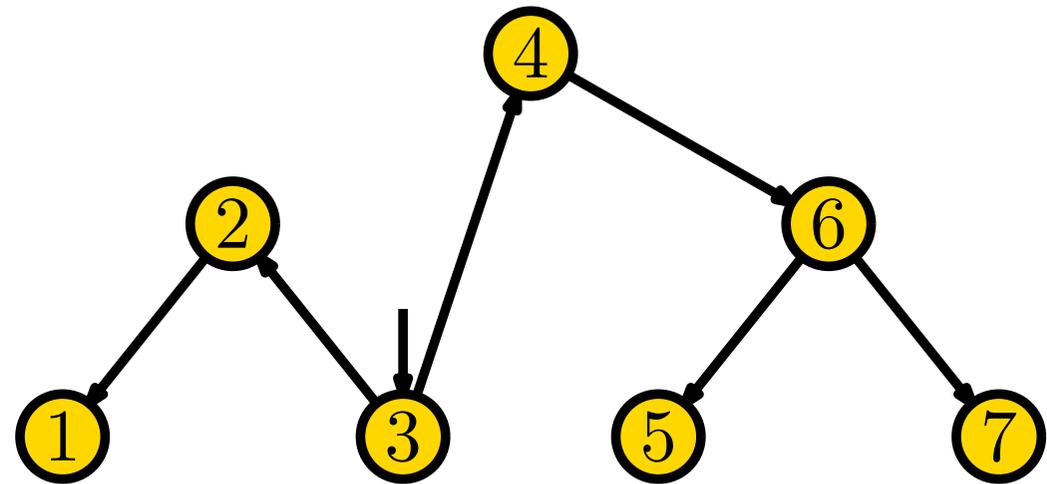
- Keys
- Pointer Structure
- Key Access
- Adjustments
- Online

The Goal:

- Dynamic Optimality
- Minimise:  $\frac{\text{Best Online Cost}}{\text{Best Offline Cost}}$

Towards Networks:

- Pairwise Access



# Dynamic Dictionaries

The Setting:

- Keys
- Pointer Structure
- Key Access
- Adjustments
- Online

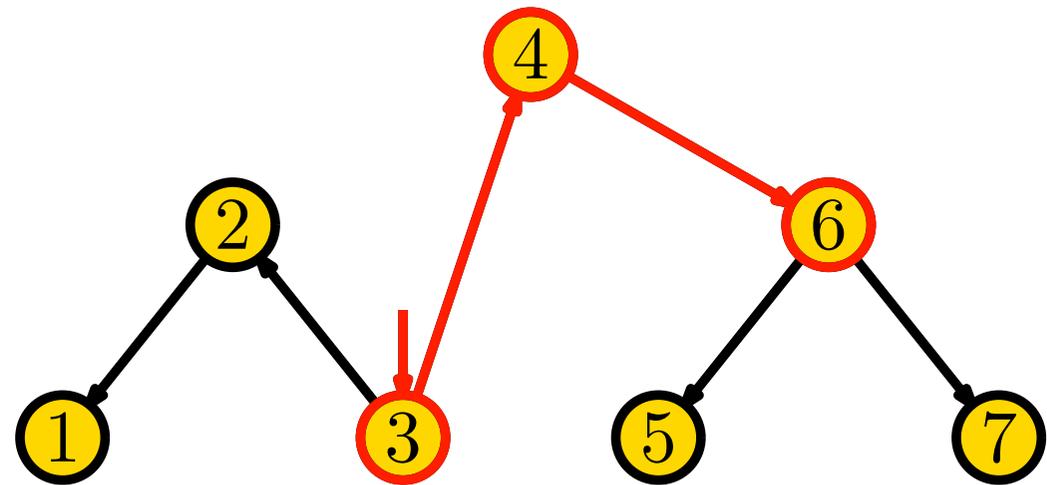
The Goal:

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- Minimise:  $\frac{\text{Best Online Cost}}{\text{Best Offline Cost}}$

Towards Networks:

- Pairwise Access

Access 6



# Dynamic Dictionaries

The Setting:

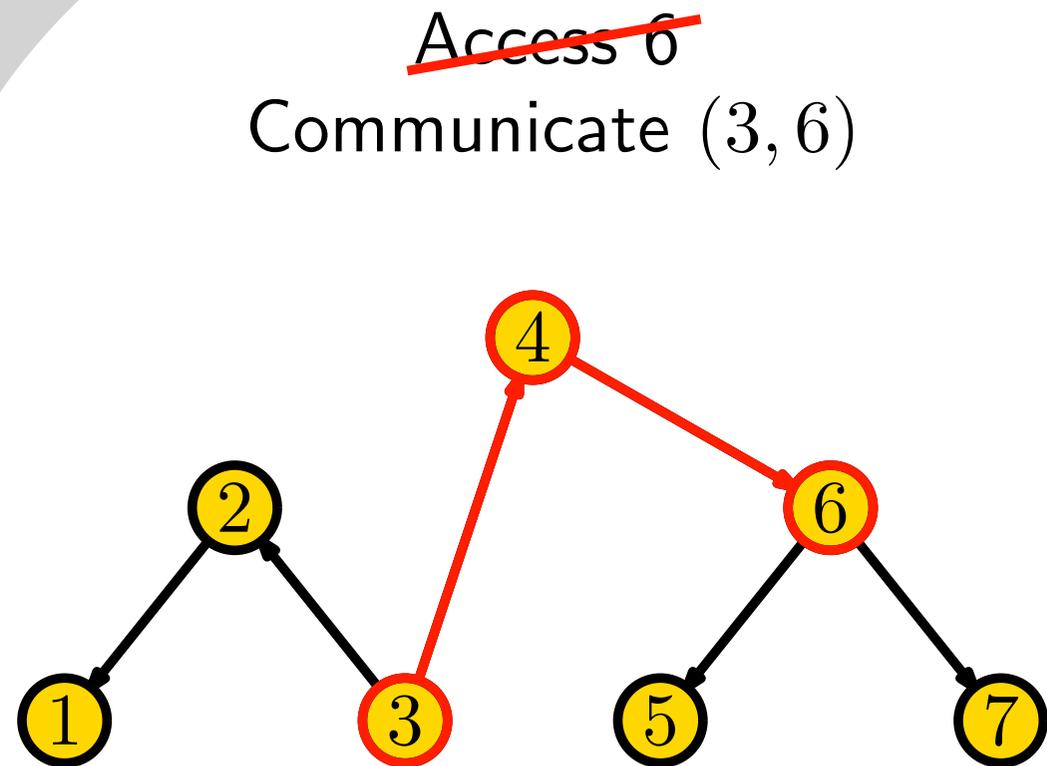
- Keys
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The Goal:

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Towards Networks:

- Pairwise Access



# Dynamic Dictionaries

## The Setting:

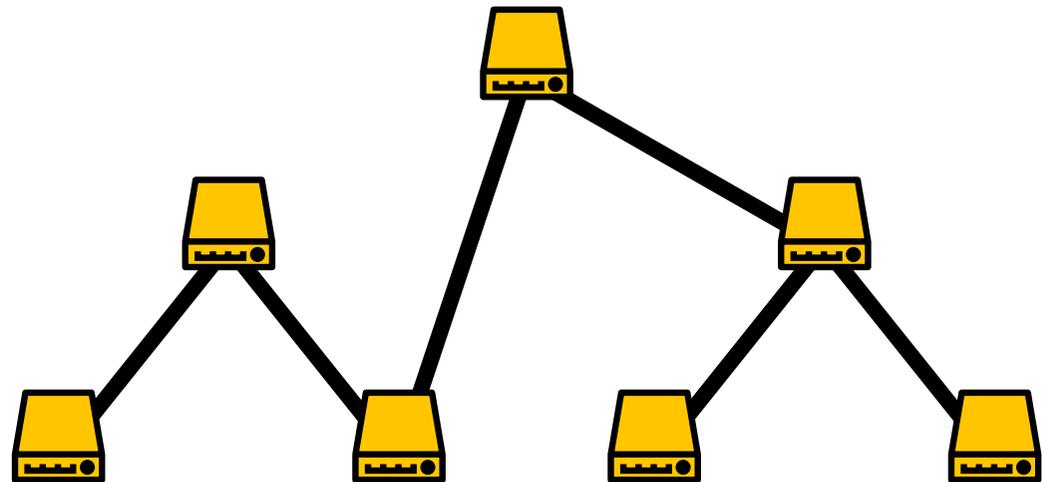
- Keys
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## The Goal:

- Dynamic Optimality
- Minimise:  $\frac{\text{Best Online Cost}}{\text{Best Offline Cost}}$

## Towards Networks:

- Pairwise Access
- Ignore Searching/Routing



# A Simple? Network

# A Simple? Network

The Dictionary:

■ Linked List



# A Simple? Network

The Dictionary:

- Linked List
- Access in Front

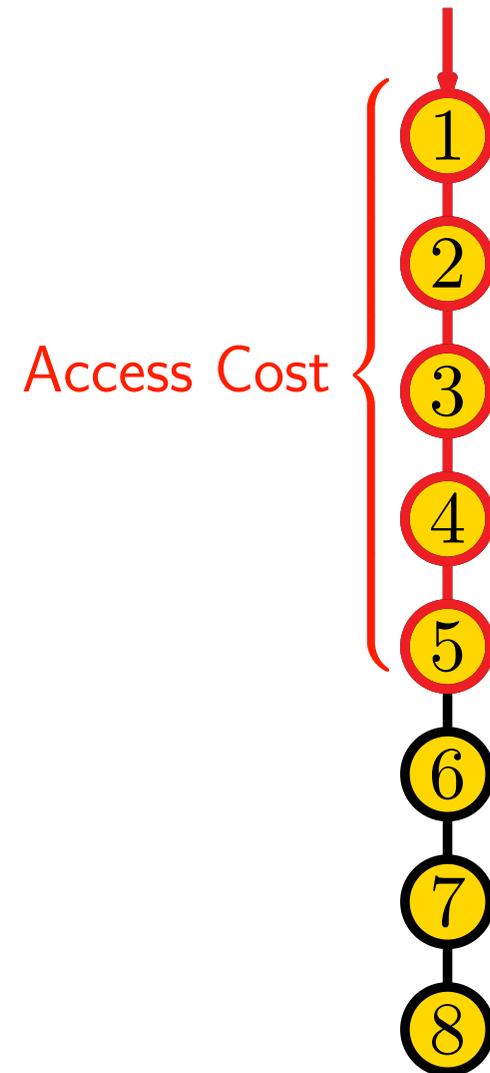


# A Simple? Network

The Dictionary:

- Linked List
- Access in Front

Access 5

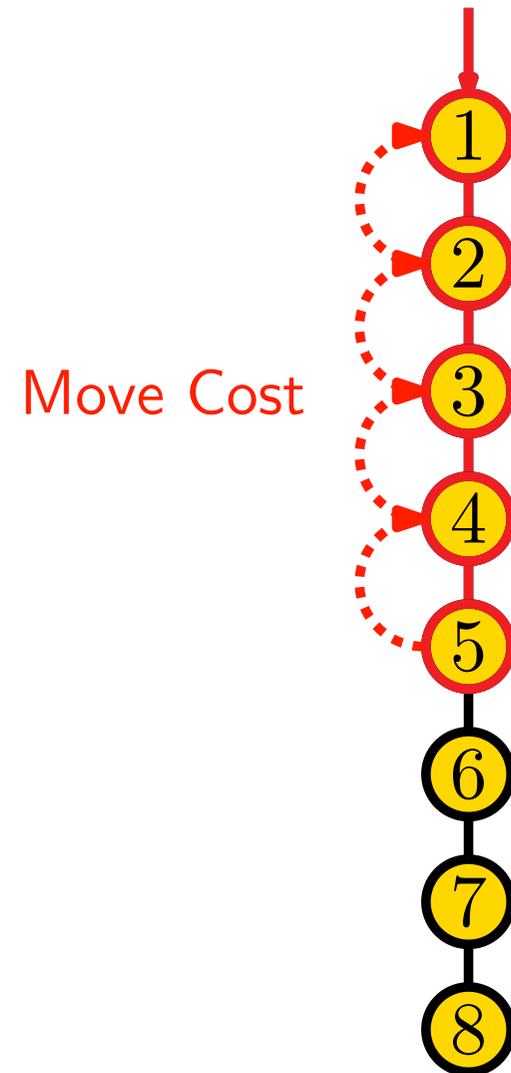


# A Simple? Network

The Dictionary:

- Linked List
- Access in Front
- Move-to-Front Algorithm

Access 5



# A Simple? Network

The Dictionary:

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- Access in Front
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Access 5



# A Simple? Network

The Dictionary:

- Linked List
- Access in Front
- Move-to-Front Algorithm  
→ Dynamically Optimal!



# A Simple? Network

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The Network:

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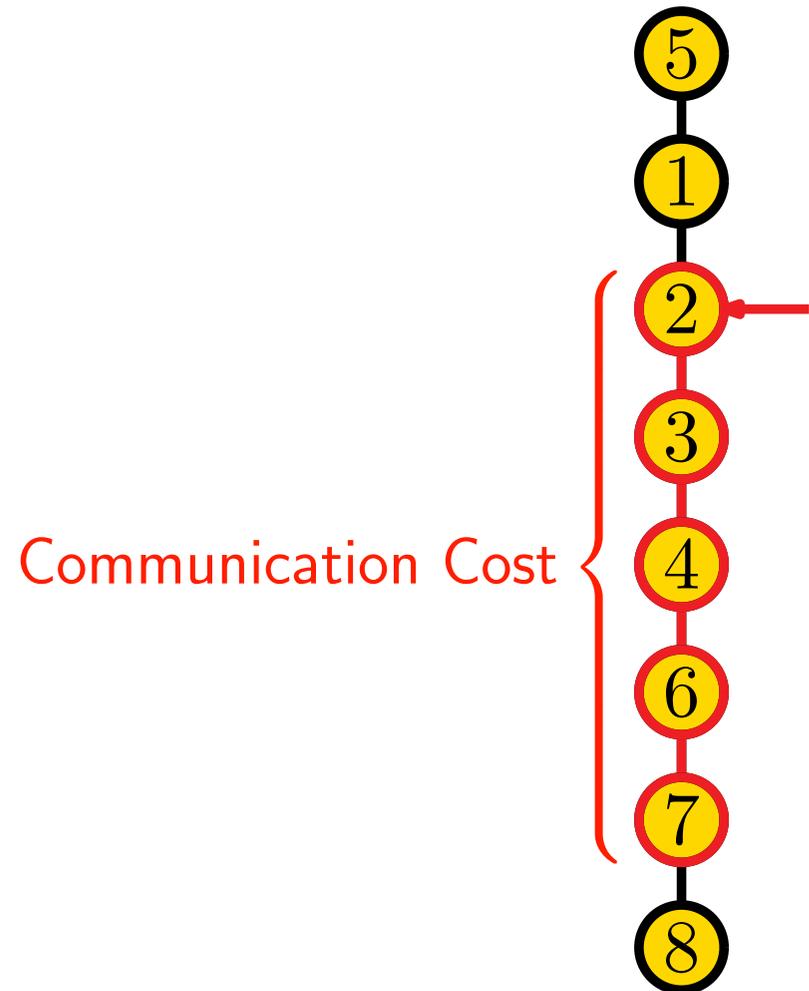
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The Network:

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Communicate (2,7)



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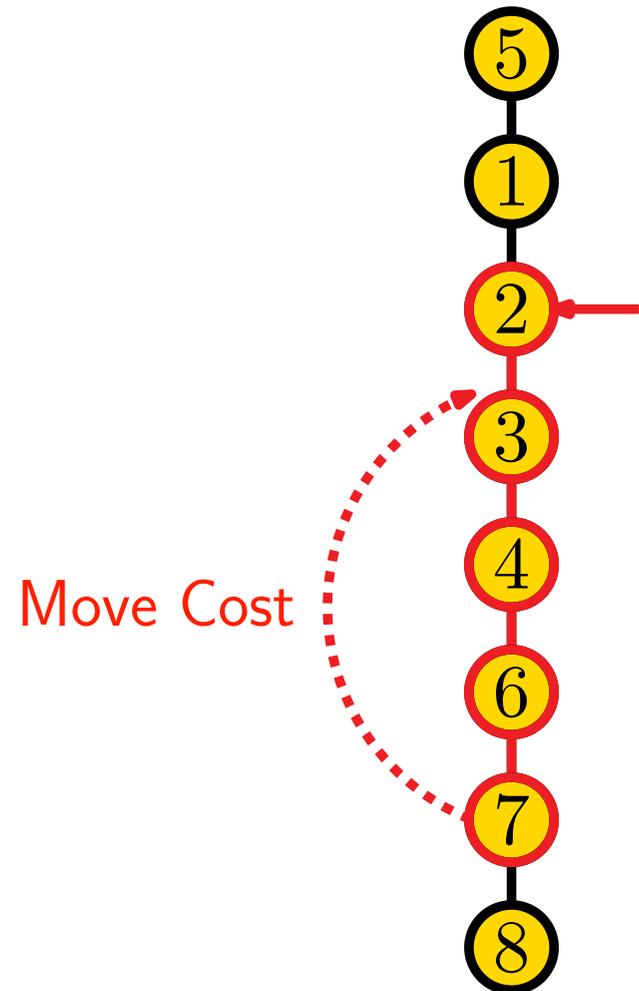
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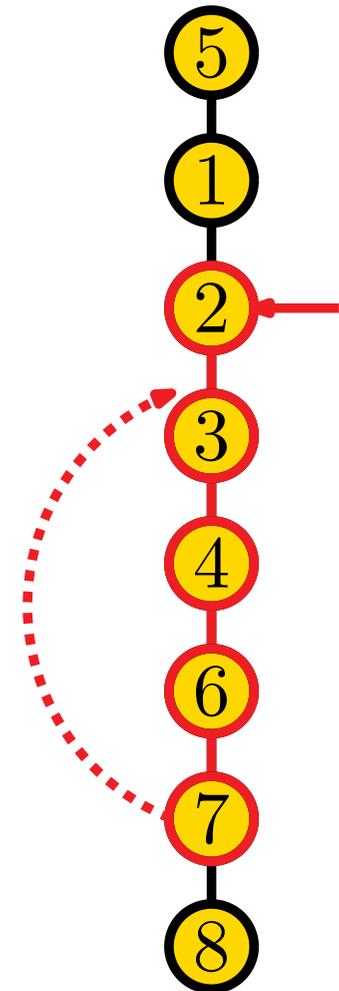
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The Network:

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- Move-to-Where?

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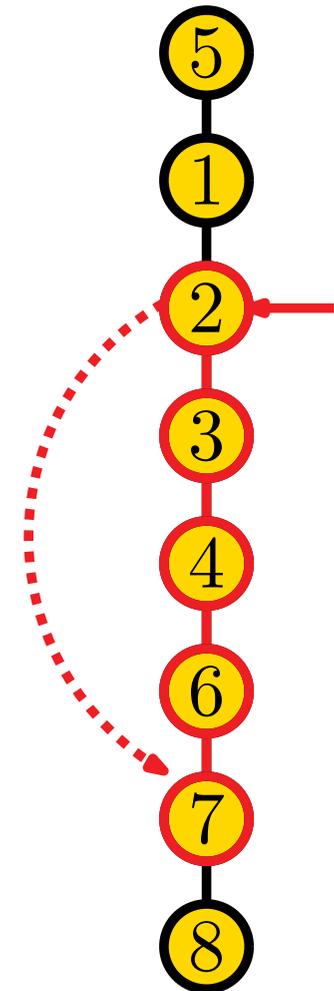
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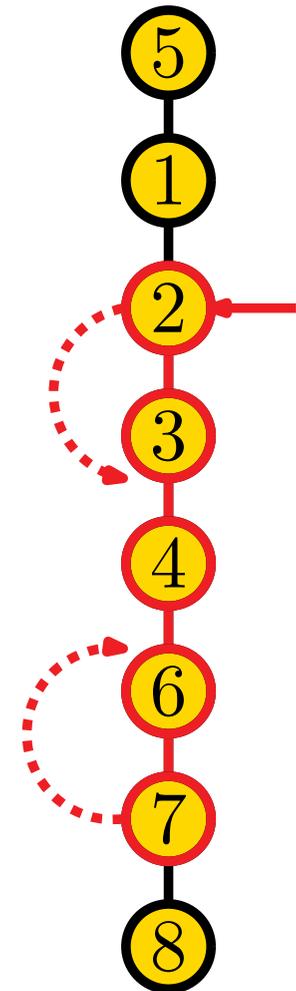
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The Network:

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Not So Simple:

- Co-locating can Backfire



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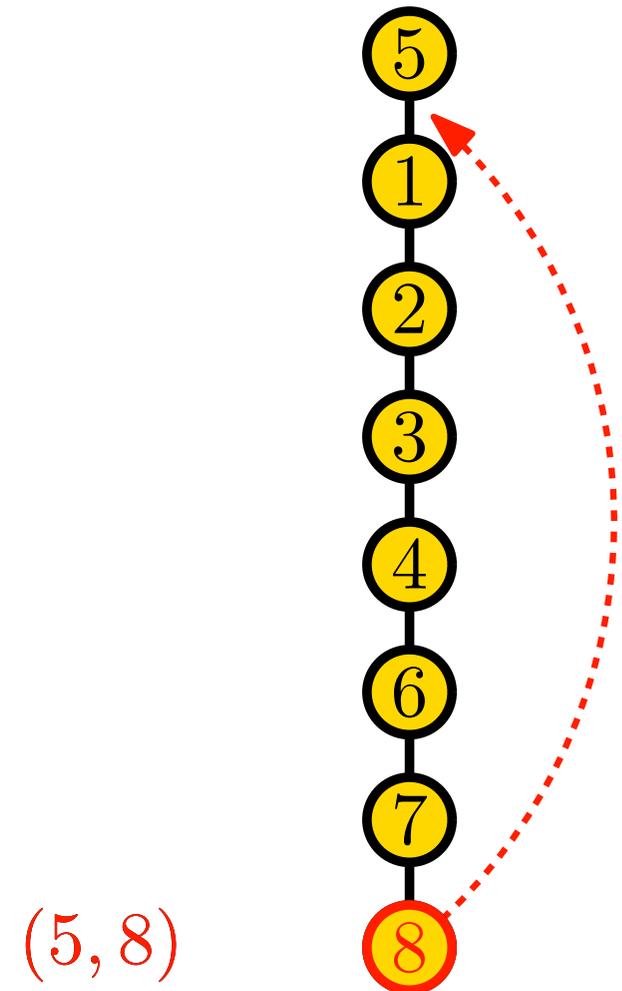
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Communication:



# A Simple? Network

The Dictionary:

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- Move-to-Front Algorithm  
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The Network:

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- Move-to-Where?

Not So Simple:

- Co-locating can Backfire

Communication:

(5, 8)



# A Simple? Network

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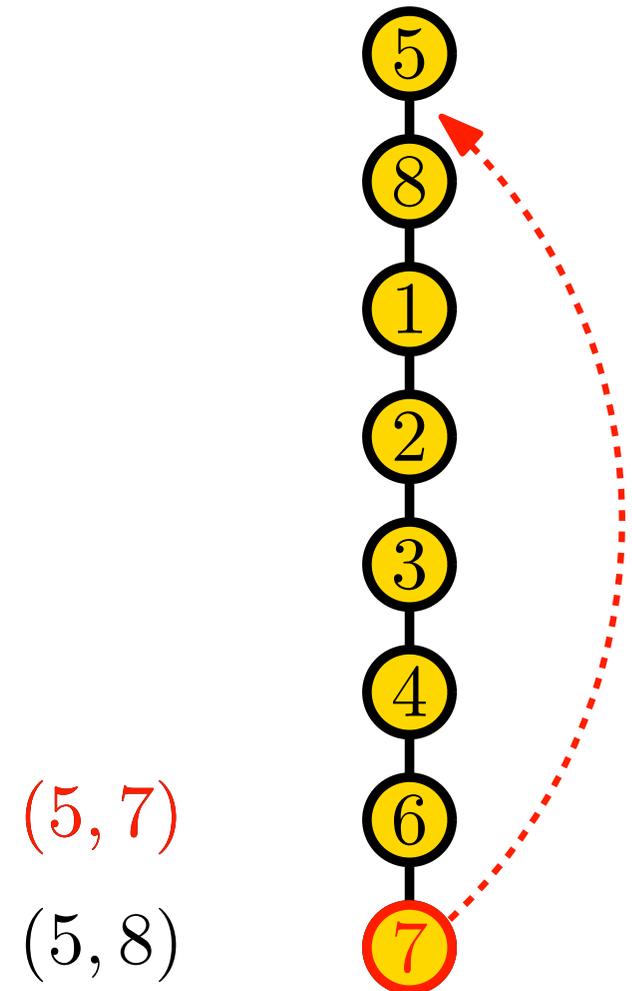
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Communication:



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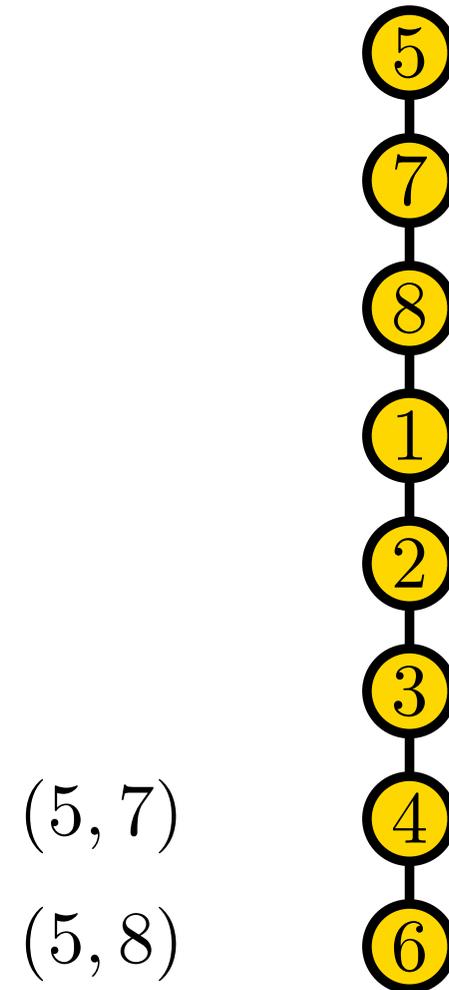
The Network:

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Not So Simple:

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Communication:



# A Simple? Network

The Dictionary:

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- Move-to-Front Algorithm  
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The Network:

- Pairwise Access
- Move-to-Where?

Not So Simple:

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Communication:

(5, 6)

(5, 7)

(5, 8)



# A Simple? Network

The Dictionary:

- Linked List
- Access in Front
- Move-to-Front Algorithm  
→ Dynamically Optimal!

The Network:

- Pairwise Access
- Move-to-Where?

Not So Simple:

- Co-locating can Backfire

Communication:

(5, 4)

(5, 6)

(5, 7)

(5, 8)



# A Simple? Network

The Dictionary:

- Linked List
- Access in Front
- Move-to-Front Algorithm  
→ Dynamically Optimal!

The Network:

- Pairwise Access
- Move-to-Where?

Not So Simple:

- Co-locating can Backfire

Communication:

(5, 3)

(5, 4)

(5, 6)

(5, 7)

(5, 8)



# A Simple? Network

The Dictionary:

- Linked List
- Access in Front
- Move-to-Front Algorithm  
→ Dynamically Optimal!

The Network:

- Pairwise Access
- Move-to-Where?

Not So Simple:

- Co-locating can Backfire

Communication:

(5, 2)

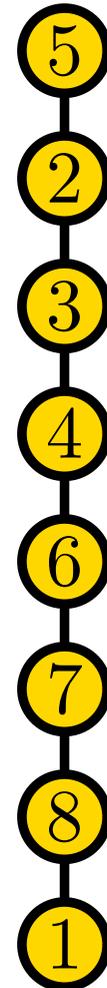
(5, 3)

(5, 4)

(5, 6)

(5, 7)

(5, 8)



# A Simple? Network

The Dictionary:

- Linked List
- Access in Front
- Move-to-Front Algorithm  
→ Dynamically Optimal!

The Network:

- Pairwise Access
- Move-to-Where?

Not So Simple:

- Co-locating can Backfire

Communication:

(5, 1)

(5, 2)

(5, 3)

(5, 4)

(5, 6)

(5, 7)

(5, 8)



# A Simple? Network

The Dictionary:

- Linked List
- Access in Front
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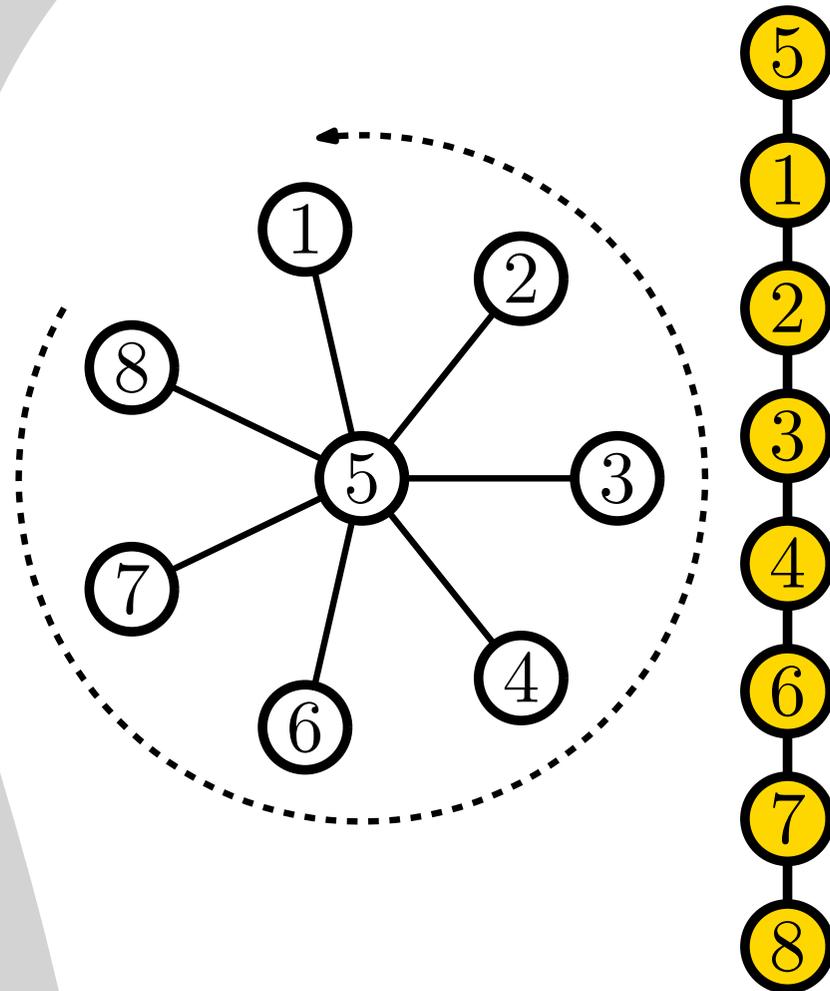
The Network:

- Pairwise Access
- Move-to-Where?

Not So Simple:

- Co-locating can Backfire
- Wheel can be Everywhere

Communication:



# A Simple? Network

The Dictionary:

- Linked List
- Access in Front
- Move-to-Front Algorithm  
→ Dynamically Optimal!

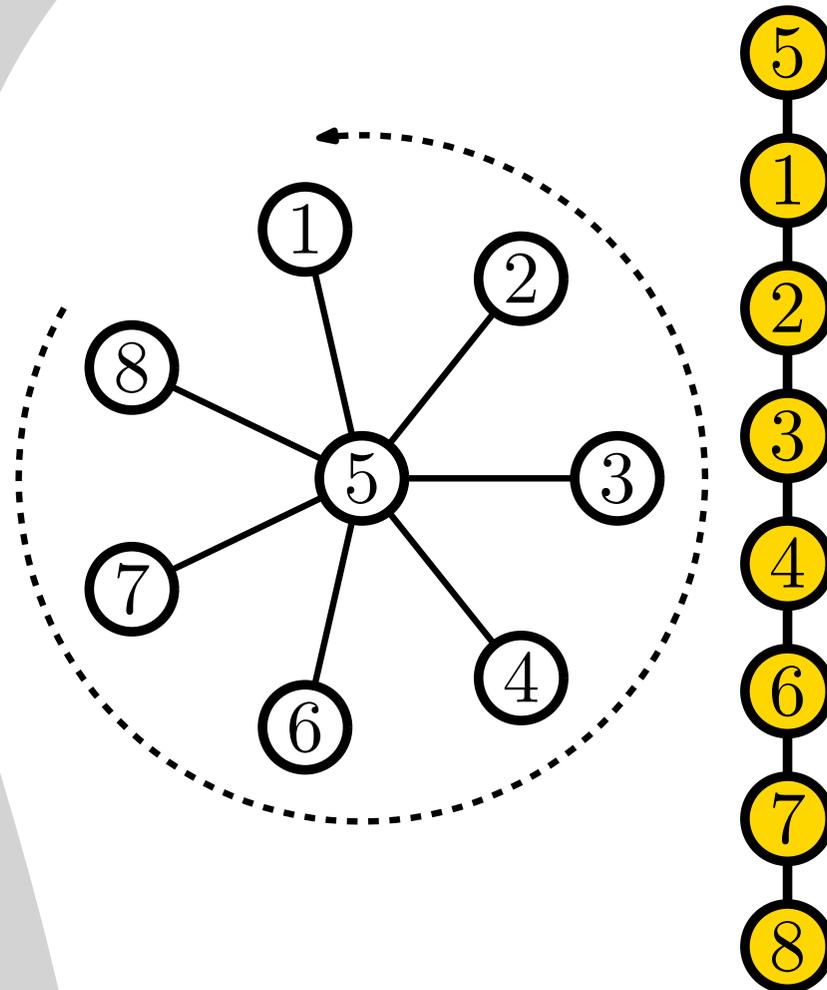
The Network:

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- Move-to-Where?

Not So Simple:

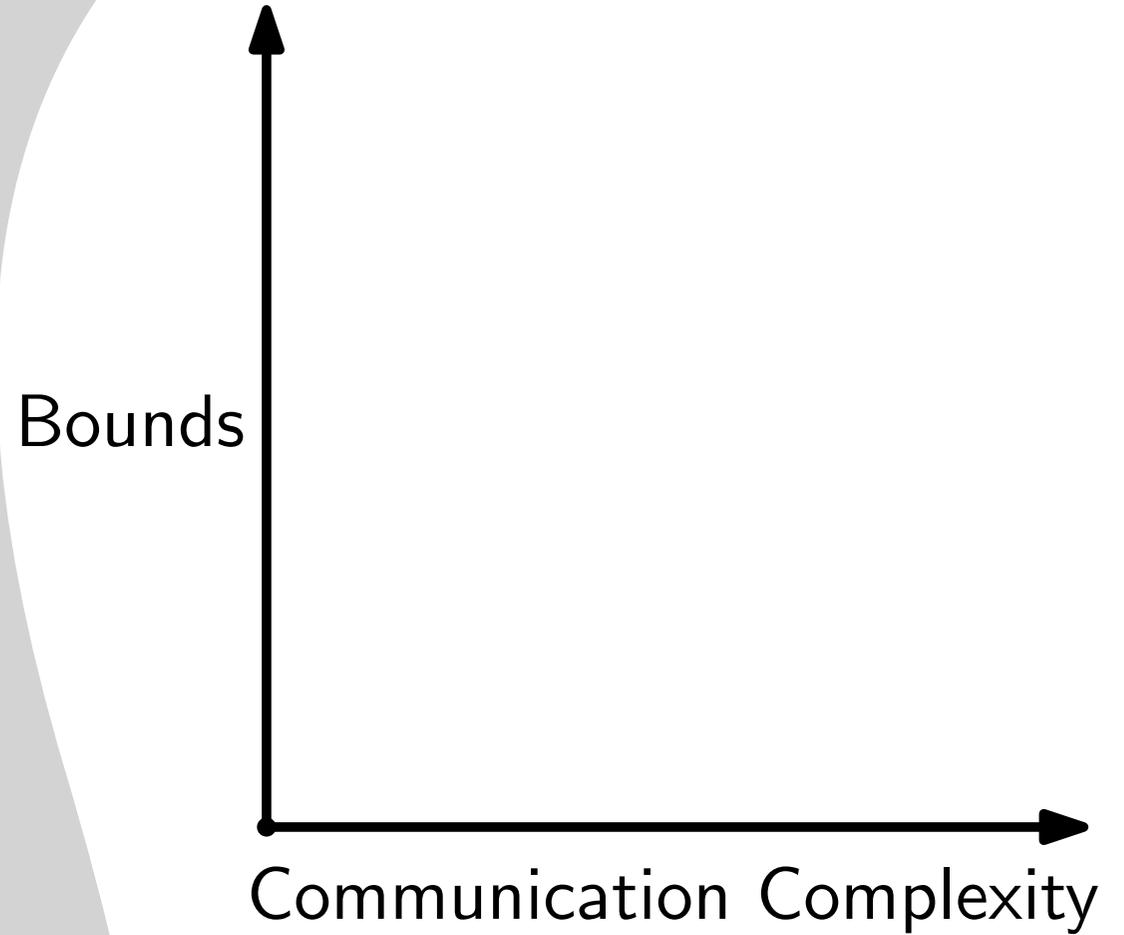
- Co-locating can Backfire
- Wheel can be Everywhere
- Same Conclusion by Olver et al.

Communication:



# Results

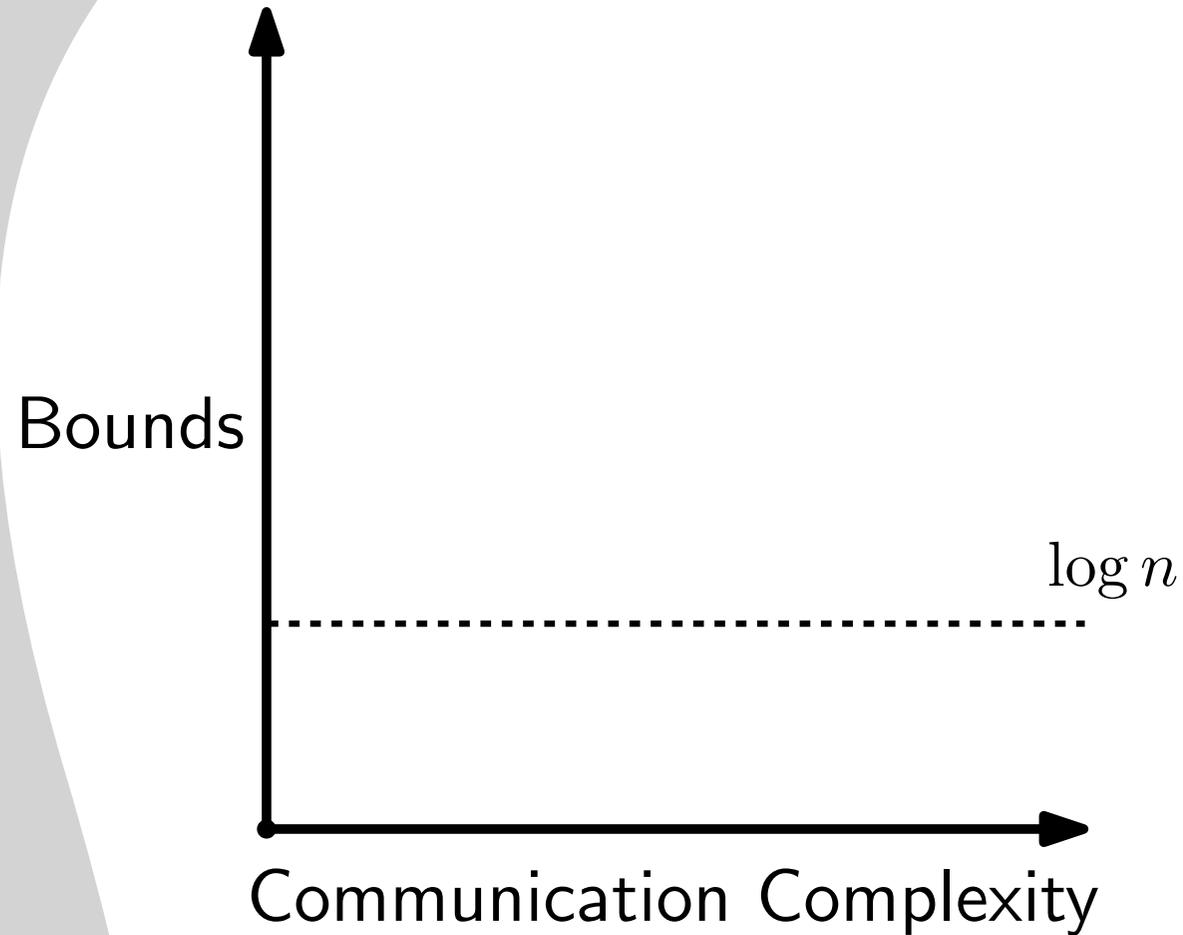
# Results



# Results

Grid Networks:

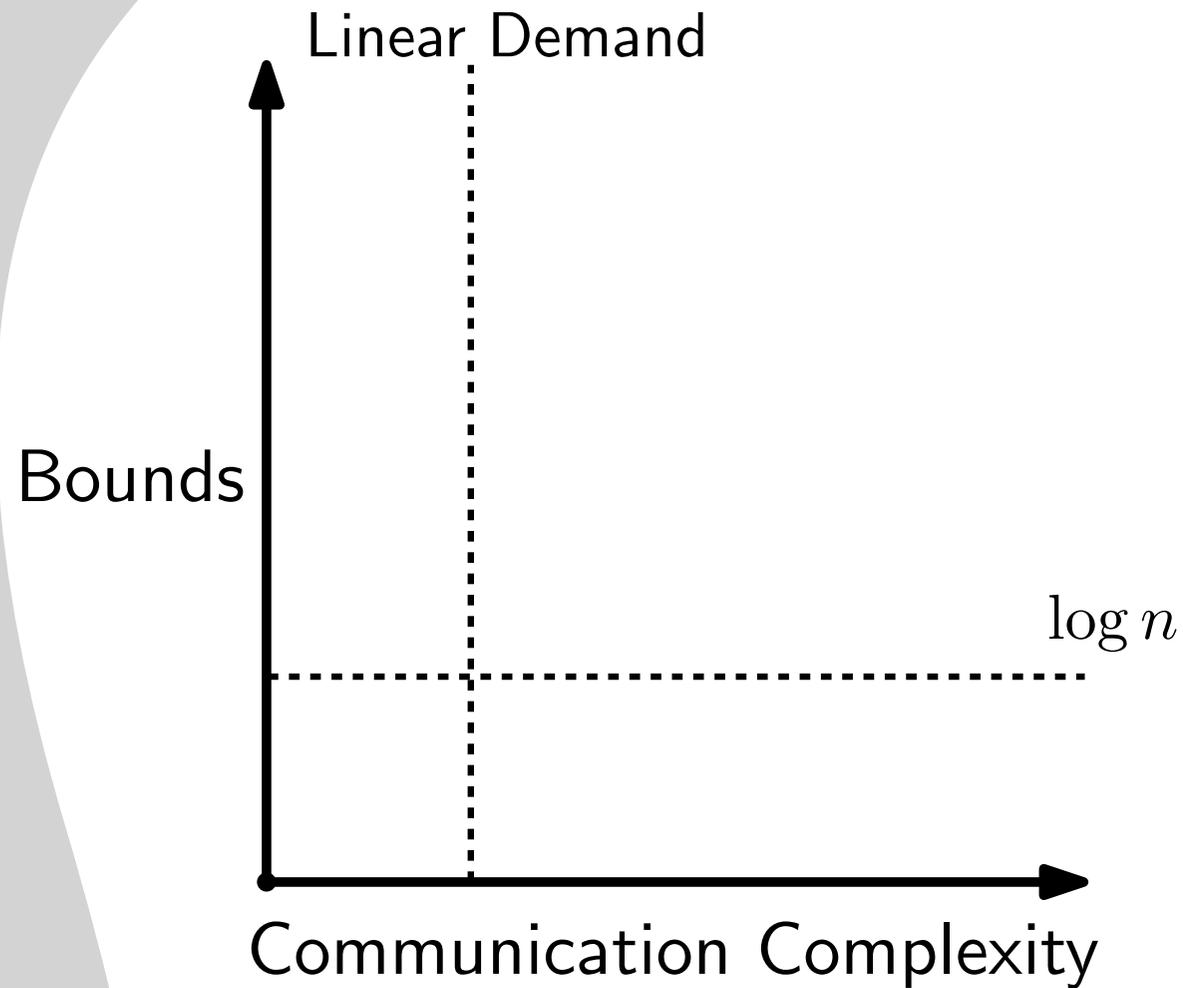
- Bounds on Competitive Ratio



# Results

Grid Networks:

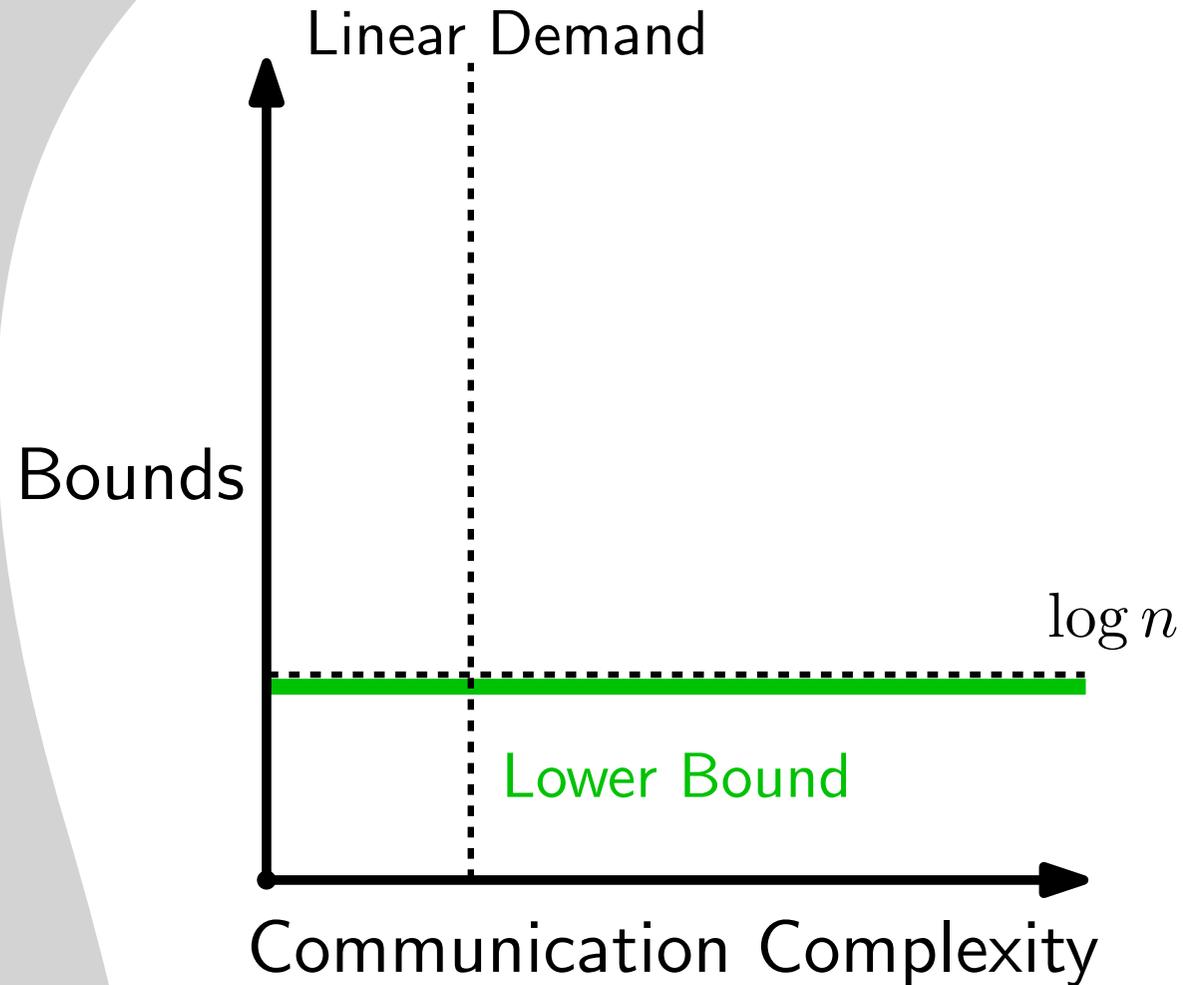
- Bounds on Competitive Ratio
- Restricted Communication



# Results

Grid Networks:

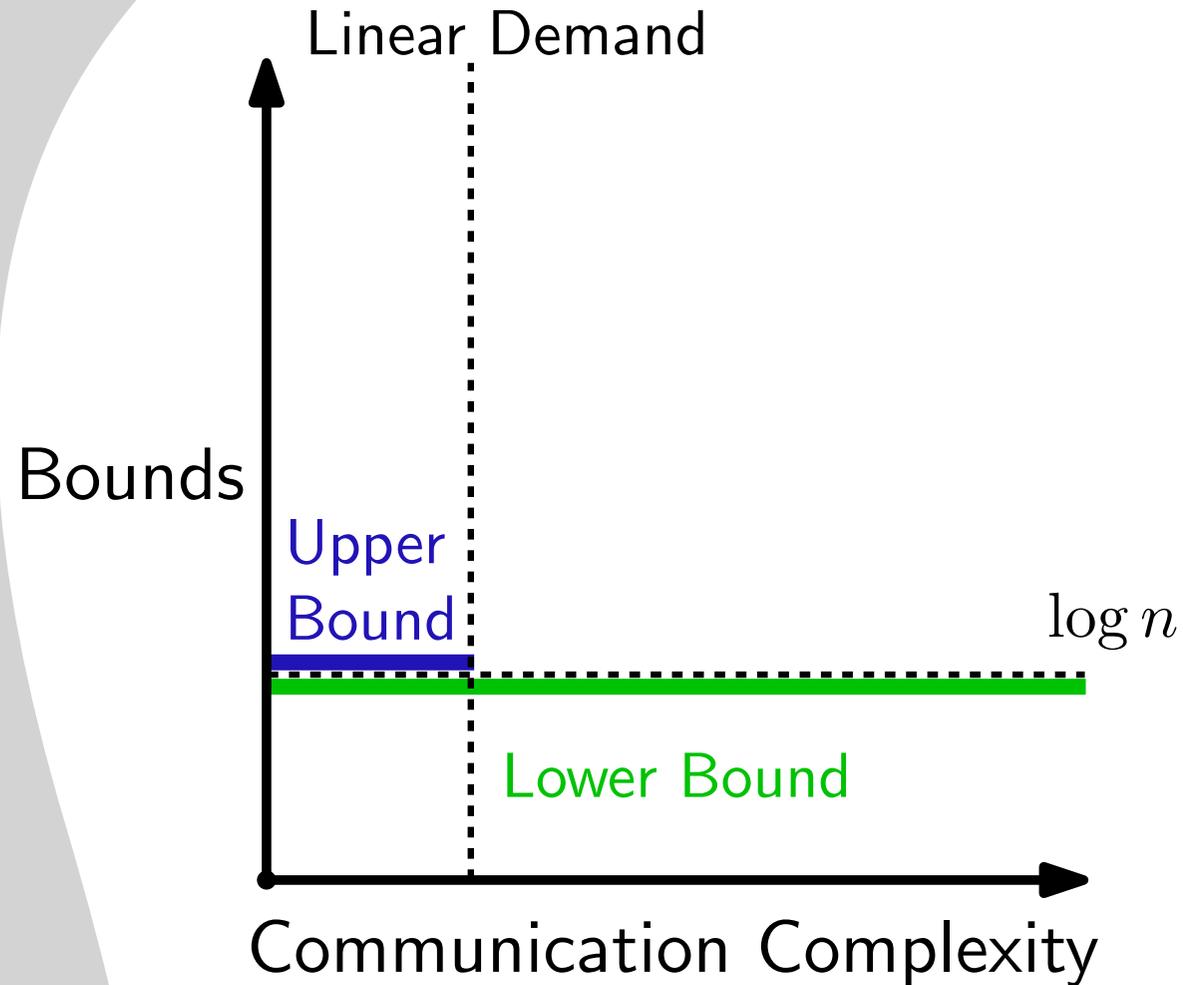
- Bounds on Competitive Ratio
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# Results

Grid Networks:

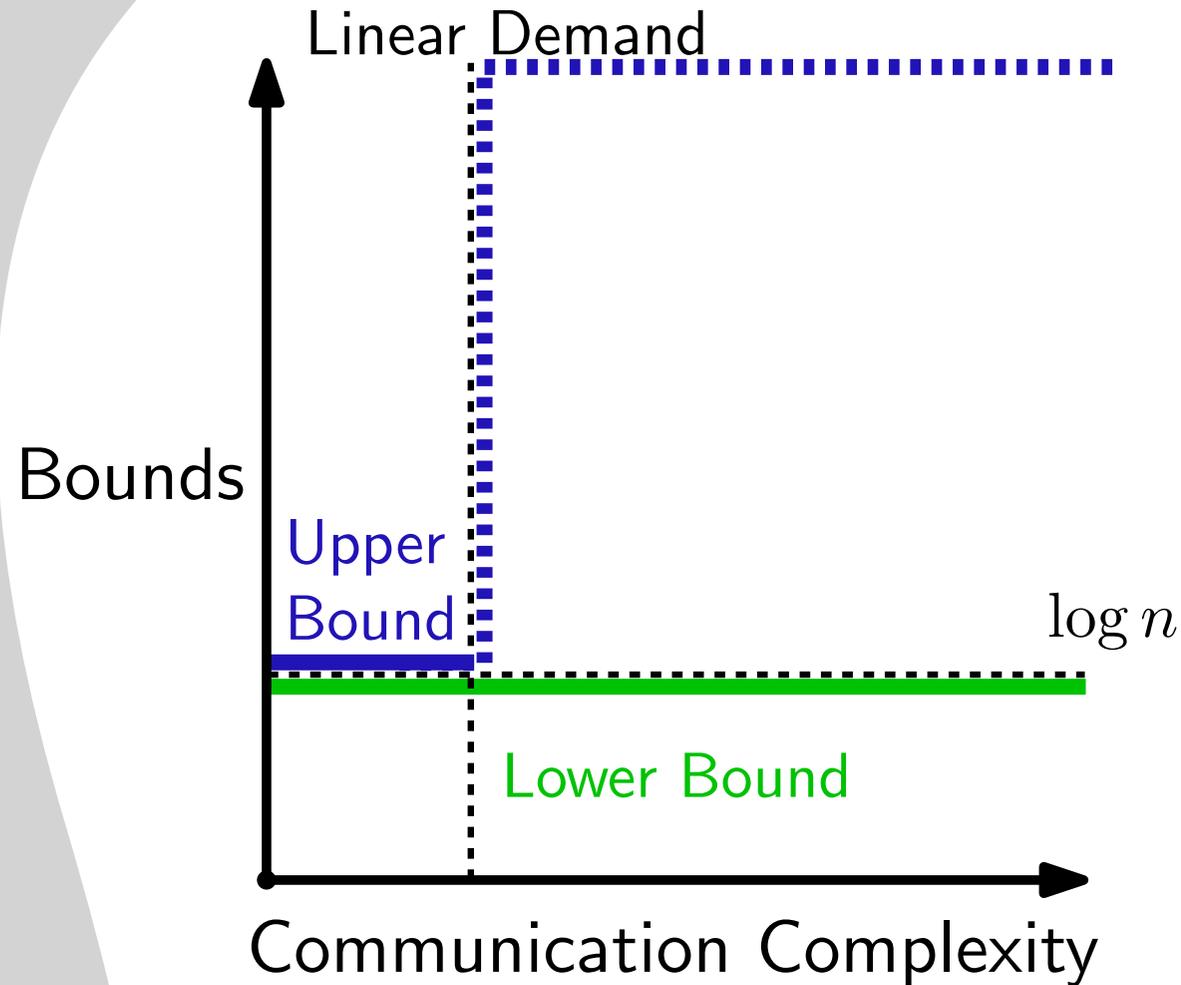
- Bounds on Competitive Ratio
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# Results

Grid Networks:

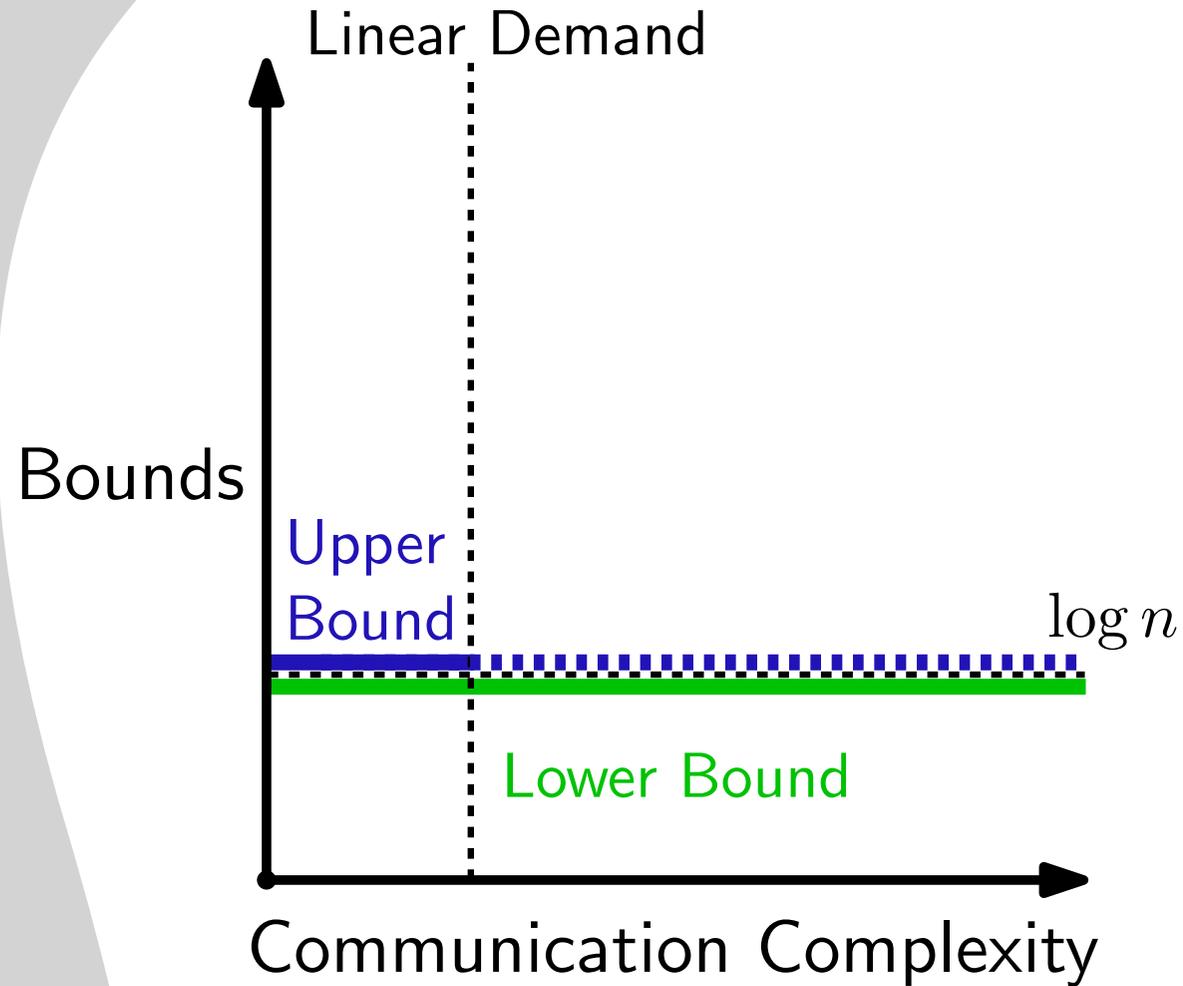
- Bounds on Competitive Ratio
- Restricted Communication



# Results

Grid Networks:

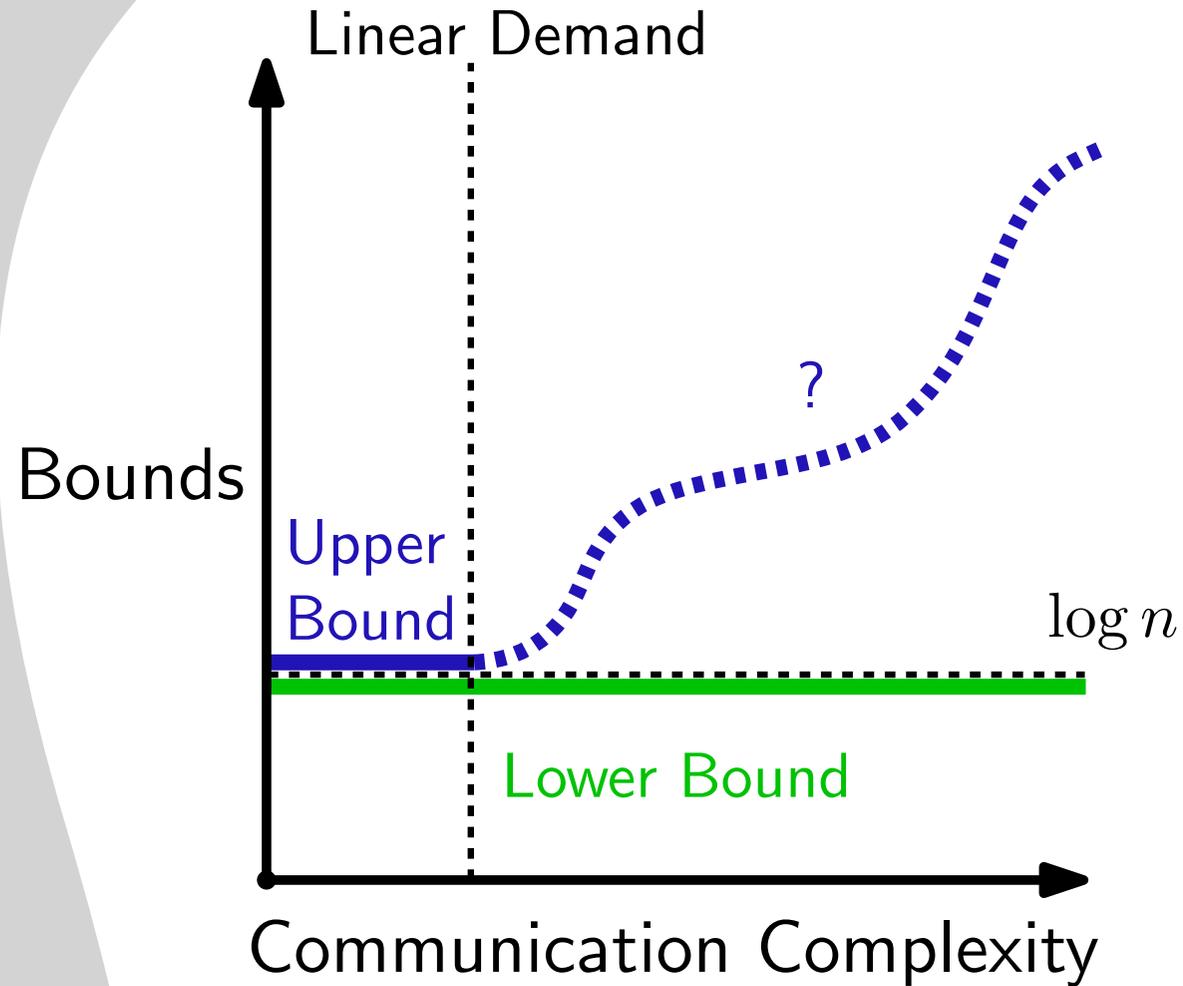
- Bounds on Competitive Ratio
- Restricted Communication



# Results

Grid Networks:

- Bounds on Competitive Ratio
- Restricted Communication



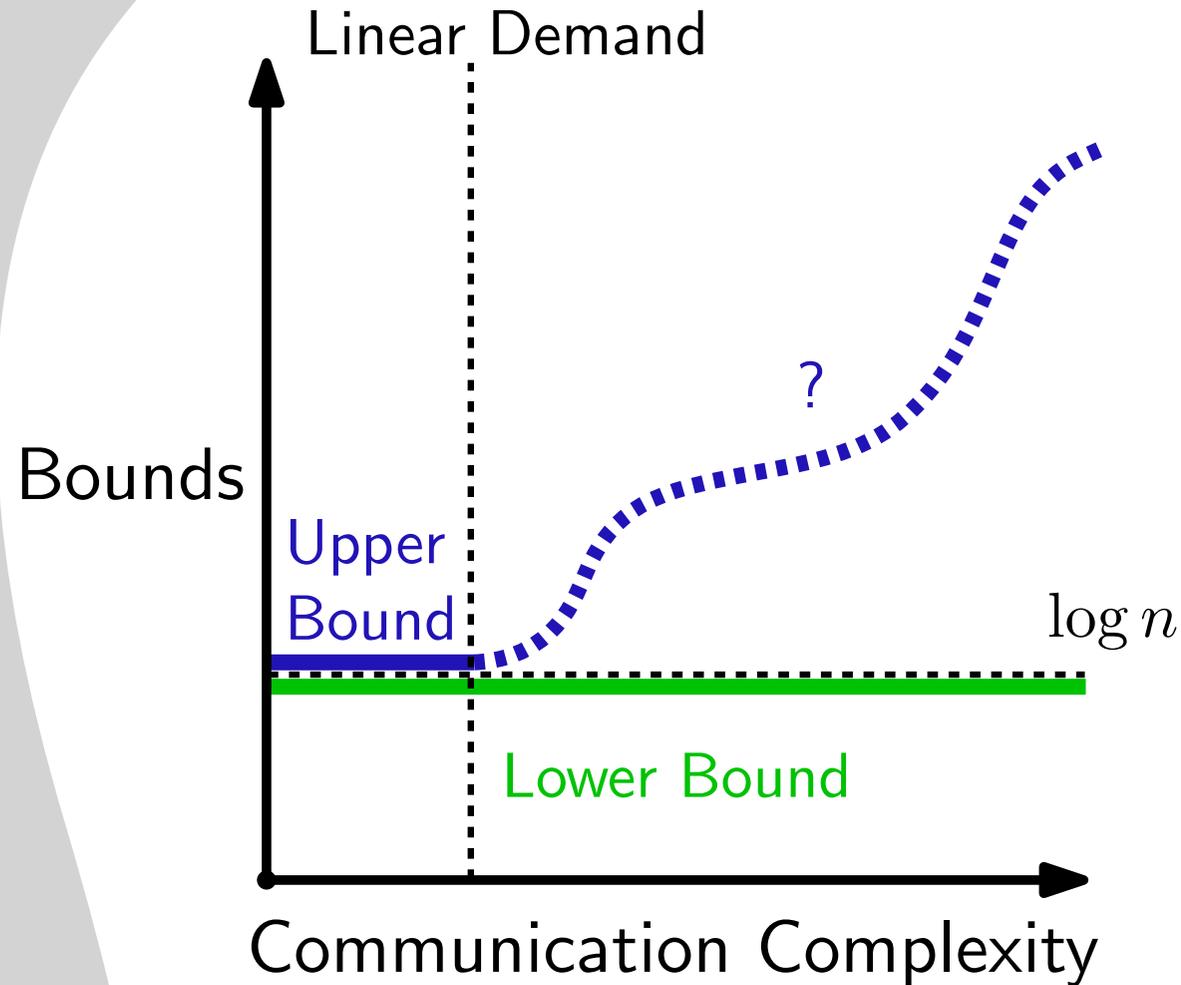
# Results

Grid Networks:

- Bounds on Competitive Ratio
- Restricted Communication

Line Networks:

- Distributed Implementation

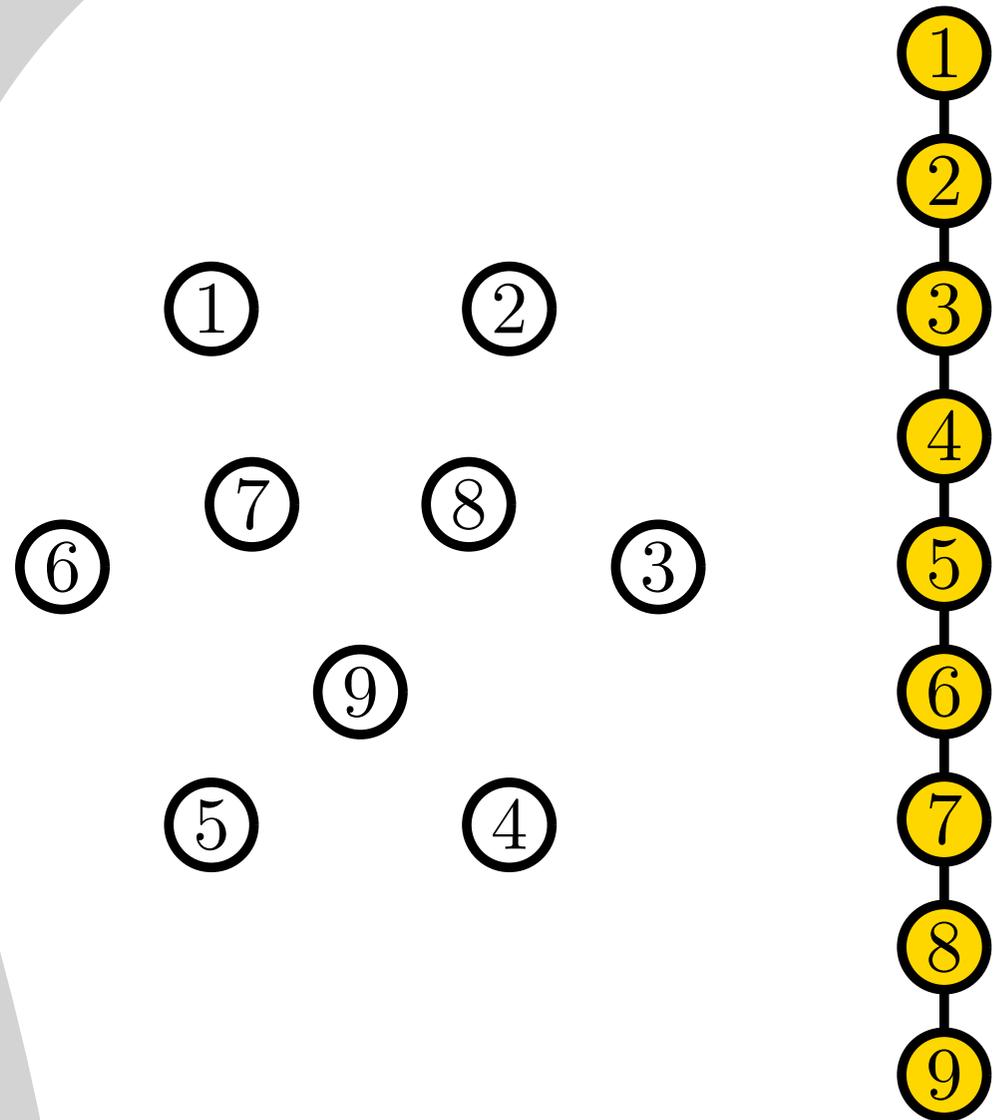


# Lower Bound

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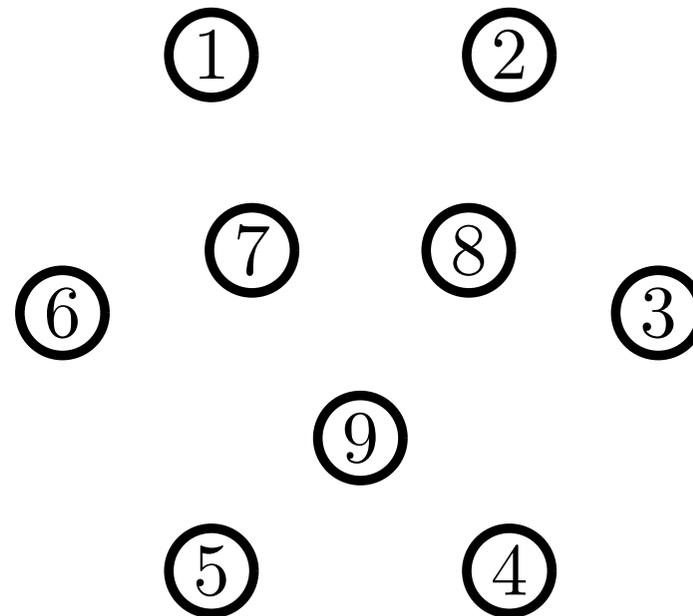


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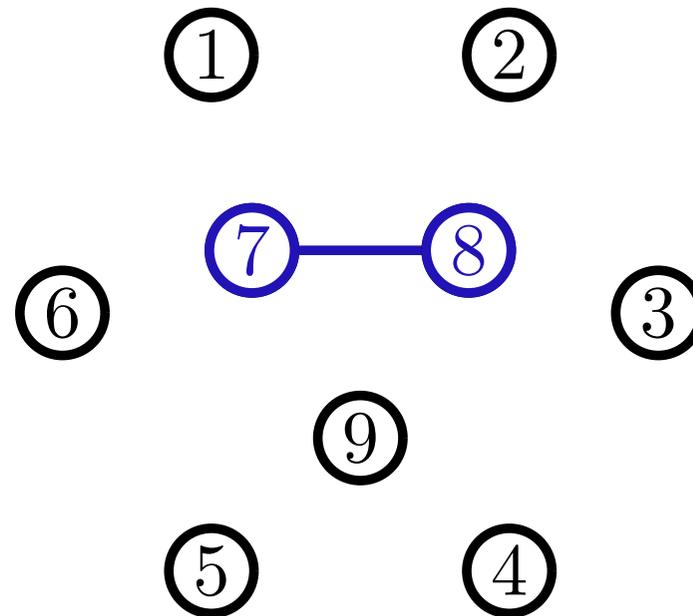
# Lower Bound

$\sigma =$



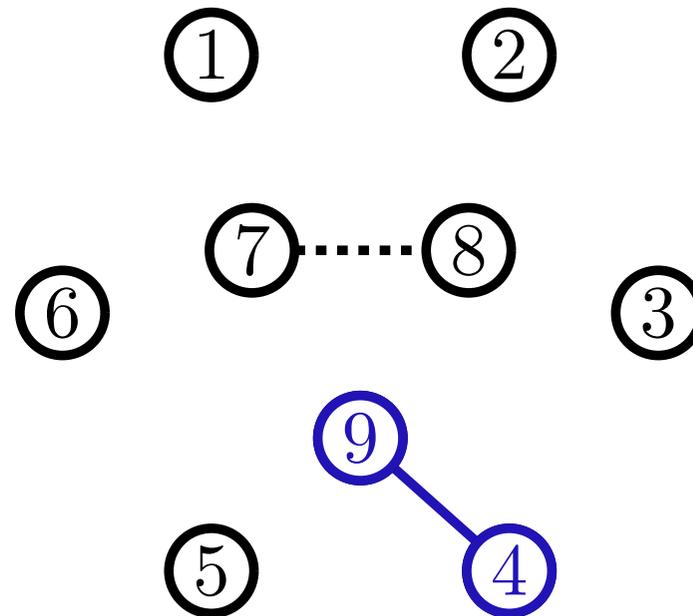
# Lower Bound

$$\sigma = (7, 8)$$



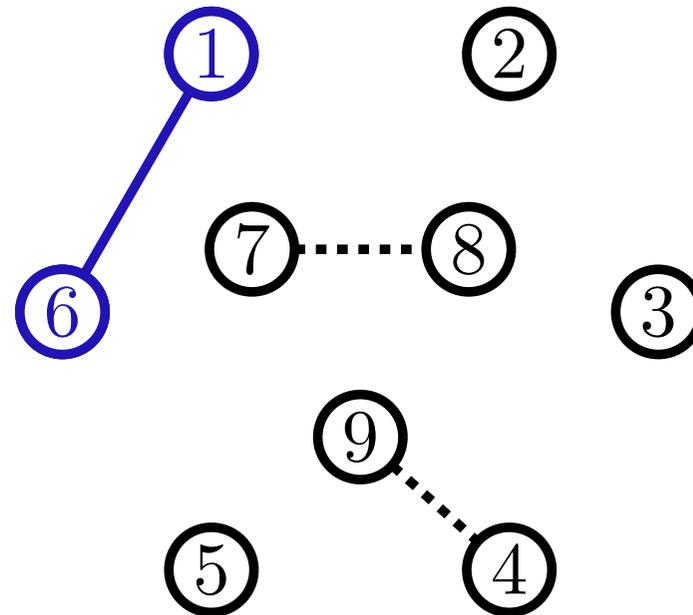
# Lower Bound

$$\sigma = (7, 8)(4, 9)$$



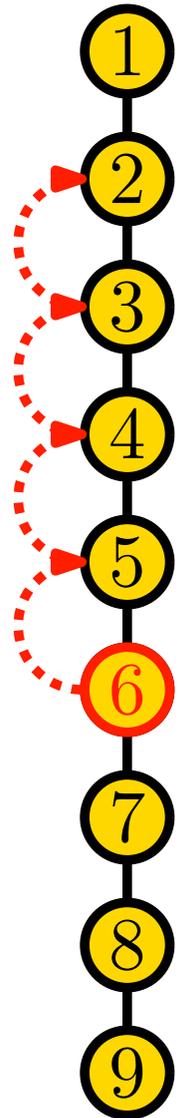
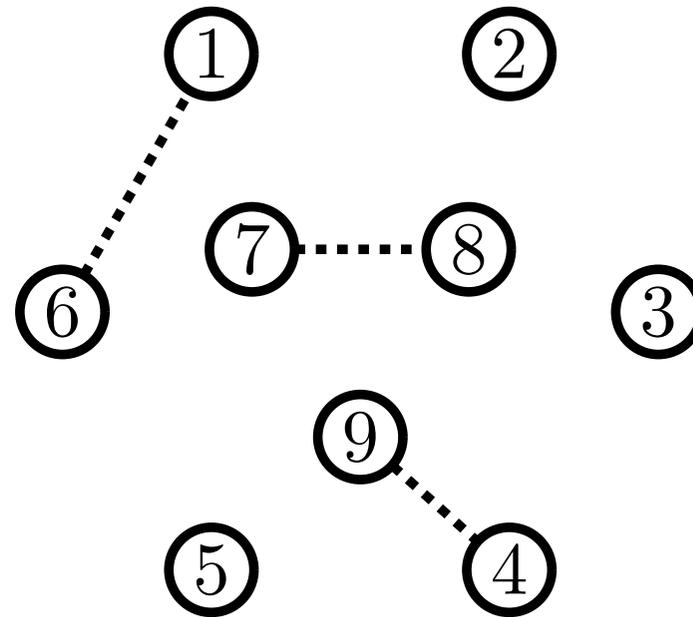
# Lower Bound

$$\sigma = (7, 8)(4, 9)(1, 6)$$



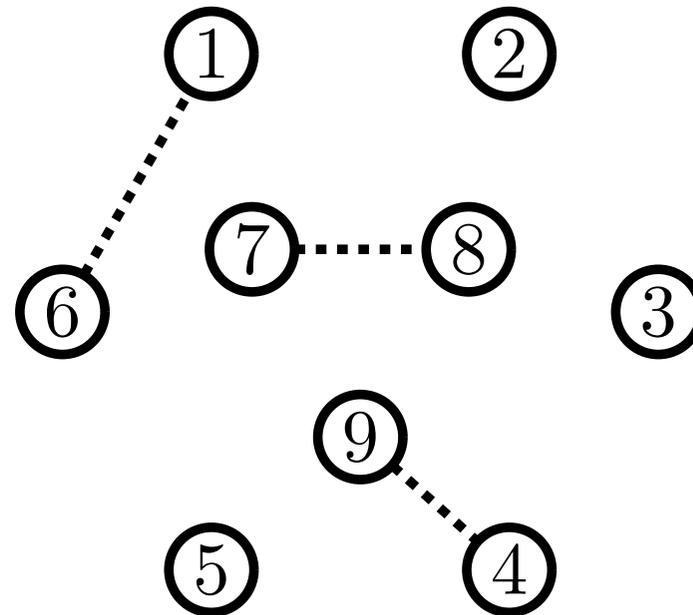
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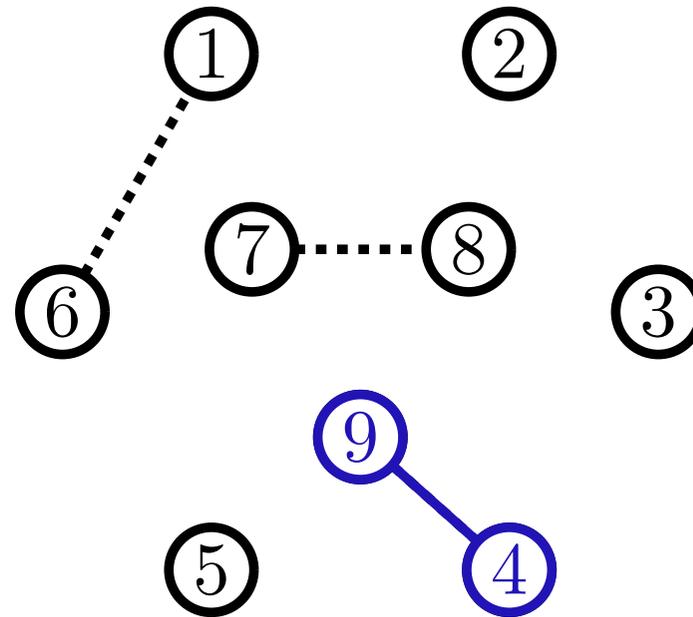
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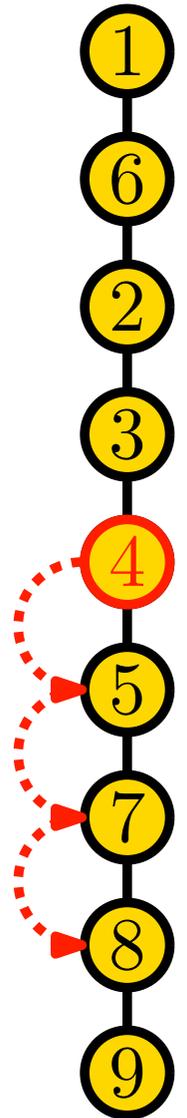
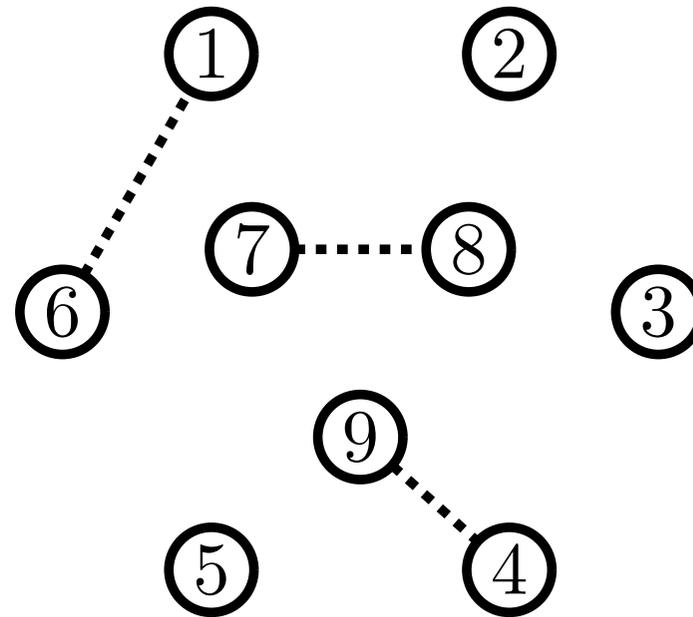
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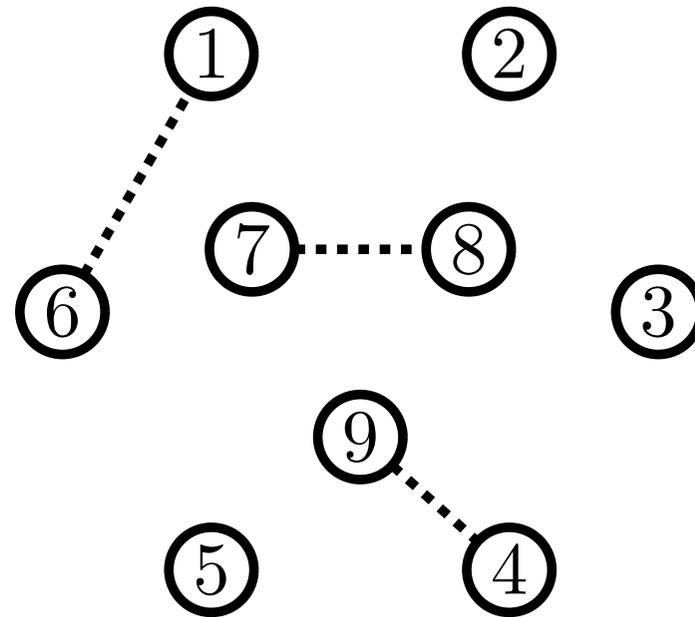
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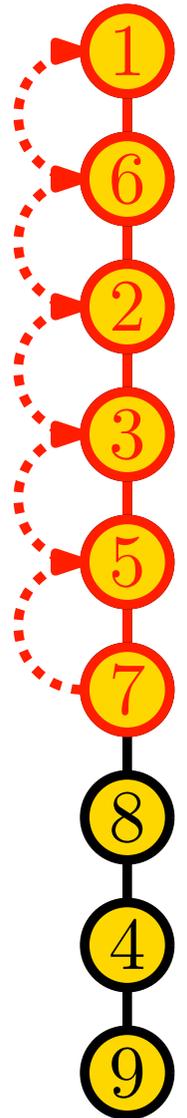
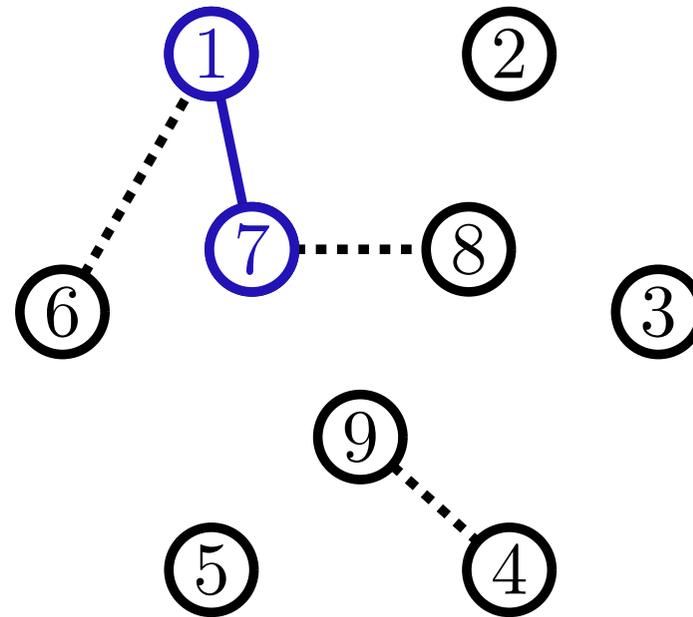
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$$\sigma = (7, 8)(4, 9)(1, 6)(4, 9)$$



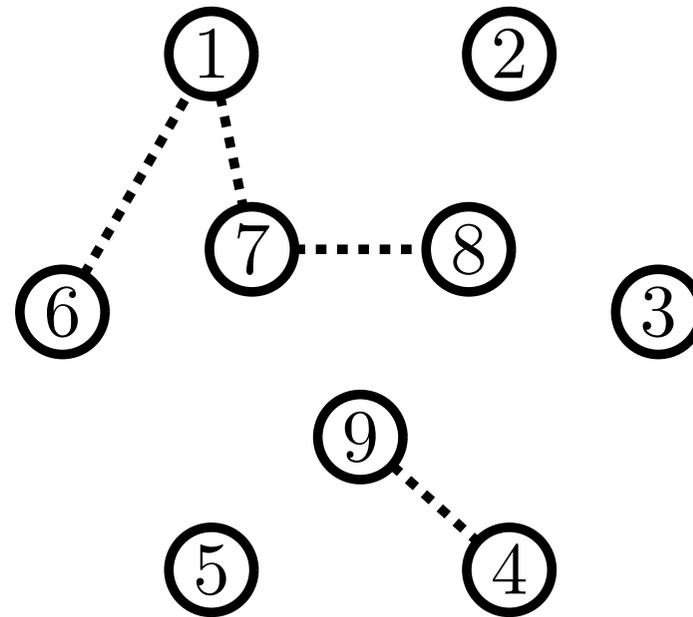
# Lower Bound

$$\sigma = (7, 8)(4, 9)(1, 6)(4, 9)(1, 7)$$



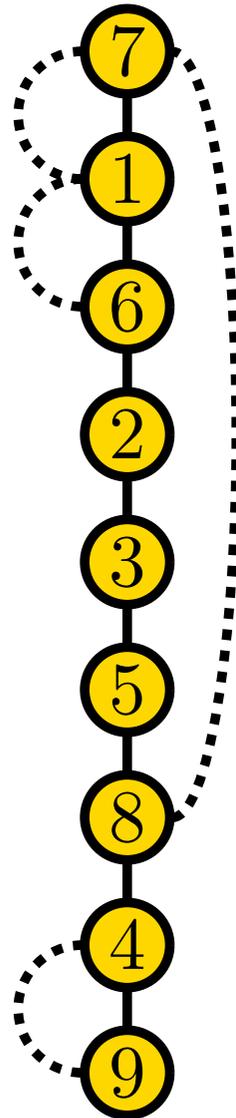
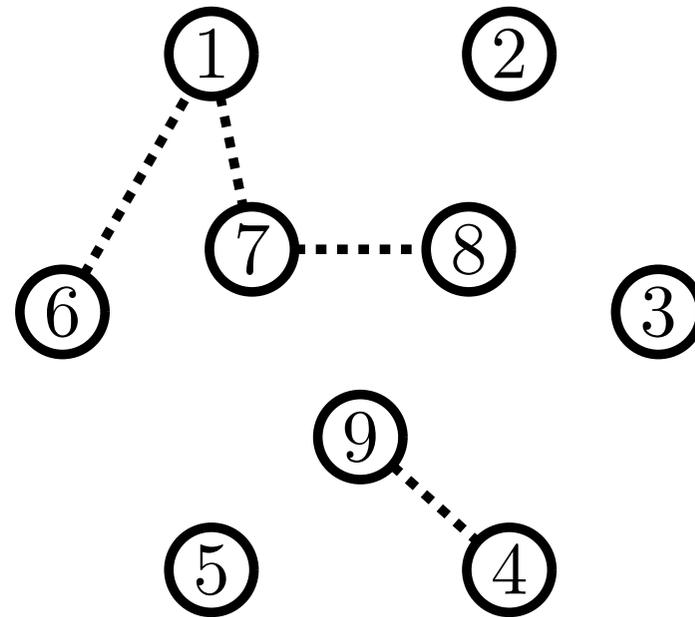
# Lower Bound

$$\sigma = (7, 8)(4, 9)(1, 6)(4, 9)(1, 7)$$



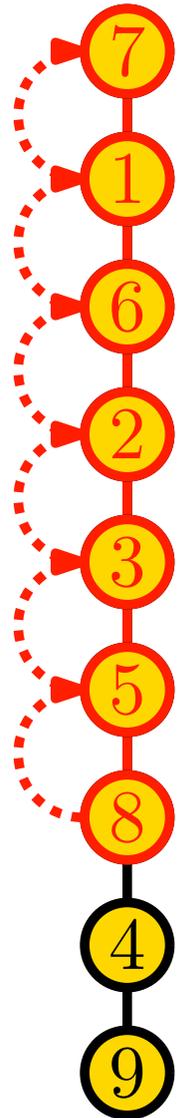
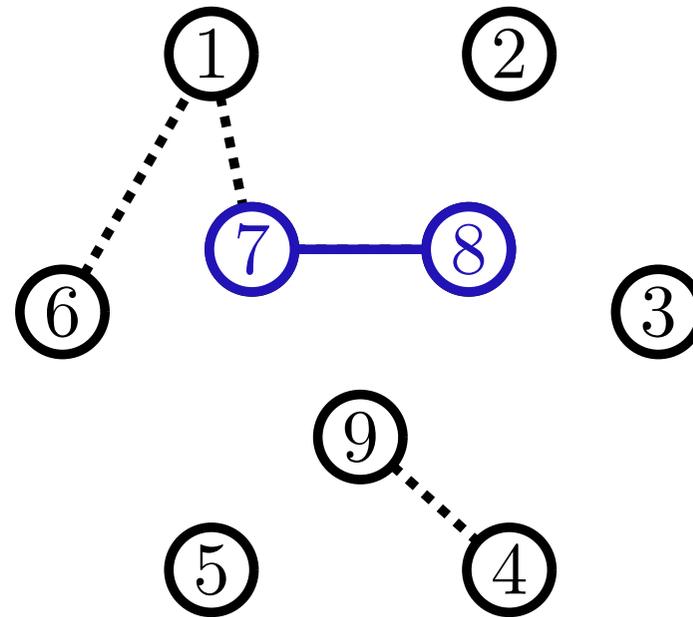
# Lower Bound

$$\sigma = (7, 8)(4, 9)(1, 6)(4, 9)(1, 7)$$



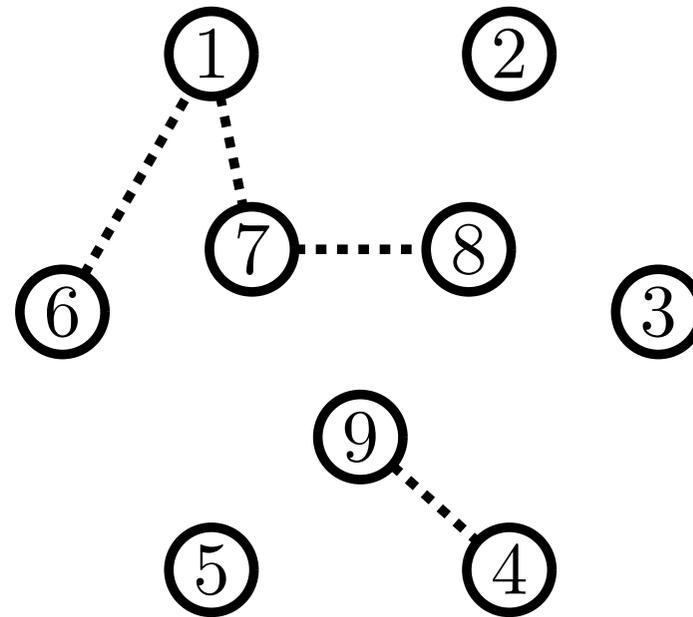
# Lower Bound

$$\sigma = (7, 8)(4, 9)(1, 6)(4, 9)(1, 7)(7, 8)$$



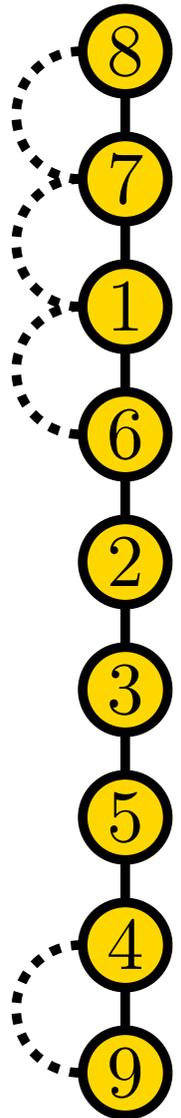
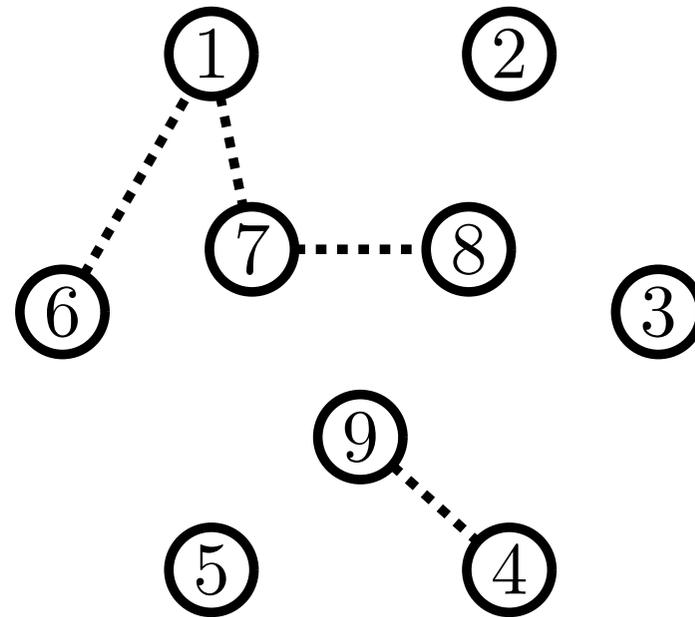
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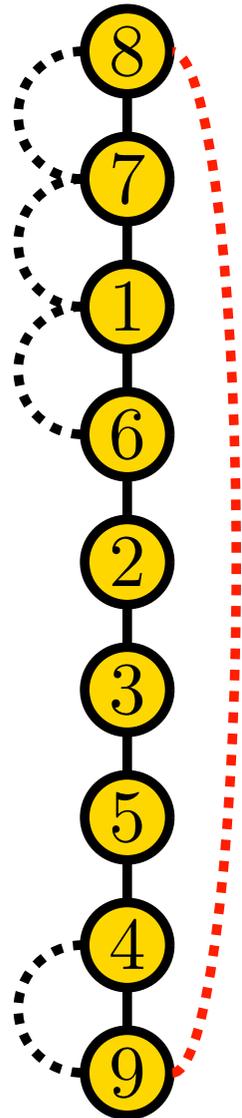
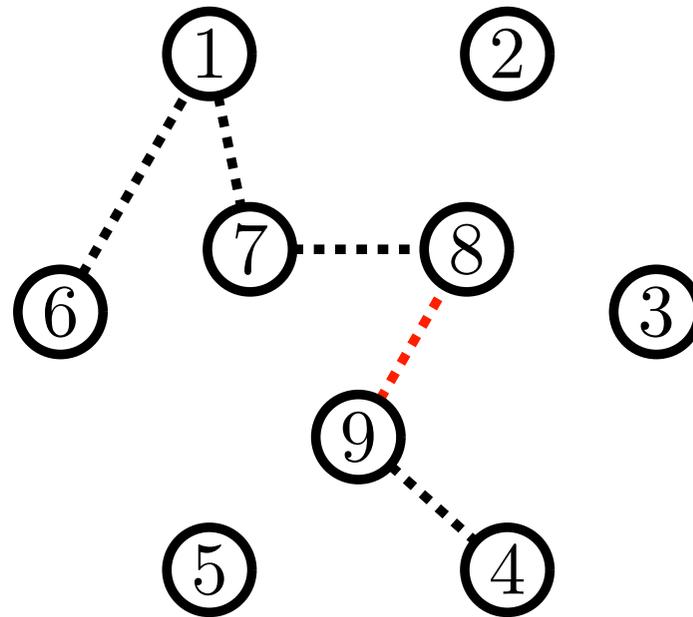


# Lower Bound

The Strategy:

- Exploit Bad Edges

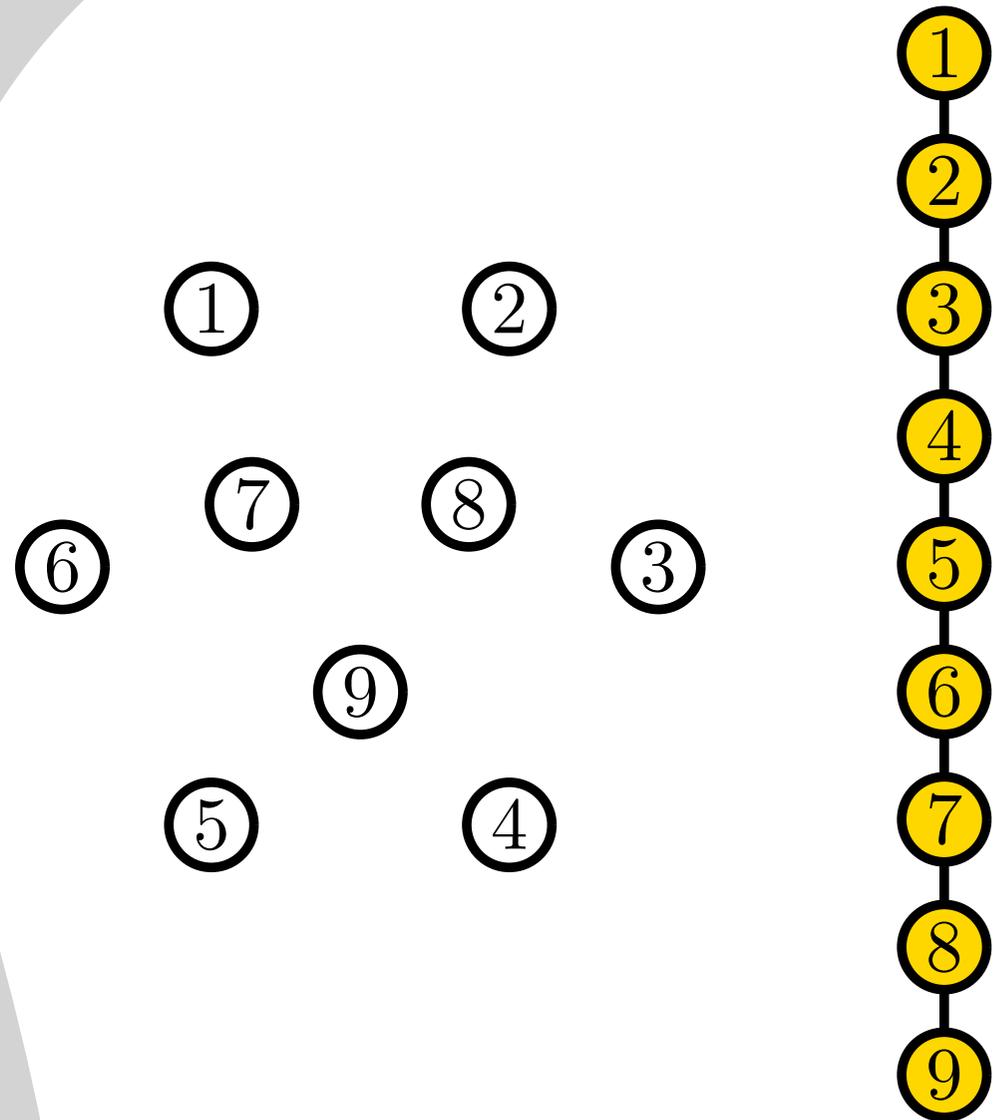
$$\sigma = (7, 8)(4, 9)(1, 6)(4, 9)(1, 7)(7, 8)$$



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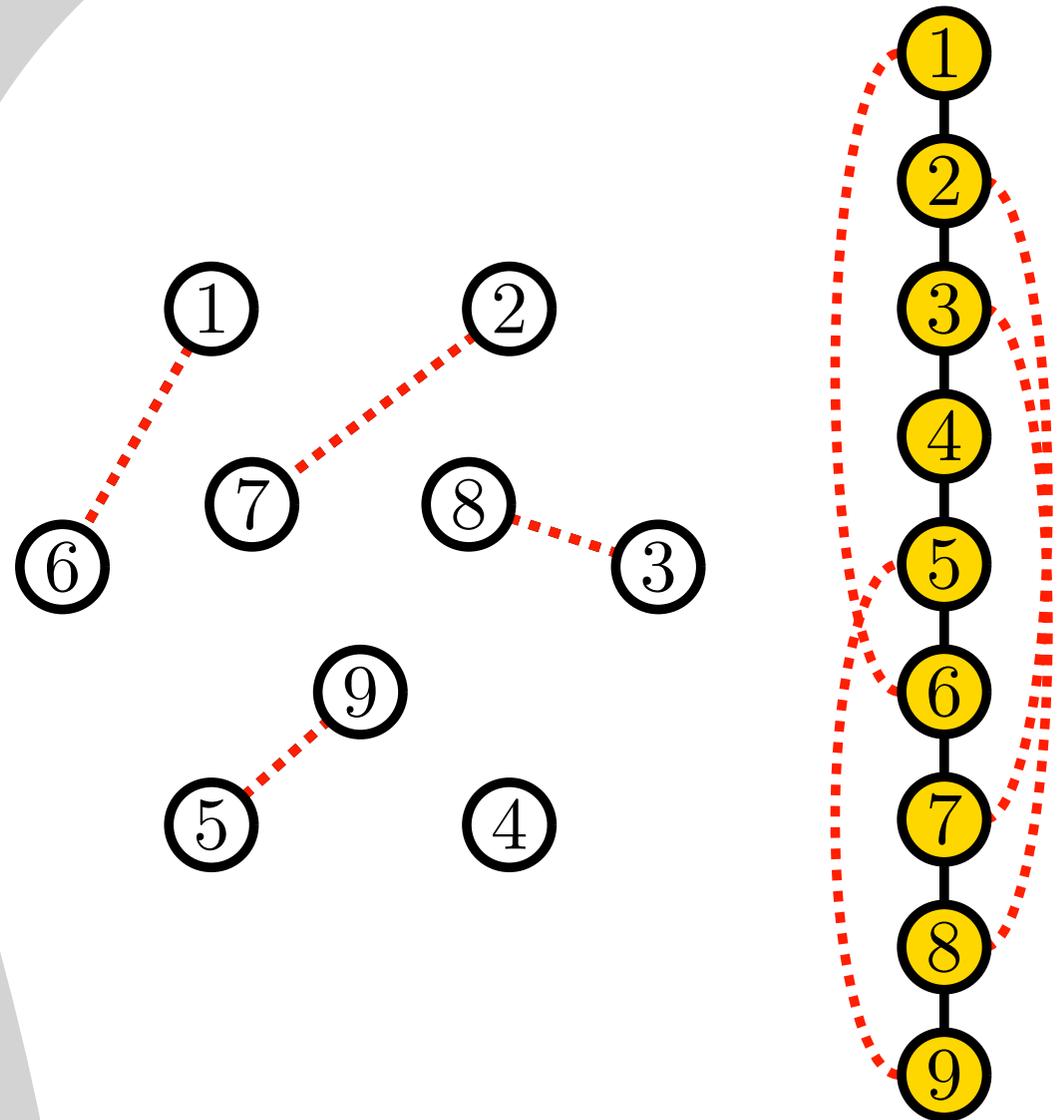
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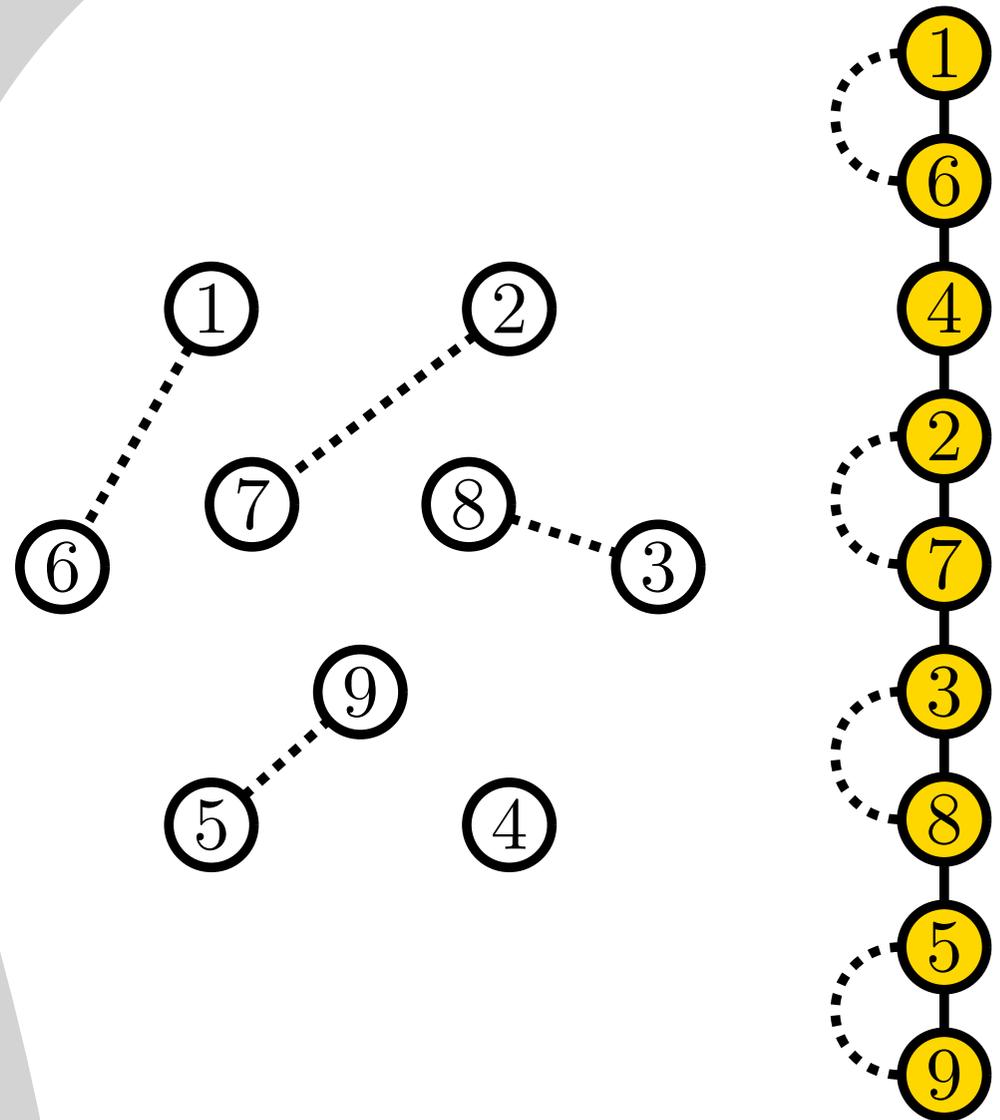
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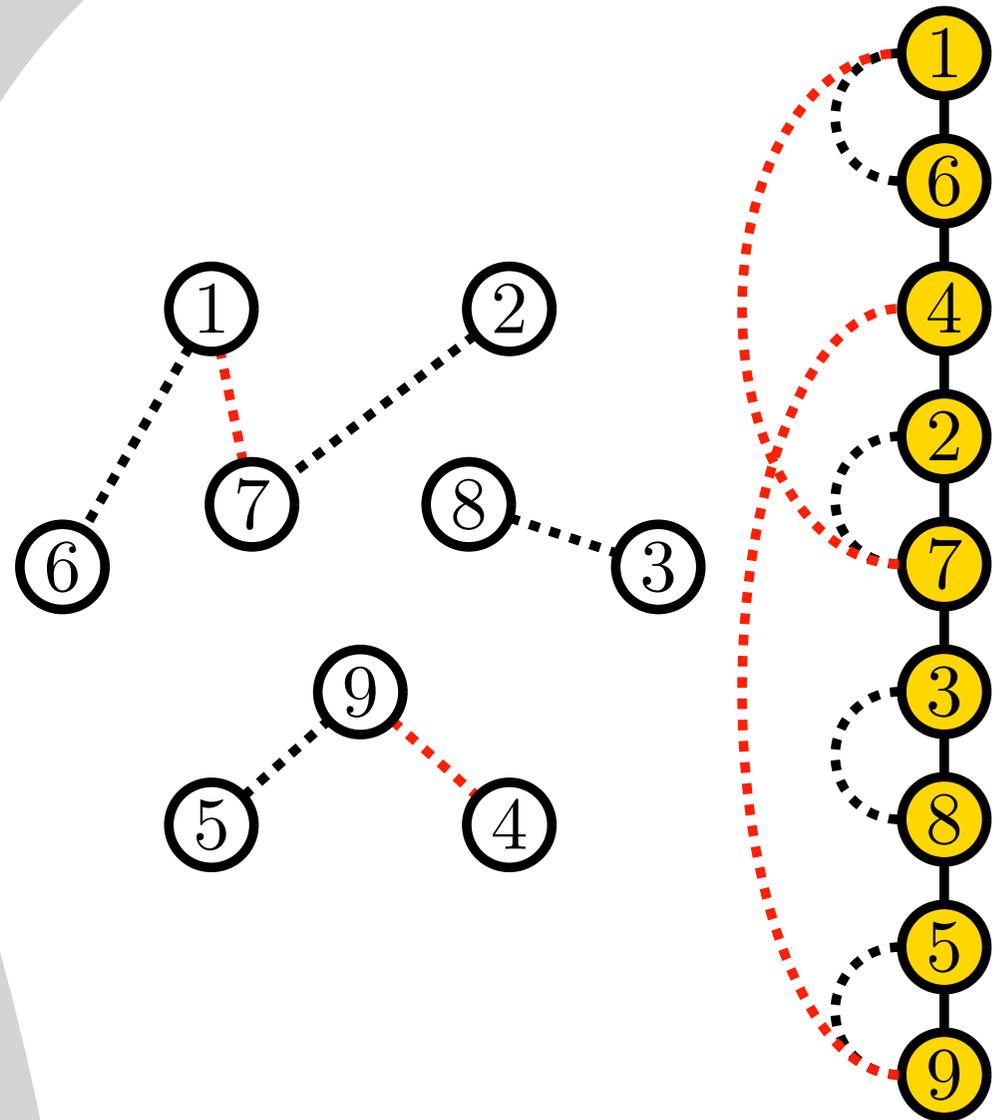
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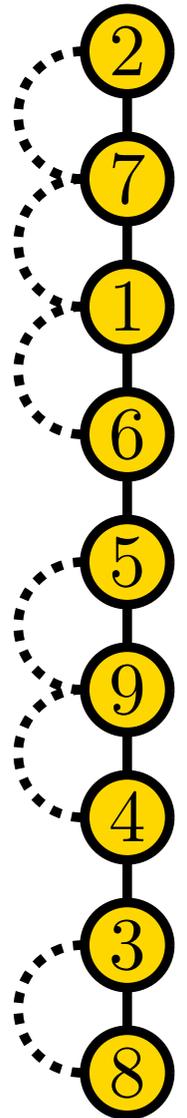
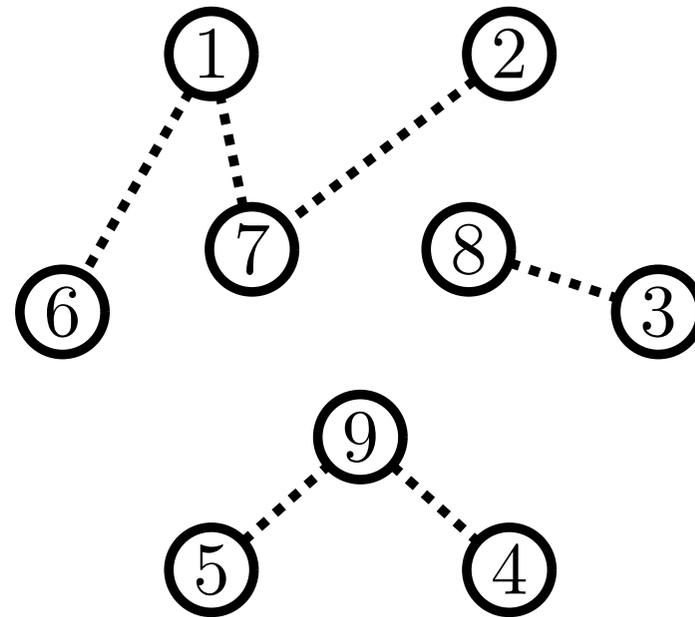
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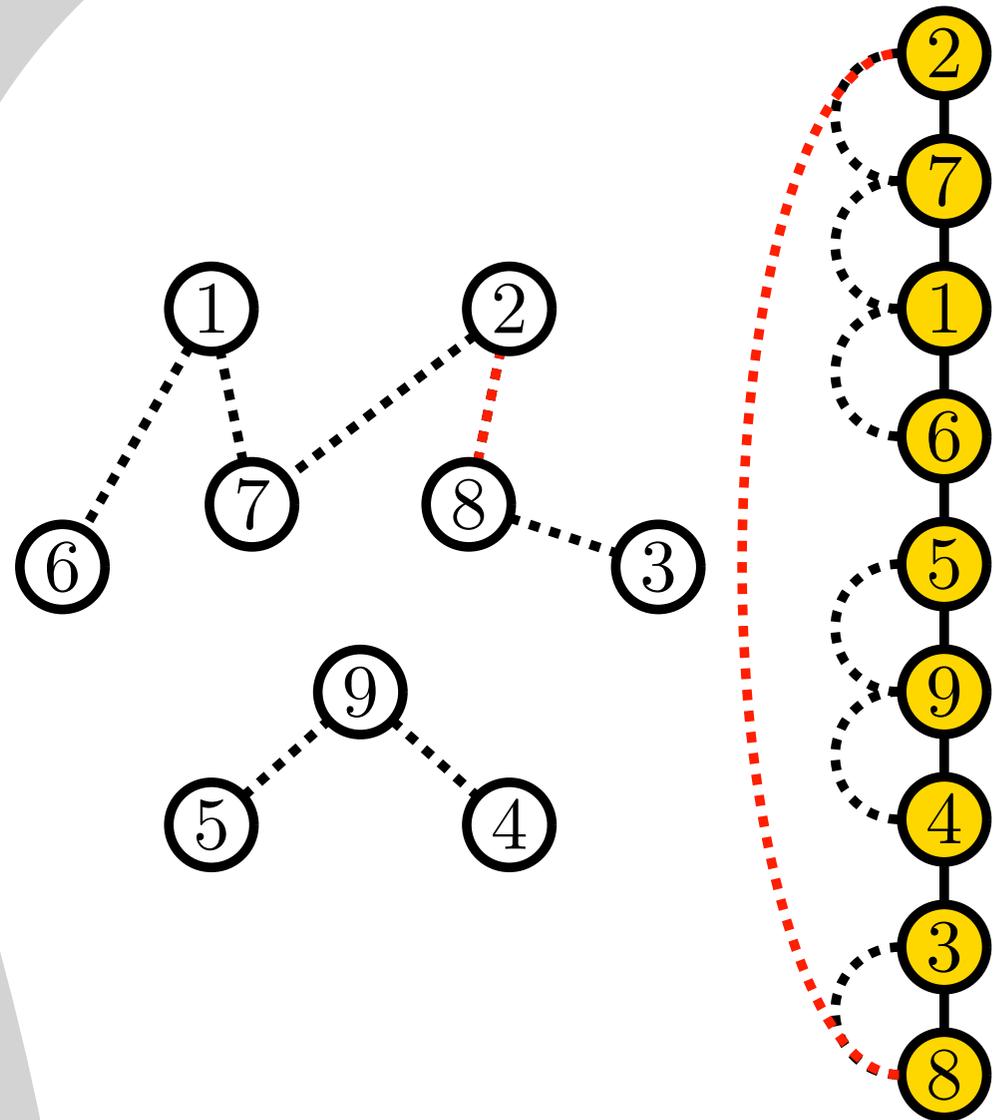
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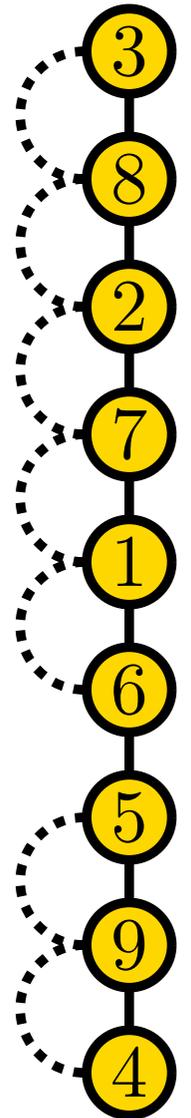
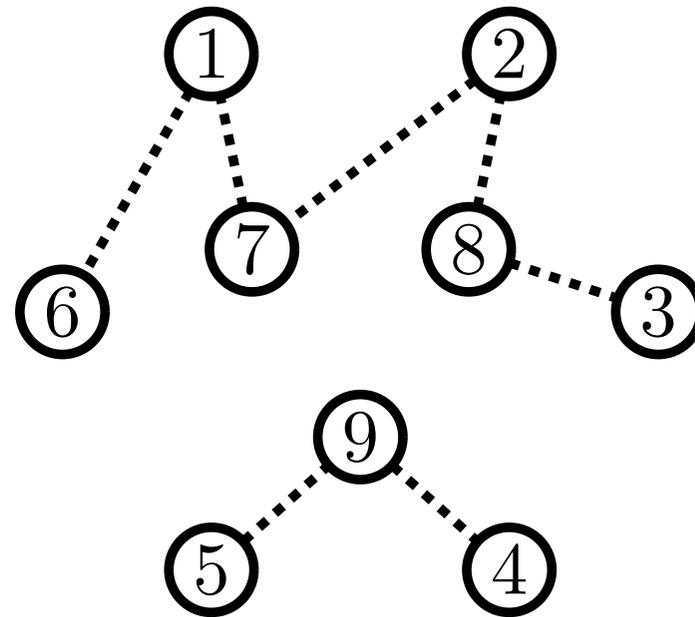
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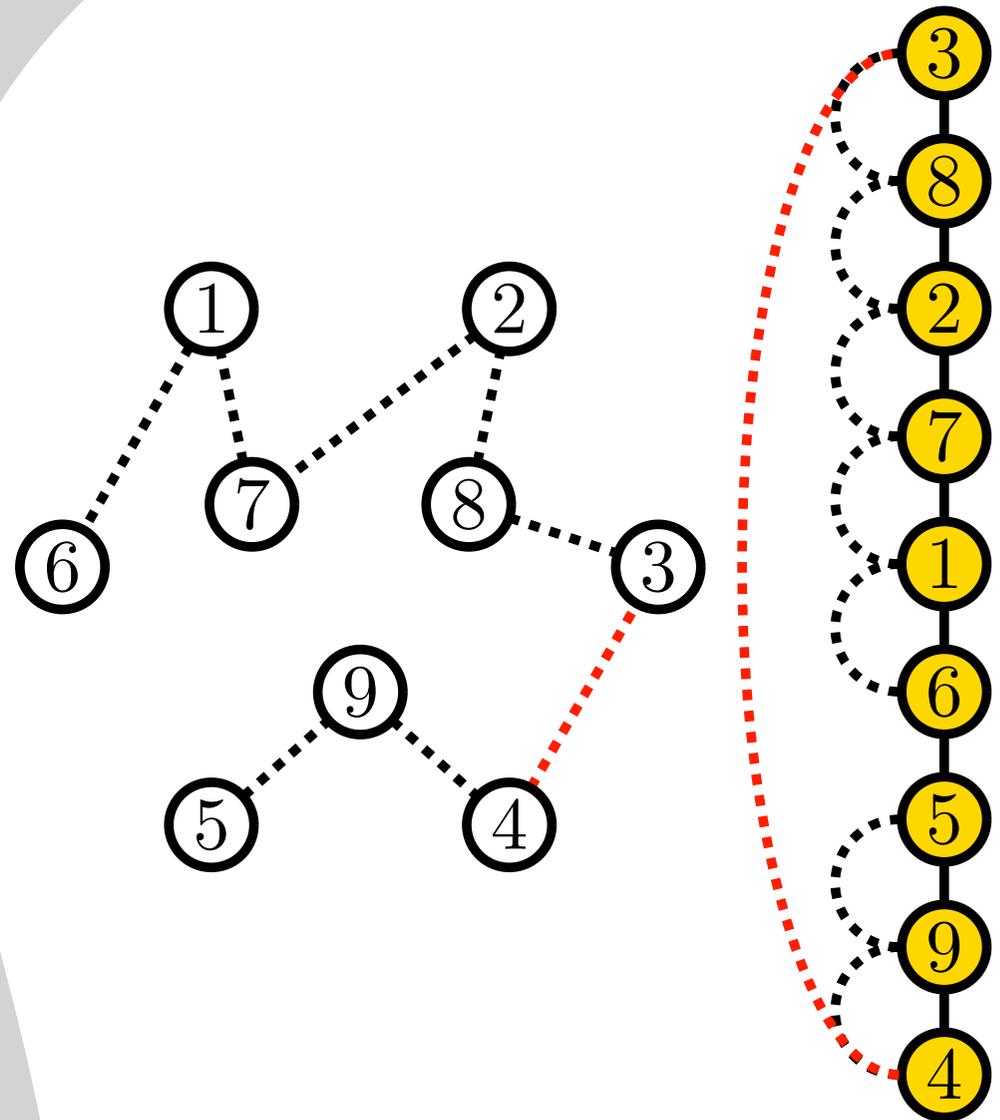
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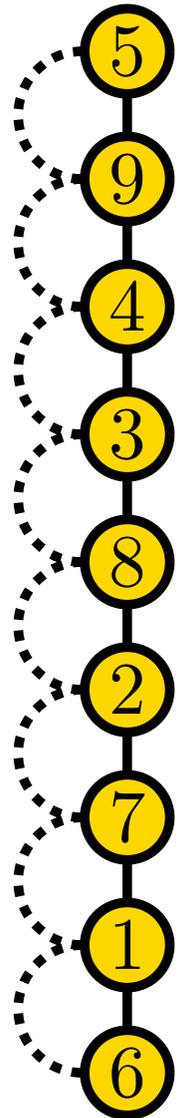
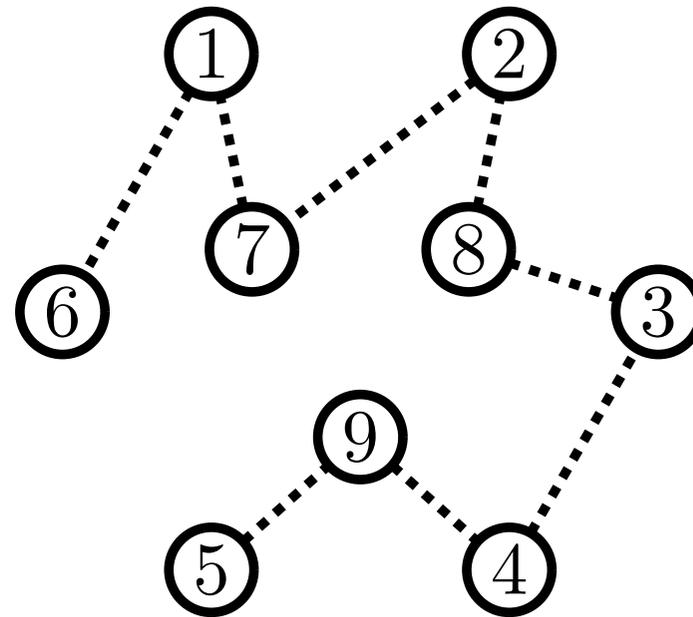
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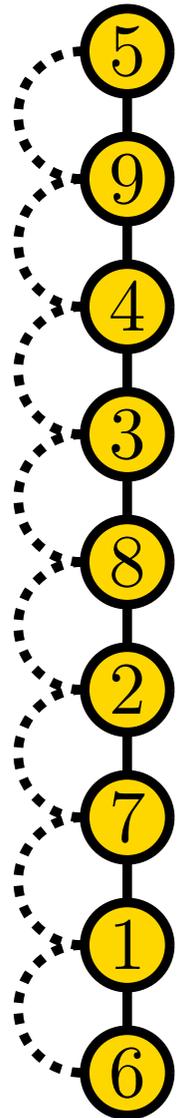
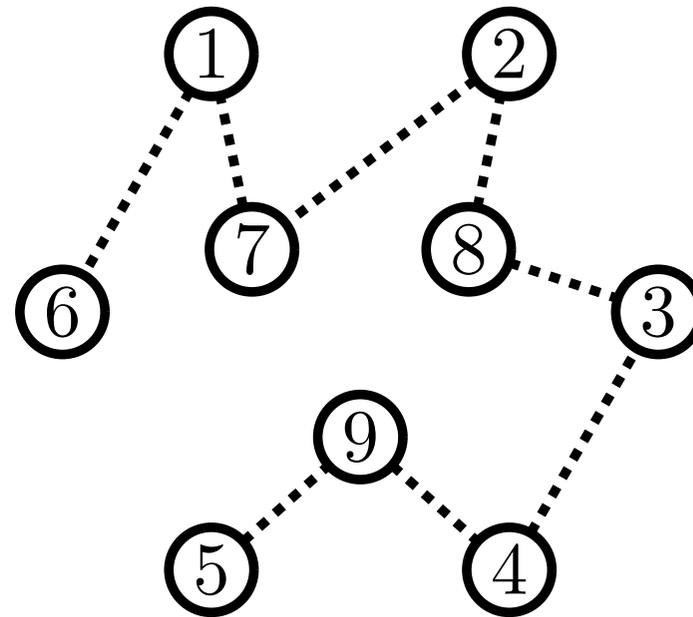
# Lower Bound

The Strategy:

- Exploit Bad Edges
- Introduce Bad Matching

The Basic Math:

- $\log n$  Bad Matchings
- $n^2$  Cost per Matching
- Online Cost:  $\Omega(n^2 \log n)$



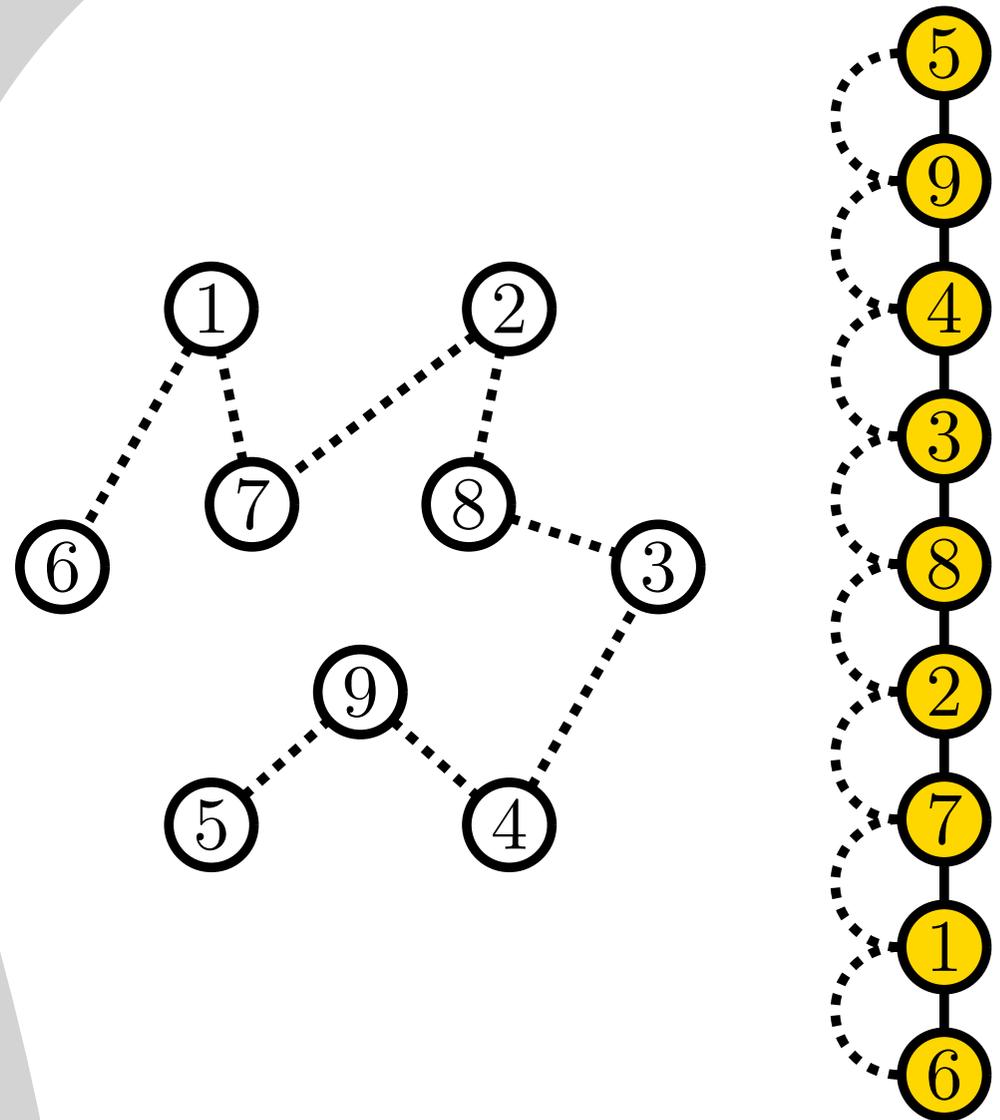
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The Strategy:

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- Maintain List Graph

The Basic Math:

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- Competitive Ratio:  $\Omega(\log n)$



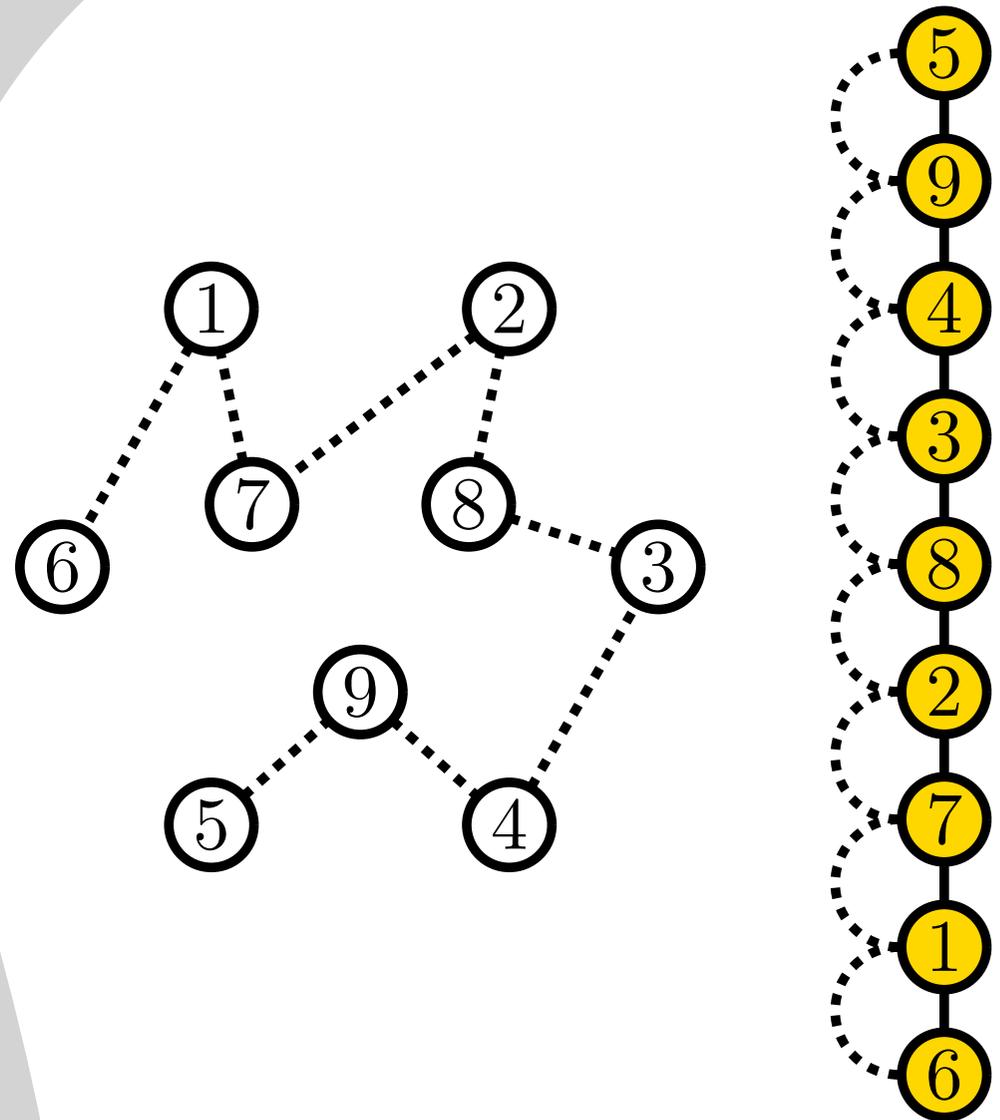
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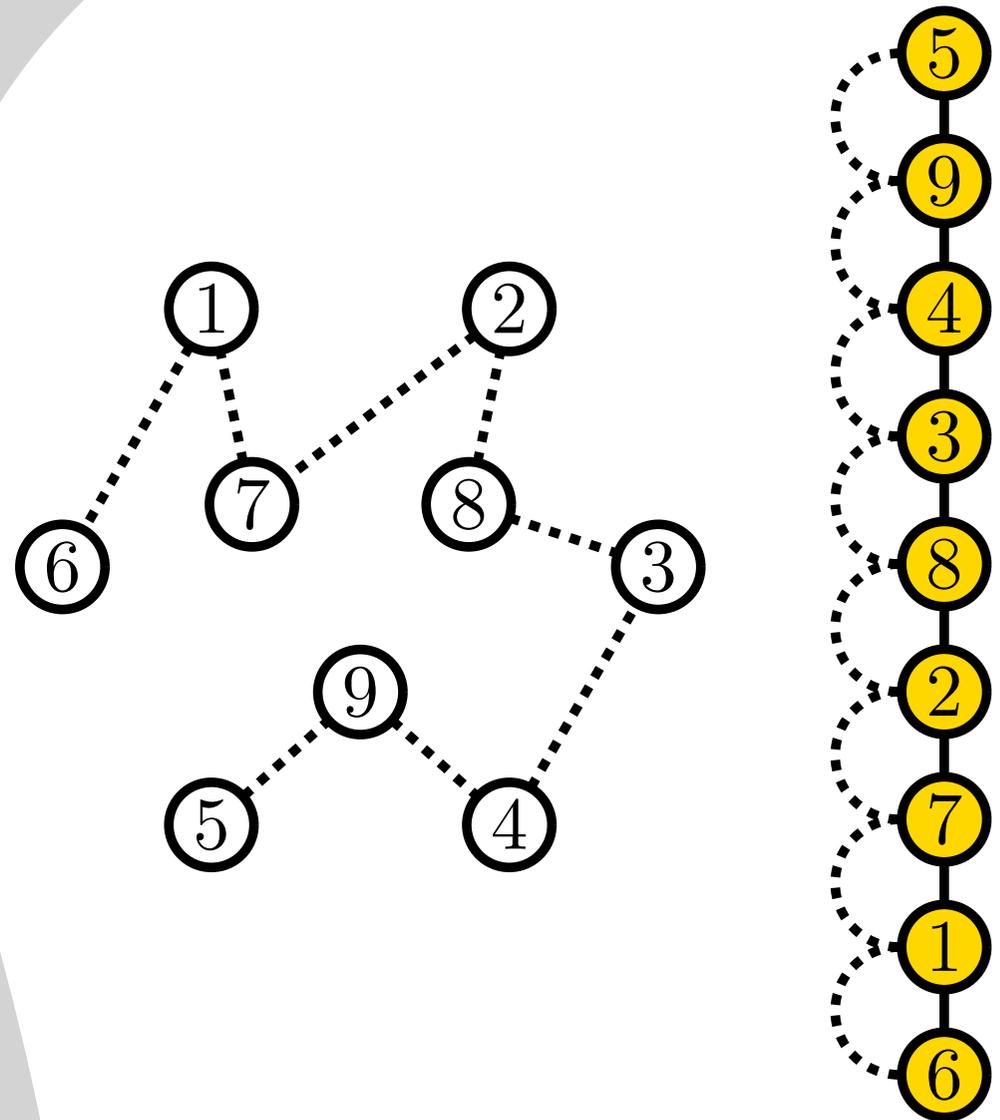
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$n^2$  per Matching:

- Quantify Distortion of Matching



# Lower Bound

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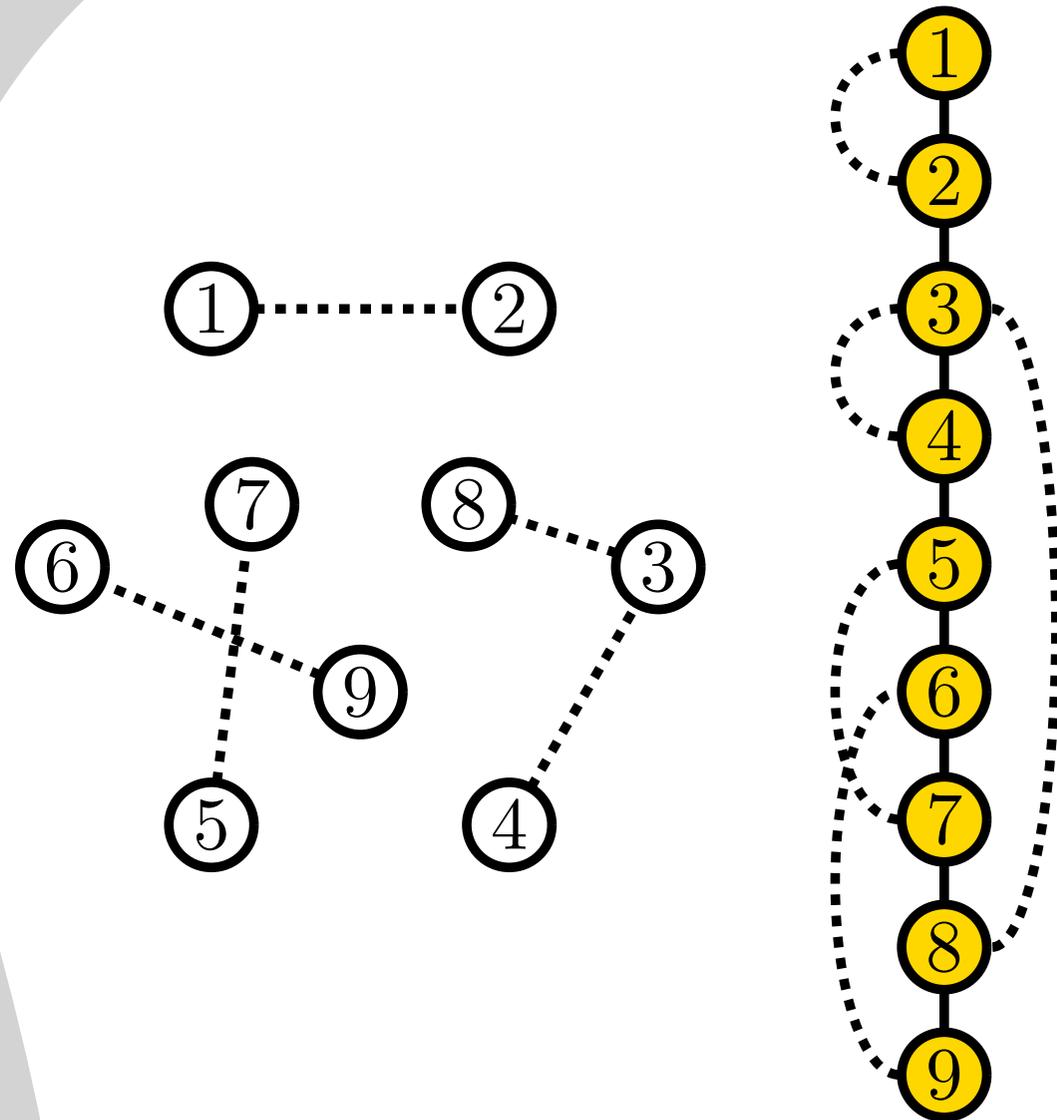
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# Lower Bound

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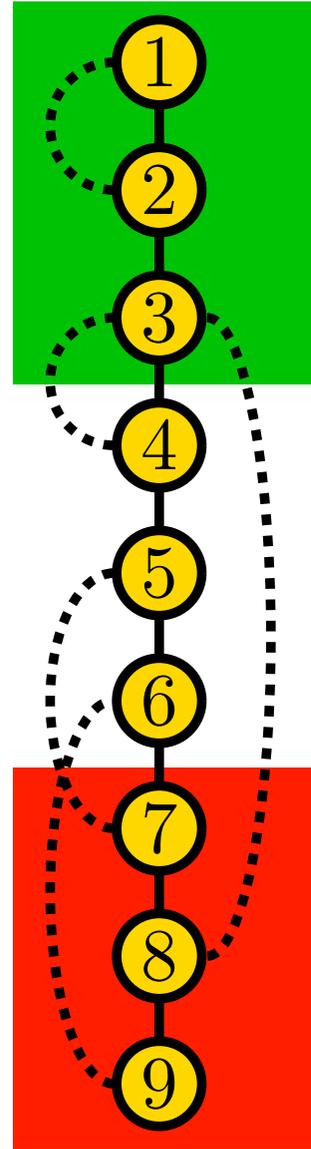
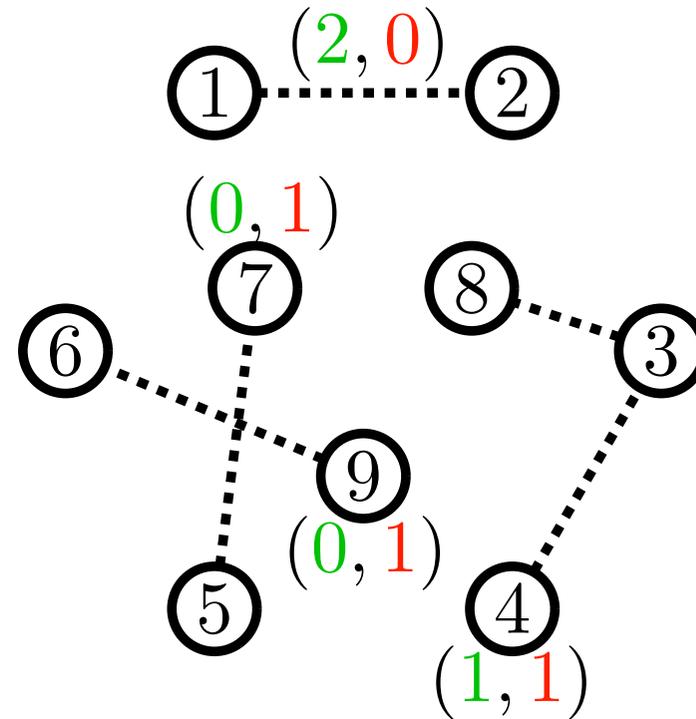
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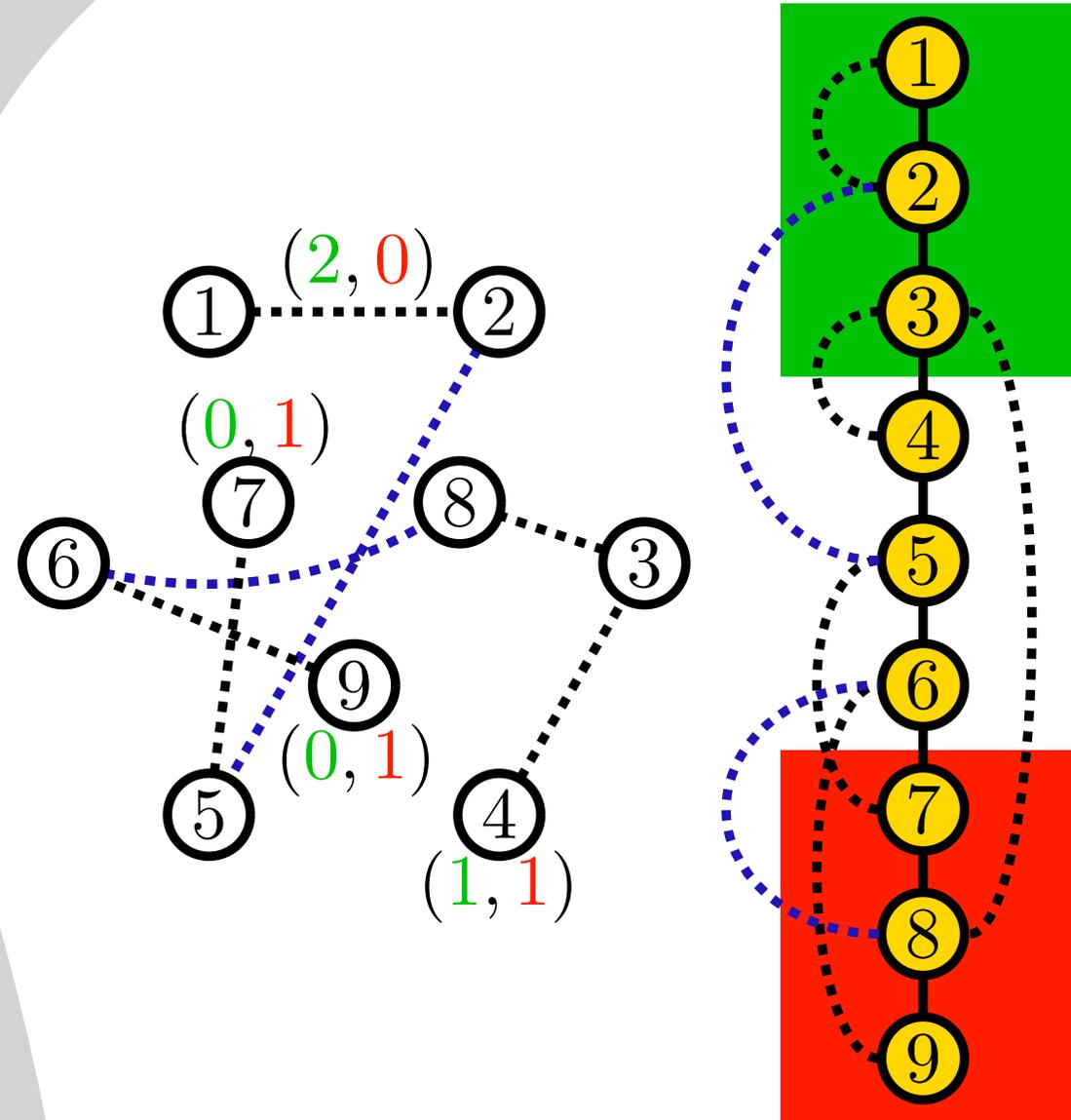
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# Lower Bound

The Strategy:

- Exploit Bad Edges
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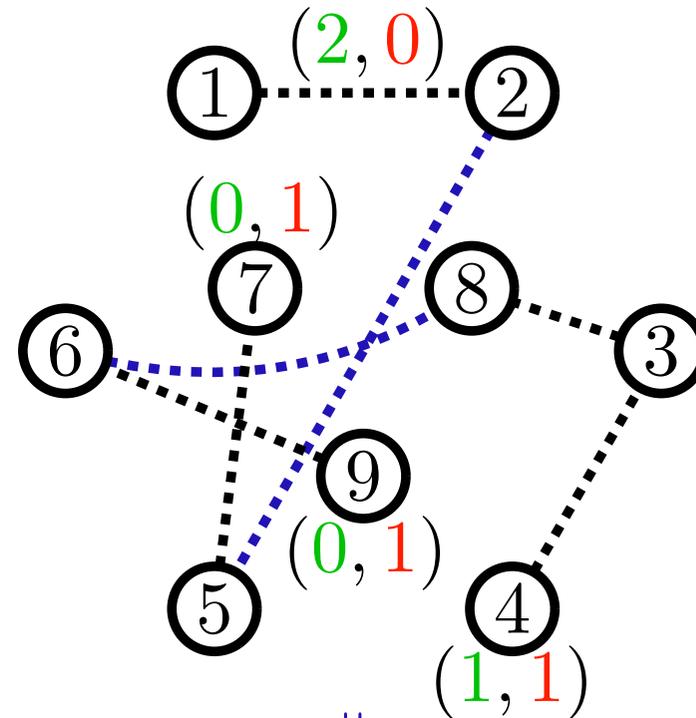
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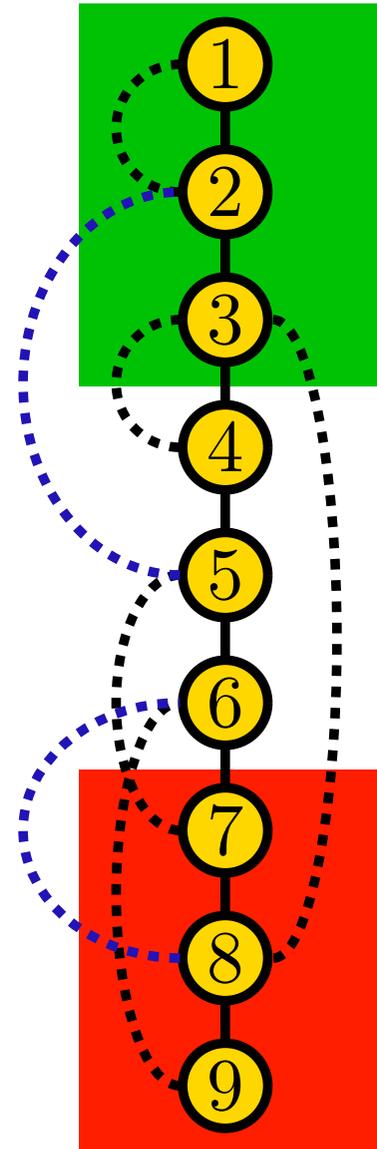
$n^2$  per Matching:

- Quantify Distortion of Matching

$$2 \times 1 + 0 \times 0 = 2$$



$$1 \times 1 + 0 \times 1 = 1$$



# Lower Bound

The Strategy:

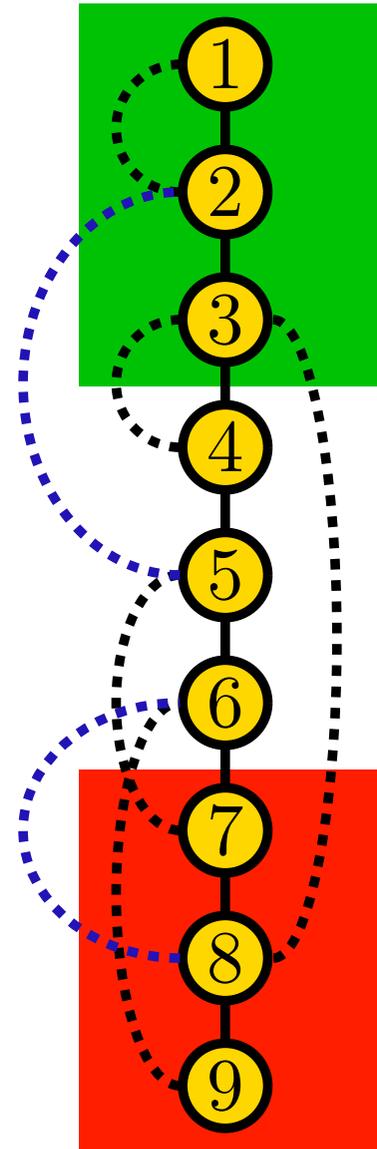
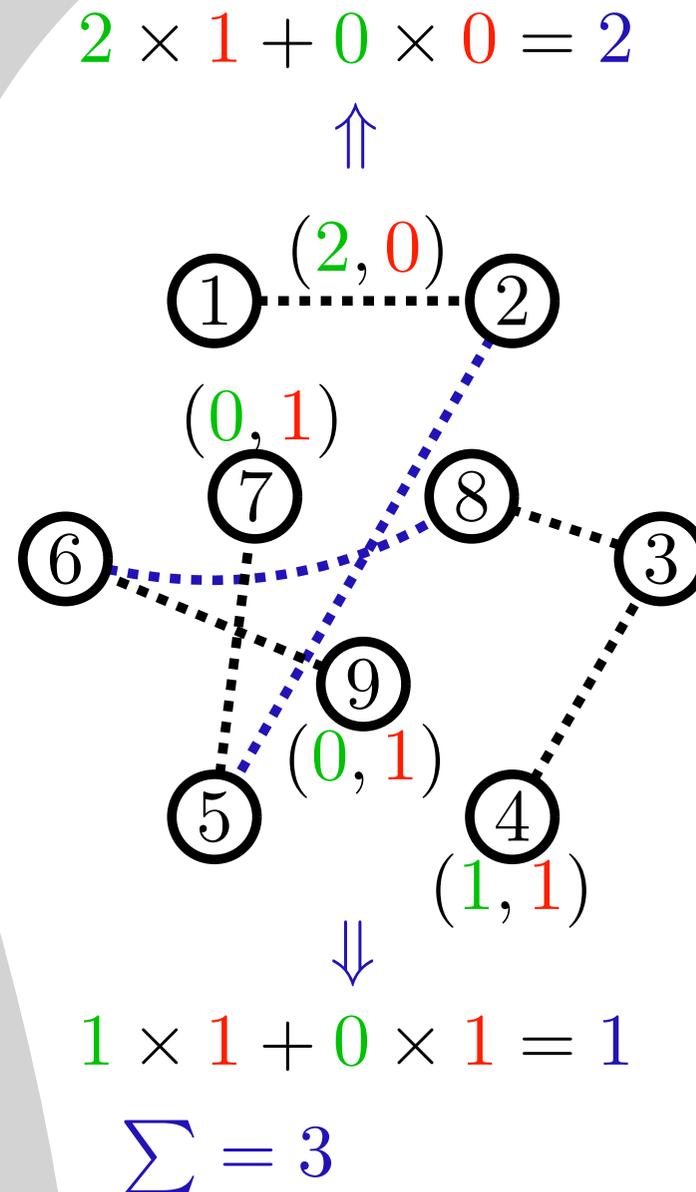
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# Lower Bound

The Strategy:

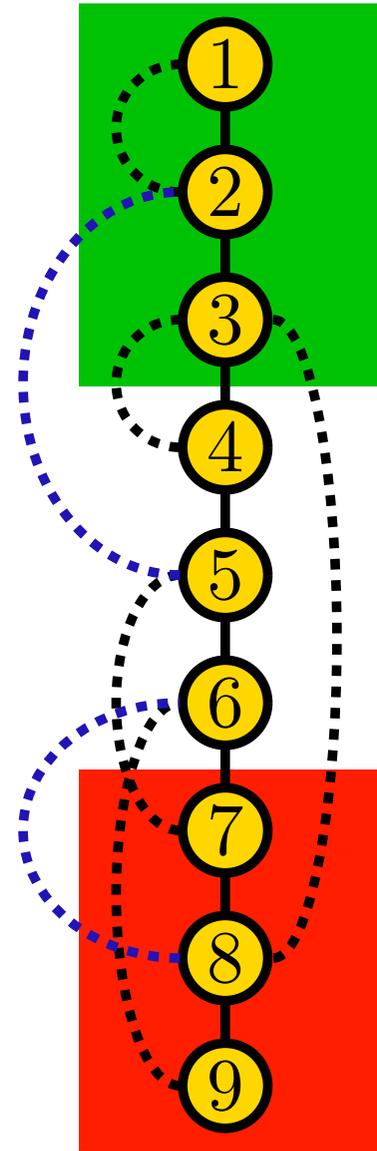
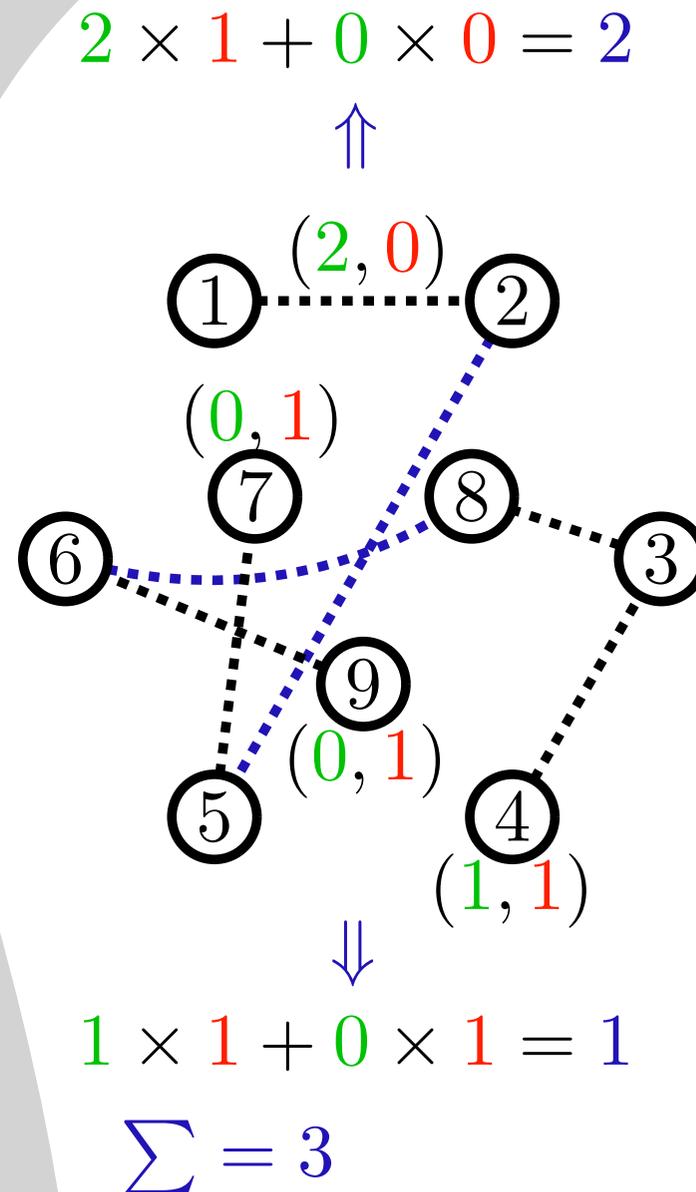
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- Competitive Ratio:  $\Omega(\log n)$

$n^2$  per Matching:

- Quantify Distortion of Matching
- Compute Sum of Distortion of all Matchings
- Average Distortion is Sufficient

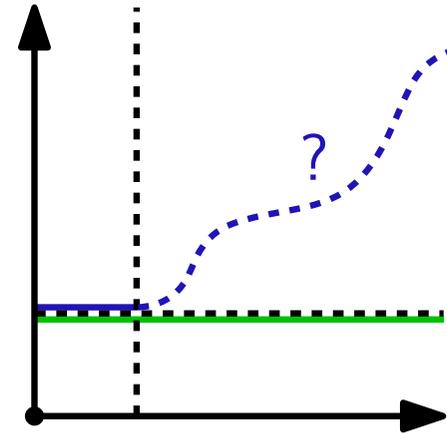


# Open Questions & Future Work

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Better Bounds for Line Networks:

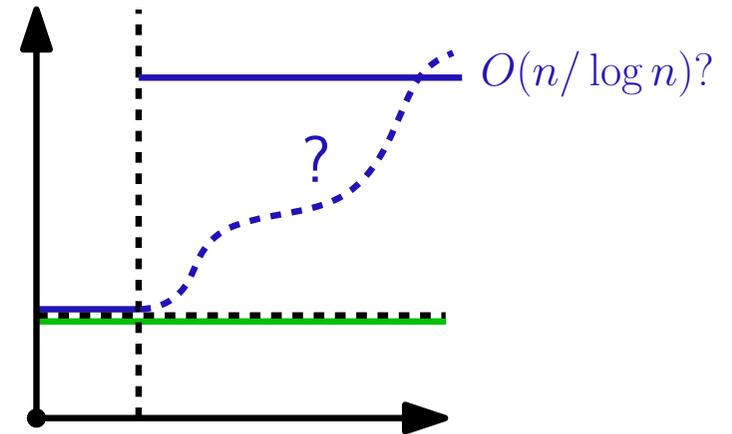
- Nontrivial Upper Bound?
- Better Lower Bound?
- Dynamic Offline Algorithms?



# Open Questions & Future Work

Better Bounds for Line Networks:

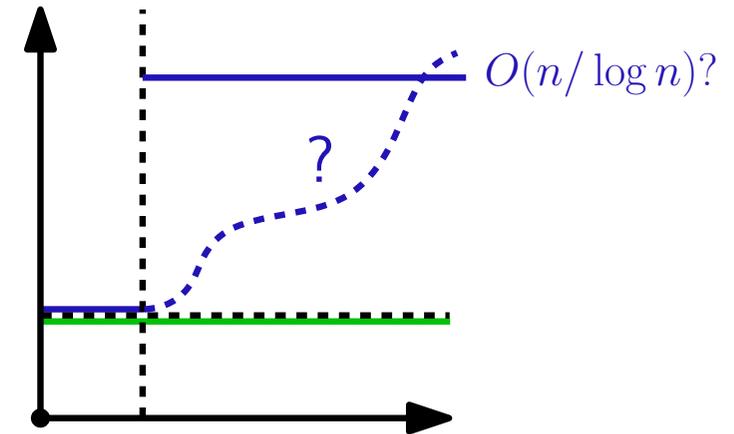
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# Open Questions & Future Work

Better Bounds for Line Networks:

- Nontrivial Upper Bound?
- Better Lower Bound?
- Dynamic Offline Algorithms?



Different Networks:

- Binary (Splay) Trees
- No BST Property

